A NEW FRAMEWORK FOR MINIMISING HANDOVER IN MULTICAST MOBILITY

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Submitted in Partial Fulfilment of the Requirements of the Degree of Doctor of Philosophy, 2015

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ACKNOWLEDGEMENTS

First of all, I would like to deliver my greatest appreciation to my supervisor, Prof. Nigel Linge, for all his support and guidance throughout the research period; without his support this research could not have been completed successfully.

I would also like to thank the Thai government for providing me with the scholarships and living allowances in order to support my quest for a PhD degree. Appreciation would also be given to my office, King Mongkut's University of Technology North Bangkok for granting me the study leave and the opportunity to further my studies in this PhD scholarship.

I would also like to thank all the staff and members at CNTR (Computer Networking and Telecommunications Research) centre for all their support and suggestions. I have really enjoyed my time there with everyone.

Last but not least, I would also like to express my gratitude to my family, parents and friends for their continuous patience and support throughout this period of study.

LIST OF ABBREVIATIONS

3G Third Generation Wireless (Communications)

3GPP 3rd Generation Partnership Project 3GPP2 3rd Generation Partnership Project 2

3GPP-PSS 3rd Generation Partnership Project – Packet Switched Streaming

AP Access Point
AR Access Router

ASM Any Source Multicast

BCMCS Broadcast and Multicast Service

BGMP Border Gateway Multicast Protocol

BSC Base Station Controller

BTS Base Station Transceiver System

BU Binding Update

CBS Cell Broadcast Service

CBT Core-Based Trees

CDMA Code-Division Multiple Access

CN Core Network

CNTR Centre of Networking and Telecommunication Research

CoA Care of Address

CPU Central Processor Units

DMM Distributed Mobility Management

DMSP Designated Multicast Service Provider

DVMRP Distance-Vector Multicast Routing Protocol

EAP-AKA Extensible Authentication Protocol method for UMTS Authentication

and Key Agreement

EoS Economies of Scale

FA Foreign Agent

FMIP6 Fast Mobile IPv6

fps Frames per Second

FSMs Finite State Machines

GGSN Gateway GPRS Support Node

GPRS General Packet Radio Service

GPS Global Positioning System

GSM Global System for Mobile Communications

GTP Generic Tunnelling Protocol

HD High Definition

HDTV High-definition TV

HMIPv6 Hierarchical Mobile IPv6

HoA Home Address

ICMP Internet Control Message Protocol

ICMPv6 Internet Control Message Protocol version 6

IETF Internet Engineering Task Force

IG Integration Gateway

IGMP Internet Group Management Protocol

IGMPv2 Internet Group Management Protocol version 2

IMS Internetworking Management System

IP Internet Protocol

IPv4 Internet Protocol version 4
IPv6 Internet Protocol version 6

ISHO Inter-System Handover

ITU-T International Telecommunications Union - Telecommunication

Standardization Sector

LMA Local Mobility Anchor

MAC Media Access Control

MAG Mobile Access Gateway

MBMS Multimedia Broadcast/Multicast Service

MBONE Multicast Backbone

Mbps Megabit per second

MDP Markov Decision Process

MIH Media Independent Handover

MIPv6 Mobile Internet Protocol version 6

MLD Multicast Listener Discovery

MOM Mobile Multicast Protocol

MOSPF Multicast Open Shortest Path First

MPEG-2 Motion Picture Experts Group 2

MPEG-4 Motion Picture Experts Group 4

MR Multicast Router

MSP Multicast Service Provider

ND Neighbour Discovery

NIA Network Interoperating Agent
NS-2 Network Simulator version 2

OMNET++ Objective Modular Network Testbed in C++

PIM Protocol Independent Multicast

PIM-DM Protocol-Independent Multicast Dense Mode
PIM-SM Protocol-Independent Multicast Sparse Mode

PMIPv6 Proxy Mobile IPv6
PTM Point to Multipoint

PTP Point to Point

QoE Quality of Experience
QoS Quality of Service

RFC Request for Comments

RNC Radio Network Controller

RP Rendezvous Point

RSSI Received Signal Strength Indicator

SAP Session Announcement Protocol

SDTV Standard-definition TV

SGSN Serving GPRS Support Node

SIP Session Initial Protocol

SNMP Simple Network Management Protocol

SPT Shortest Path Tree

SSM Source-Specific Multicast

STP Spanning Tree Protocol

TCP/IP Transmission Control Protocol / Internet Protocol

TV Television

UMTS Universal Mobile Telecommunications System

UTRAN UMTS Terrestrial Radio Access Network

VoD Video on Demand

WiFi Wireless Fidelity

ABSTRACT

Nowadays, mobile devices support a variety of multimedia applications such as live video, radio or online gaming. People spend their time on mobile devices for entertainment more and more via the internet. Due to the requirements of multimedia applications over wireless communication those applications require a huge bandwidth on the network to support them, which creates problems for the network provider. However, one pattern that is appropriate for the efficient delivery of multimedia messages is multicast delivery.

Multicast services do, however, introduce challenges within the network when the recipients of the service are moving. Powerful multicast routing protocols are designed for static client IP addresses. Hence, when the mobile node changes the location, it introduces the problem of access network handover. Therefore, this is the aim of the research where a new framework will be developed for multicast mobility within WiFi network to reduce and provide smooth mobility when handover occurs. This research is focused on techniques to reduce handover latency, minimize packet loss and provide connection when a user moves between network zones.

To achieve these aims, this designed framework lets mobile nodes send the message to register to foreign agents in advance for addressing IP address of the new zone and to establish the multicast tree earlier. Moreover, there are processes that keep the connection of the path alive.

The framework is being simulated on OPNET Modeler for evaluating the performance in terms of handover latency time, the number of packet loss and so on. There are many scenarios that have been tested. According to the results, it shows that the new framework has reduced handover latency time around 60% on average and minimized packet delay approximately 0.7 - 150 ms on mobile node depending on network topology. This framework can provide IP address reconfiguration, binding update, joining multicast group and distribution path of multicast tree in advance. However, there are some overheads and cost that this framework has to pay for such as IP address database, increasing broadcast within networks and keeping connection path alive.

Chapter 1 Introduction

1.1 Background

Today's mobile device supports a variety of multimedia applications such as live video, radio or online gaming. People spend their time on their devices for personal entertainment more and more. There are statistics which show that on average, people are spending 2.7 hours per day on the internet via smart phone [1].

The smart phones are characterized by small screens, limited CPU power and memory. Due to the requirements of multimedia applications over wireless communication those applications require a huge amount of bandwidth from the network to support them, which creates problems for the network provider and also affects to QoE (Quality of Experience) of end user for multimedia services.

However, one pattern that is appropriate for the efficient delivery of multimedia messages is multicast. For instance, transmitting live video data from a media server using multicast allows for a single data stream to be simultaneously sent to several users thereby offering a considerable bandwidth saving over sending multiple separate data streams.

Multicast services do, however, introduce challenges within the network when the recipients of the service are moving. Current multicast routing protocol standards are designed for static client IP addresses. Not only does mobile node movement introduce the problem of access network handover, but also, when considering devices could result in them switching between access network hotspots. Support for this requires an efficient solution to be found for managing multicast services where mobile devices have the capability of operating over several different networks.

Therefore, this research is investigating multicast mobility handover within WiFi networks. The aim of this PhD research is to propose a new method and framework that tries to provide smooth mobility in homogeneous networks such as within WiFi zones.

1.2 Research Aim

The research has the aim to propose a new technique that provides efficient mobility for multicast services in WiFi networks. This research is focused on techniques to reduce handover latency, multicast address management and provide a connection when a user moves between networks.

This research is focused on techniques to reduce handover latency, which comprises multicast handover delay, end-to-end delivery delay and to minimize packet loss. For multicast routing, the research will focus on how to maintain the multicast session for mobile nodes and manage multicast group memberships thereby also minimising packet delay.

1.3 Research Objectives

To achieve this research aim, four objectives have been acknowledged. These research objectives are:

- 1. To understand, analyse and identify major issues in multicast mobility and focus on issues, which are related to the handover process.
- 2. To propose the technique that can support multicast mobility when a mobile node changes the zone within WiFi network environments.
- 3. Implement and evaluate, within a simulation environment, that part of the framework that handles WiFi zone handover.
- 4. To publish the results of the research and write up a PhD thesis.

1.4 Research Questions

During the research process, the following key research questions will be addressed:

- 1. How can we reduce latency time in a multicast stream when handover occurs?
- 2. What performance advantages can be gained over existing schemes for mobile multicast delivery?

1.5 Contribution to Knowledge

Designing a new framework for minimizing multicast mobility within WiFi networks by focusing on techniques to reduce handover latency, packet delay and maintain multicast services for mobile nodes. The key protocols that had been modified in the framework are Mobile IP, IGMP, PIM and ICMP protocols. The key novel features of this new framework are that firstly, mobile nodes register with each foreign network that is within range, secondly a Care of Addresses is obtained for each foreign network, thirdly that the associated multicast trees are established but finally, that the chosen foreign network and multicast tree is only activated once handover is confirmed. In this way, the new framework is able to minimise handover delay and loss of connectivity when handover is taking place.

1.6 Overview of Research Methodology

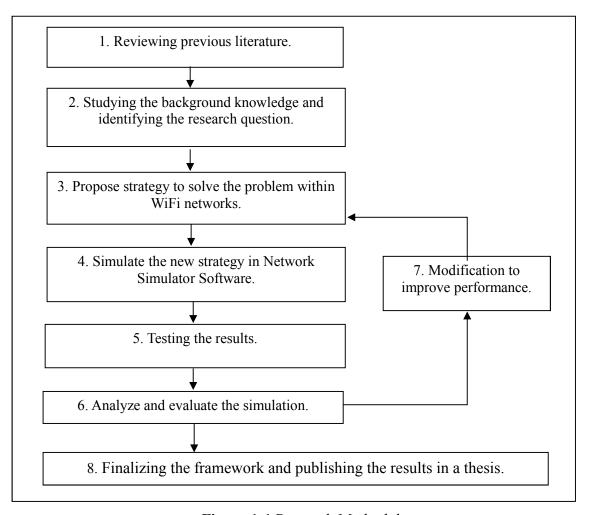


Figure 1-1 Research Methodology

In order to complete this research, the following research methodology has been adopted as shown in Figure 1-1.

The details of each step are described as follows:

- 1. Reviewing previous literature More than 150 research articles have been reviewed. At this stage, the researcher has gained more knowledge about the problems in multicast mobility, and the mechanism of handover.
- 2. Studying the background knowledge and identifying the research question The literature review confirms the importance and relevance of the research field and shows where gaps exist within the knowledge. These gaps then help to formulate the research question for this research.
- 3. Propose strategy to solve the problem within WiFi networks From the background knowledge, literature review, and research question it is then possible to start to propose and design a new framework for improving the performance and QoS guarantees of multicast mobility.
- 4. Simulate the new strategy in network simulator software The first stage of the framework will be evaluated for use within WiFi networks only. This allows the basic concepts of the framework to be evaluated and as necessary, modified.
- 5. Testing the results The detailed simulation results will allow the performance of the framework to be quantified and areas for improvement identified across a broad range of networking scenarios.
- 6. *Analyze and evaluate the simulation* The outputs from the simulation will be analysed in detail to assess performance and to compare to standard framework.
- 7. *Modification to improve performance* Depending upon the simulation results, further refinement and modification of the scheme will follow as necessary.
- 8. Finalizing the framework and publishing the results in a thesis The results of all of the simulation studies will allow for a final design to be confirmed and this, together with the contribution to knowledge of the research, will be presented within a PhD thesis.

1.7 Structure of the Thesis

The rest of the thesis is organized in the following chapters.

- ➤ Chapter 2 presents the literature review concentrating on the concept of the multicast, multicast mobility, multicast handover issues within wireless networks.
- ➤ Chapter 3 presents the concept and theory idea design of the research, which is to combine network architecture, protocol overview, connection management and modified protocol message.
- ➤ Chapter 4 consists of framework simulation details which have been implemented within OPNET Modeler.
- ➤ Chapter 5 will present the network scenario, simulation result from OPNET Modeler software and performance evaluation.
- **Chapter 6** consists of conclusion and recommendations for further research.

Chapter 2 Literature Survey

2.1 Introduction

This chapter provides an introduction and a summary of the key literature that has been consulted for this research. Creating a framework to provide multicast services for mobile nodes in a wireless communication network is a challenging issue.

The relevant published research has been surveyed in the fields of multicast delivery, multicast mobility and multicast handover within WiFi networks. This framework can support both IPv4 and IPv6 WiFi environments, so the key protocol information and literature review in this chapter will be discussed and focuses on IPv4 and IPv6 environments.

This chapter begins with an overview of multicast delivery concept and protocols. It will then go on to a review of multicast mobility protocols and problems within WiFi networks. After that, the processes and problems of multicast handover are examined. Following this, details about multicast handover issues and approaches will be investigated. Finally, a brief summary is given for the chapter.

2.2 Multicast Delivery

In the IP network system there are three types of communication in a network:

- ➤ Unicast Delivery one source is sending a single transmission of data directly to one destination.
- ➤ **Broadcast Delivery** one source is sending the data to all destinations in the network.
- ➤ Multicast Delivery one source is sending the data to a select group of destinations.

In traditional computer networks, data is typically sent from a source node to a destination node known as unicast delivery, which is suitable for most applications in the network; alternatively, one source node can transmit a copy data to all end point destination nodes and then the destination nodes will decide if they want to use those data or not known as broadcast delivery. One benefit of broadcasting is that it reduces loads at source node from duplicating data that are sent for multiple destination nodes. The source node sends only one copy of the data to the broadcast address and then the network devices on the network will duplicate the data and transmit to cover the network. However, there are some kinds of application that need to send the same data to multiple destination nodes such as multipoint videoconferencing, online gaming and live TV. It will waste bandwidth on the network if it uses the process of unicast delivery to support those applications. Hence, the aim of multicast delivery is to deliver data packets between one source to multiple destinations more efficiently.

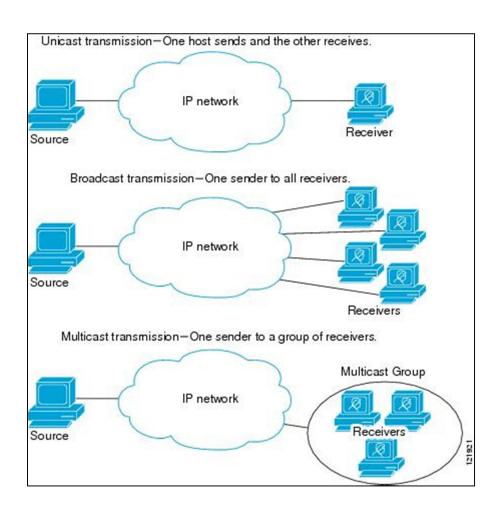


Figure 2-1 Comparison between Unicast, Broadcast and Multicast transmission [2]

The concept of multicast delivery was introduced in the late 1980s by Stephen Deering [3, 4]. The idea is to transmit a single copy of IP packet to a group of destinations, which is identified by a same multicast IP address. A main factor in multicasting is bandwidth efficiency in the network. The multicast functions and protocols have evolved over time as refined in RFC 3376 [5], RFC 4604 [6] and so on.

Multicast is the delivery of data packet to a group of user devices using a common IP multicast destination address. When a multicast tree is set up, the source starts sending IP datagram to the host group address. Then, the network devices take on the responsibility for sending the IP datagram to all destinations within multicast group. Multicast routers along the path are responsible for ensuring that datagrams are transmitted over the appropriate links to ensure they reach all hosts of the multicast group.

The process of coping datagrams occurs only when paths diverge at a multicast router, thus reducing the bandwidth consumption on the network. Figure 2-2 is shown an example of multicast delivery in a TCP/IP (Transmission Control Protocol / Internet Protocol) network. The arrows represent the direction of multicast delivery packets that are sent to host B, host D and host E in the network.

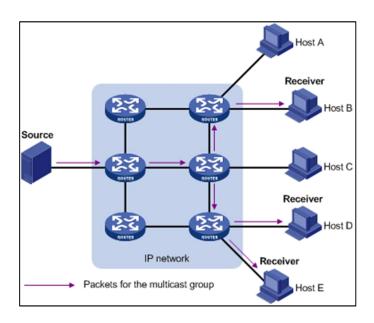


Figure 2-2 Multicast delivery [7]

In networking, the term multicast is synonymous with IP multicast. Multicast delivery is a technique for delivery datagrams from one source to multiple destinations over the network. The optimal distribution paths are created after the member in the multicast group joins in the multicast tree in real time. Multicast delivery technique and group management can support a large number of destination nodes without needing to know how many destination nodes there are. Multicast group requires the source node to send a request to join message only once, even if there are a large number of receive nodes to be sent. IP multicast has an efficient process to maintain the members within the multicast group.

The model of the multicast group is known as an open and dynamic group. This means that, the source node can send multicast packets at any time when it is ready, with no need to announce, register or schedule transmission. The only one thing that the source node needs to know is a multicast address. Also, it is not necessary for the source node to know about group membership in advance. The host members in the multicast group can join or leave a multicast group at any time. However, source and destination multicast nodes can communicate to each other via IP multicast group address. Sources use the multicast IP address as the destination address in their sending data stream. Destination hosts use multicast IP address to notify the network device that they want to receive those packets that are sent to this multicast tree. However, the destination host has to send "join message" to be a member of the multicast group first.

In order to deliver multicast data stream, the network creates *a "multicast tree"*. The multicast tree construction is begun with network devices close to the destination nodes and is thus receiver-driven. The multicast tree is built for that group once there are members in a particular IP multicast group and maintained by "multicast routing protocol". There are several kinds of multicast routing protocol depending on the network. Also, each one of them has its own unique method and technique. More information about multicast routing protocols will be given in a later section.

2.2.1 Benefits of Multicast Delivery

The main benefit of multicast delivery is a reduced bandwidth requirement on the network, because this technique sends a single copy of data stream along any link that forms the route.

Live multimedia services such as video, radio, television etc., are delivered as multimedia streaming services. These services require large bandwidth allocations and if they are delivered as separate streams for each user, this will place huge demands on network resources. Hence, multicast delivery can help to control network traffic and reduce this problem [8]. This advantage is known as optimized performance, because multicast delivery eliminates traffic redundancy on the network. Figure 2-3 illustrates a multicast tree, which is represented in Figure 2-2.

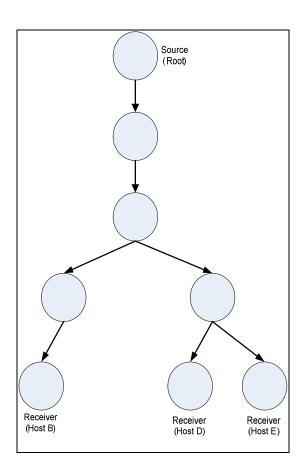


Figure 2-3 Multicast Tree

Bandwidth requirements for multicast delivery are widely dependent on applications such as news, which requires around 10 kbps as a lightweight data. However, it will be up to 384 kbps for delivery of video content and video streaming. Similarly, audio streaming will require a bandwidth of 48 to 64 kbps [9, 10], while low-quality video with audio will require a bandwidth up to 128 kbps [9].

Moreover, in 1994, a host in Japan was found to require a bandwidth over the entire MBONE (Multicast Backbone) of 650 kbps for video streaming. However, if compared with the bandwidth requirement for digital TV, there is a big difference. SDTV (Standard-definition TV), with MPEG-2 compression, will require 2-5 Mbps and 1.5-2 Mbps with MPEG-4 compression. On the other hand, the bandwidth requirement for HDTV (High-definition TV) is increased about five times, 15-20 Mbps with MPEG-2 compression and 5-10 Mbps with MPEG-4 compression [11].

In addition, another advantage of multicast is enhanced efficiency by reducing loads on central media servers in terms of CPU (Central Processor Units) power, memory usage and protocol management. Network routers also only need to manage and maintain one data stream per multicast service [12]. However, routers do have the added burden of needing to maintain knowledge of which devices belong to which multicast tree and to manage routing along these trees.

By using multicast delivery, video streaming application can offer a higher quality service because the load for single multicast delivery is less than multiple unicast delivery [13].

The network bandwidth saving is the main point expected of multicast's effect on the network. Hence, network providers normally must consider the cost of managing, maintaining and implementing multicast. However, the main factor is bandwidth. R. Chalmers and K. Almeroth [14] defined a metric to compare the performance between unicast and multicast delivery. The metric compares the total number of links traversed by unicast and multicast datagrams on a given network infrastructure. The metric refers to each links as a hop in the route path of a single multicast or unicast datagram.

$$\delta = 1 - \frac{\text{multicast hops}}{\text{unicast hops}} \tag{2.1}$$

Where

Multicast hops are the total number of multicast links in the distribution link in the network.

Unicast hops are the total number of unicast hops in the network.

 δ is the multicast metric which is a fraction in the range $0 \le \delta \le 1$.

 δ marks the percentage increase in the bandwidth utilization achieved by using multicast more than unicast delivery. For this metric if the value of δ is zero, this means the number of hops is the same. If the δ has a value of nearly one, it means the performance of using multicast delivery is higher compared to unicast.

Due to the number of multicast members being dynamic, new nodes and destinations can join and leave the multicast group at will. This means the multicast group size always changes over time, especially in a real time application. J. Chuang and M. Sirbu [15] presented a cost function of path, which is related to equation 2.1 and defines a direct relationship between the hop counts and the size of multicast group.

$$\frac{L_m}{\bar{L}_u} = N^k \tag{2.2}$$

Where

 L_m is the total multicast distribution tree length

 \bar{L}_u is the average length of unicast routing path

N is the size of multicast group

k is the economies of scale (EoS) factor.

The range of value k is between 0 and 1, which is the value of the slope of the graph relationship between normalized tree cost and multicast group size. For most of the topologies investigated in their experiment, which is in real and generated networks, it is shown that $k \approx 0.8$.

Now assuming $\bar{L}_u = \frac{L_m}{N}$ will get:

$$\delta = 1 - N^{\varepsilon} \tag{2.3}$$

Where

$$\varepsilon = k - 1 \approx 0.8 - 1 \approx -0.2 \tag{2.4}$$

So

$$\delta \approx 1 - N^{-0.2} \tag{2.5}$$

Equation 2.5 gives us an estimate for the multicast performance, which is based on the number of destination nodes in the multicast group.

Another advantage of multicast is distributed application, as multicast makes multipoint applications possible. The example applications of multicast services are provided in the following.

- ➤ Multimedia Conferencing
- > Online video/audio streaming
- ➤ Interactive distance learning
- ➤ Online TV
- > Group online gaming
- ➤ Video on Demand (VoD)
- ➤ Commercial stock exchanges

However, a major problem with this kind of application is the lack of reliable delivery of data because multicast delivery is UDP protocol based, not TCP protocol. There are some multicast disadvantages such as [13]:

- *Best Effort Delivery*: dropping packets are to be expected. Therefore, multicast applications should be designed accordingly and should not expect high reliability of packet transmission. Packet losing should be accepted on those applications.
- Duplicate Packets: some multicast protocol techniques such as registers, asserts and STP transitions may result in the occasional production of duplicate datagrams.

- Same Video Stream: only one copy of multicast stream is sent to all users in the same multicast group, so every user receives the same data stream at the same time. Hence, individual users cannot choose the content they want, and also they cannot pause the video stream, rewind it or skip some parts.
- > Specific content: in the multicast network, it is more complicated to control the users' access to specific streaming content because all users in the same group access the data at the same time.
- > Out of Order Delivery: as multicast is based on UDP protocol, some routing protocol techniques may also result in packets being out of order. Hence, multicast applications should be designed to solve the issue of out of order packets arriving.

2.2.2 Functions of Multicast Delivery

This section provides a brief overview of the important functions of multicast delivery.

- ➤ *Multicast Address Management:* This function is about the assignment and the scope of multicast address.
- Multicast Service Announcement and Discovery: These services allow destination hosts to discover the availability of multicast source. In TCP/IP network, the SAP (Session Announcement Protocol) protocol handles this function [16].
- ➤ Multicast Group Management: Handles the collection and maintenance members of the multicast group. In IPv4 network, IGMPv2 (Internet Group Management Protocol version 2) protocol deals with this function. The counterpart in IPv6 network is MLD (Multicast Listener Discovery) protocol.
- Multicast Routing: Multicast Routing protocols are responsible for the multicast tree for example to build and maintenance the multicast trees. Also they are connect multicast group members and for the forwarding of data stream on the multicast tree. A number of multicast routing protocols have been standardized such as MOSPF (Multicast Open Shortest

Path First), PIM-DM (Protocol-Independent Multicast Dense Mode), PIM-SM (Protocol-Independent Multicast Sparse Mode), DVMRP (Distance-Vector Multicast Routing Protocol) and Core-Based Trees (CBT) [17].

- **Reliable Multicast Transport:** This function is to ensure the reliable delivery of multicast stream to a potentially large destination group.
- Multicast Mobility: Multicast delivery mechanism designs for static receiver. However, the mobile source and destination moves create issues for delivering the multicast data such as how to keep connection, managing IP address in a new location and so on. Solving these problems is a big challenge. This research has investigated this function in some depth and proposed a new framework to achieve the required functionality of this function.

According to multicast delivery, it is not only the way to send one message source to many destinations but also it can send from many sources to many destinations. Moreover, in a mobile environment this means the source could possibly move. In terms of a mobile source, these are divided into Any Source Multicast (ASM) and Source-Specific Multicast (SSM) [18].

ASM is where a mobile node submits data to any source in a multicast tree via an ASM group and either creates the root of a source-specific shortest path tree (SPT) forwarding datagram to a rendezvous point (RP) or destinations, or it distributes packets directly down a shared tree.

Source-Specific Multicast or SSM has been designed for multicast sources with static source addresses. Normally, the source addresses in a mobile node subscribed to an SSM group are directly used for route identification. However, the address of the SSM source possibly changes under mobility. So, mobile node implementations of SSM source filtering must be MIPv6 (Mobile Internet Protocol version 6) aware in the sense that a logical source identifier, which is HoA (Home Address), is correctly mapped to its current location represented by CoA (Care of Address) address.

2.3 Multicast Protocols in IP Networks

The aim of this section is to review the standard protocols, which are related to the multicast services within the network. In a TCP/IP network, the protocol responsible for multicast group management is MLD protocol for IPv6 network and IGMP protocol for IPv4 network. Hence, in the next section MLD protocol will be the focus.

Multicast routing protocols use algorithms to build multicast distribution trees. Each routing protocol is different in regards to how they share content, information and create paths [19]. Figure 2-4 is shown the classification of the intra-domain multicast protocols.

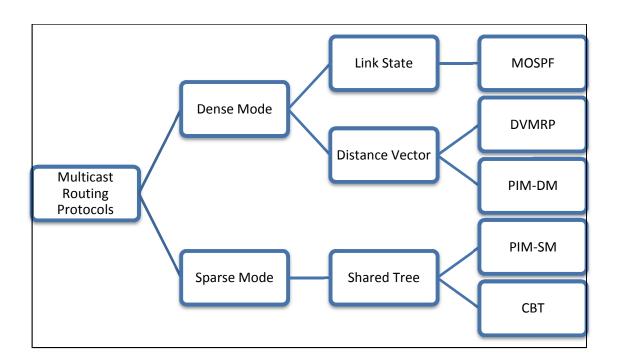


Figure 2-4 Intra-domain routing protocols [20]

A brief summary of these multicast routing protocols are as follows:

MOSPF: This protocol is an extension of the OSPF routing protocol in unicast delivery. The MOSPF protocol uses the link-state algorithm the same as OSPF to create the paths with minimum protocol traffic overhead [20]. Moreover, if the number of multicast groups and group member size increase the computational complexity of the link-state algorithm will increase also [21].

- ➤ **DVMRP:** This protocol defined in 1988 by Waitzman, Partridge and Deering, was the first multicast routing protocol, which was deployed over the MBONE network [22]. This protocol is an extension of the RIP routing protocol in unicast delivery and extends to multicast network [21].
- ➤ **PIM:** It is a popular multicast routing protocol standard in TCP/IP network. This research has modified the protocol in the framework. In the following section PIM protocol will be discussed in detail.
- > **CBT:** This protocol was defined in 1997 [23]; it is a shared tree protocol and cannot create a source-based tree.

2.3.1 MLD Protocol

Multicast data delivery comprises both local and global techniques in which local techniques are responsible for multicast group management. Global techniques are responsible for multicast routing. For LAN multicasting each node can choose whether it wants to receive multicast data or not. Multicast destination nodes will inform the network device that they want to receive data packets that are sent to the multicast group.

This process is called a "join" and is controlled by the IGMP protocol in an IPv4 network and controlled by the MLD protocol in an IPv6 network. In this research, global mechanisms are managed by the PIM protocol (Protocol Independent Multicast), which is multicast routing protocol in intra-domain network. For local mechanisms we chose to work with MLD protocol [24].

The MLD protocol was defined in 2004 by [25]. There are two versions of the MLD protocol, which are MLD version 1 (MLDv1) and version 2 (MLDv2). In this research we concentrate on the MLDv2 protocol, which is used by IPv6 routers to discover the receivers who wanted to join in a multicast tree [12].

Receiver nodes on a route keep an interface state for every IPv6 multicast address from which they want to receive data stream per interface [26]. The interface state on a receiver node contains a record including a filter mode and also source list for each multicast IP address [20].

2.3.2 PIM Protocol

To support multicast communication, a multicast routing protocol is required. For this research we consider the PIM protocol because PIM is one of the most popular shared tree multicast routing protocols. There are two types of PIM protocol, namely PIM-DM, which was defined in 2005 [27] and used in environments where multicast trees are populated densely within the network. Another type is PIM-SM defined in 1998, which is better suited to sparsely populated networks.

In this research we are focusing on the PIM-SM protocol because it is designed to support large region networks such as the internet. Sparse mode is activated when multicast groups are thinly populated across a large network region. This mode is designed for that situation. However, PIM-DM builds a separate source-based tree for every source, while a shared tree has been used for all sources within a multicast group.

All multicast sources in the tree transmit all multicast stream traffic to the root, and then the root forwards the multicast traffic to all destination nodes in the network. Normally, multicast sources encapsulate their multicast data in unicast data packets addressed to the RP router within the multicast tree. When the RP receives those packets, the RP will decapsulate these packets and then forward them over the multicast delivery tree to all members in the multicast group [20].

Normally, destination hosts join a multicast tree by sending a join message towards the RP. The routers on the route towards the RP will store status information for the multicast group while passing the join request to the RP router, thereby building the multicast delivery tree in the direction of the new destination.

There are many control messages, which are used within the PIM-SM protocol. The PIM-SM message can show the following parts:

- Bootstrap message
- > PIM-SM control-message encapsulation
- > Hello message
- > PIM-SM packet header
- Encoded Unicast Address field
- > Join/Prune message
- > Encoded Group Address field
- > Candidate RP advertisement
- Encoded Source Address field
- Register message
- > Assert message
- > Register-Stop message

2.4 Multicast Mobility in WiFi Network

In this section, a review of mobility within the wireless network will be presented, including multicast mobile and related protocols.

2.4.1 Overview of WiFi Network

Wireless networks support different data rates and coverage area sizes. For example, IEEE802.11g supports a data rate of 54 Mbps but GPRS on 3G in practice only manages a data rate of around 9.6-384 kbps, such as specifying a minimum data rate of 144 kbps high mobility application like a car, while slow mobility like a pedestrian application requires 384 kbps and up to 2 Mbps stationary applications. Typical network data rates are summarised in Table 2-1 [28].

Table 2-1 A data rate and coverage area of wireless technologies [28]

Network type	Frequency	Data Rate	Coverage
Bluetooth	2.4 GHz ISM	Max 721	0.1-10m
Biuctootii	band	kbps	0.1-10111
IEEE 802.11a	5 GHz	20 Mbps	50-300 m
IEEE 802.11b	2.4 GHz	11 Mbps	Up to 100 m
IEEE 802.11g	2.4 GHz	54 Mbps	30-150 m
IEEE 802.16	10.66.CH-	Max 70	Oxyan 50 Irms
(WiMAX)	10-66 GHz	Mbps	Over 50 km
IMT2000, UMTS	2 GHz	Max 2	30m-20km
IM12000, UM15	2 GHZ	Mbps	30III-20KIII
GPRS (GSM),	900, 1800,	9.6-384	Un to 25 lem
EDGE (HSCSD)	1900 MHz	kbps	Up to 35 km

2.4.2 Mobility in WiFi Networks

In IPv6 wireless network, the protocol that provides mobility support is Mobile IPv6 protocol, and this protocol will maintain the connection when the mobile node moves. Also, it includes the responsibility for the reachability of the mobile node and network, keeping records and IP address. So, in the next section the protocols that are related with mobility will be presented.

2.4.2.1 Mobile IPv6 Protocol

The mobile IPv6 protocol has been proposed by the IETF (Internet Engineering Task Force) to provide mobile support for hosts within IPv6 networks. In Mobile IP, a mobile node uses two IP addresses that are its home address and a care-of-address. The home address is a stable IP address assigned to a device and based on their point of network connection [29].

In contrast, the care-of-address is a temporary address provided for a 'foreign' network and this address will change as the device moves between different IP subnetworks. The care-of-address in IPv6 network can be formed based on stateless or stateful mechanisms [30]. When the mobile node moves, it first forms a care-of-address based on the prefix of the foreign link. Then, the mobile node sends a Binding Update (BU) message to the home agent, which is its temporary care-of-address (CoA).

After that, although the home agent wants to grant or deny the request from the mobile node, the home agent will send a registration message reply [31]. After the registration process is successful, any messages destined for the mobile node are intercepted by the home agent, which encapsulates the packets and tunnels them to the foreign agent. Then, the data streams are forwarded to the mobile node [32].

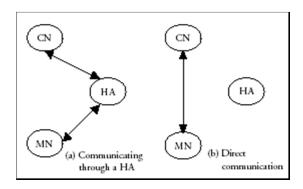


Figure 2-5 Mobile IPv6 Protocol

Mobile IPv6 offers improvements in this process compared to Mobile IPv4. For example, Mobile IPv6 can eliminate the triangular routing issue and produces route optimization. Route optimization is the process that enables the correspondent node to reroute data stream on a direct path to mobile destination [33].

Moreover, Mobile IPv6 includes embedded binding updates for the home agent and care-of-address configuration for location updates. Additional security has also been incorporated such as authentication header processing to provide validation for mobile nodes [19].

2.4.2.2 ICMPv6 Protocol

The ICMPv6 (Internet Control Message Protocol version 6) protocol [34] is used to carry IP control messages for various purposes such as destination unreachable, time exceeded and parameter problem. In addition, ICMPv6 is defined to carry information between hosts, between routers, or between hosts and routers.

There are two types of ICMPv6 message [35]:

> ICMP error message

The functions of ICMPv6 error messages are to report forwarding or delivery errors by either a router or the destination node. The ICMPv6 error messages consist of destination unreachable, packet too big, and parameter problem and so on.

> ICMP informational message

Informational messages provide simple diagnostic functions such as echo request, echo reply, and additional host functionality for example MLD and ND (Neighbour Discovery), which is a set of processes and messages that determine relationships between neighbouring devices.

2.4.3 Multicast Mobility in IPv6 WiFi Network

When a mobile node moves from one network to another network, it is a challenging problem to maintain reachability and transparency of a mobile node. In case of multicast wireless network, the scenario of handover is particularly challenging and serval issues emerge with most solutions due to the handover impacts.

The main problem when the receivers move is multicast latency problem. *Multicast latency problem*, whenever an MN moves to a foreign network, the delay experienced by executing handover process.

Moreover, if a Home Agent (HA) does not support multicast router functionalities, an MN have to discover other a multicast router (MR) and send notification by using MLD

(Multicast listener discovery) query/report messages including a request to join a multicast tree. In this case, the maximum query period of MLD is up to 125 seconds [36]. For some applications, this increased latency time is undesirable such as video conferencing.

Run-liu, W and Yun-hui [37] proposed a multicast routing algorithm trying to reduce the cost of created multicast tree combines with Mobile IP. Also this algorithm tries to reduce the bandwidth of multicast data stream. For applying to large scale of wireless network and reducing join delay time in handover process.

Holbrook, Cain and Haberman [38] proposed a new approach called the Mobile Multicast Protocol (MOM) in 2003. This approach introduces a new entity called the designated multicast service provider (DMSP) and uses a foreign agent entity. The main issue of this research is to reduce the duplicated multicast packets on home agent. However, it created the problem between the HA and FA networks.

Figueiredo, Jeon and Aguiar [39] proposed the solution to reduce vertical handover by adapting a cross layer approach and IEEE 802.21 Media-Independent Handover. There is the process of selected FA links in the network. Also the process uses the single tunnel for completing the delivery process.

2.4.4 Multicast Mobility in UMTS Network

For 3G (UMTS) mobile networks, multicast can be delivered through IP multicast, MBMS (Multimedia Broadcast/Multicast Service) and CBS (Cell Broadcast Service). CBS is a standard that allows the delivery of messages to multiple users in both GSM and UMTS networks. Similarly, MBMS is a standard that is designed to support efficient multimedia broadcast and multicast delivery in GPRS and UMTS networks.

With MBMS, there is provision for streaming services for the delivery of continuous multimedia data traffic. For CDMA2000 networks, there is the BCMCS standard for managing the capability to deliver broadcast and multicast services [20].

2.5 Multicast Mobility Problems in WiFi Networks

2.5.1 Multicast Mobility Problems

Although the research area of multicast mobility has been a concern for about 10 years, there are numerous proposals but not yet a generally accepted standard solution. One reason for this is that the standard multicast protocol was designed for stationary nodes and not for mobility. The problem and challenge of multicast mobility can be divided into 3 categories [40]:

- > *Multicast routing problems*: due to the movement of mobile nodes and the source there are routing problems, such as:
- *Network inactivity*: means when the foreign network the mobile node visited does not support multicast delivery. Hence, the mobile node has to stop receiving the multicast message.
- *Core placement*: when the mobile node moves to a foreign network, the new route will be established. If the mobile node changes zone more often, the frequent handovers can lead to a situation that those multicast routers are off centre. That results in a possible non-optimality configured path.
- *Multicast Encapsulation/Decapsulation:* a variety of methods are used for tunnels to keep connection between the mobile node and home network when the mobile node moves to foreign network.
- *Mobile receiver problems*: In multicast routing when the MNs are moving, they need both efficient IP mobility management and multicast mechanisms to provide a seamless service. Hence, multicast mobility has many constraints that should be considered. It can be classified into:

- Packet loss: a mobile node may miss some multicast data because when it moves normal multicast packets will continue to be delivered to the home network for a short period of time.
- Packet duplication: happens when the mobile node receives the packet from a different multicast router but from the same multicast tree.
 - Packet out of order: when handover occurs.
- Tunnel convergence problem: the tunnel convergence problem is concerned with the delivery of multicast packets to mobile member nodes that are located in several foreign networks over bi-directional tunnels using Mobile IP. An MN receives multicast packets from a home agent, duplicated multicast packets are transmitted to over several tunnels and these become a problem [36].

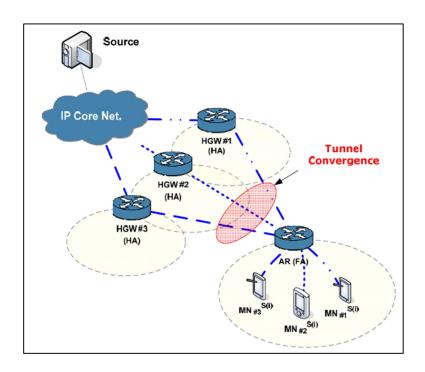


Figure 2-6 Tunnel convergence problem [36]

Mobile source problems: In multicast mobility, it is not only the mobile node that can move. Multicast sources can also move. Consequently, the following aspects must be

considered:

o *Transparency:* is a major problem when a multicast source moves. There is a problem with the CoA of the multicast source, because when a mobile source uses a

new CoA as a multicast source address, it cannot send multicast messages

immediately. It has to wait until the mobile node explicitly notifies that CoA[41].

• Reverse path forwarding (RPF): because it is specific to source location.

Packet loss.

O Source active problem: because the multicast tree has to be reconstructed

when the mobile source moves to the new network.

2.5.2 Overview of Handover Problem

Currently, there are several wireless technologies to provide ubiquitous information access to

users when they are moving. A mobile device such as a smart phone offers multiple wireless

network interfaces and can access these as it moves between different network environments.

Typically, when a mobile device is within the building they receive signals from WiFi and 3G

at the same time but choose WiFi for data services. However, when they move out of the

building, they will be receiving only the 3G signal. The smart phone will, of course,

automatically switch between these networks when WiFi connectivity is lost.

For wireless networks, there are two types of handover that can occur in the network.

Horizontal handover

Horizontal handover is the process by which mobile devices switch from one cell to another

cell within the same network technology; this will be called "intra-system handover".

Vertical handover

26

For vertical handover or "inter-system handover" the mobile device is switching between different network technologies; for example between 3G and WiFi. That means that vertical handover is different in several aspects such as data rate, bandwidth and frequency of operation.

Moreover, vertical handover can be divided into two sub types, which are upward-vertical handover and downward vertical handover [20]. For example, the device moves from WiFi to 3G. Downward vertical handover is a handover that disconnects from a cell providing broader coverage to a wireless overlay with a smaller cell size, and generally higher bandwidth per unit area such as from 3G to WiFi.

2.5.2.1 The Handover Process

One of the major problems of multicast mobility is the handover process which occurs between cells or between different technologies. The handover process can be divided into three phases [8],[28].

Network Discovery or System Discovery: This phase is where the mobile node (MN) searches for reachable wireless networks, usually based upon values of received signal strength.

Handover decision: These are the rules that determine when a mobile node should perform handover. A decision for vertical handover may depend on various parameters such as bandwidth, delay and transmitted power.

Handover Implementation or Handover Execution: This is the process by which a mobile node's connection is rerouted from their existing network to a new network in a seamless manner. Mostly, it requires the network to transfer routing information about the mobile terminal.

Recently, research has developed various techniques and approaches to deal with the multicast mobility process and handover at different layers of the network. Most of these strategies are based on a modification and implementation at the network layer for example

modifying the Mobile IP protocol [28,19,42]. Other approaches operate at the transport layer and at the application layer such as a modification to SIP (Session Initial Protocol) [43].

For the Network Discovery Phase, the MN must search for available reachable wireless networks. In this state most MNs should always keep all network interfaces on. However, keeping network interfaces on all of the time becomes a weakness because it consumes battery power without any benefit of delivering real data.

In [8], the context information about a network is stored in the context awareness database. Many parameters are collected by system discovery and are dependent on the adopted network interface card within the MN. For instance, these could include signal RSSI (Received Signal Strength Indicator), bit rates, MAC (Media Access Control), etc.

In addition, network monitoring systems adopt SNMP (Simple Network Management Protocol) to extract more relevant network information from all APs (Access Points) and ARs (Access Routers) in each subnet such as data transmission/receiving rate, network loading and multicast connections.

2.5.2.2 Handover Decision Phase

There are many research papers proposing strategies for making decisions about handover. We categorize these into three types: network-controlled handover, mobile-controlled handover and mobile-assisted handover.

Network-controlled handover is when the network makes the handover decision for a mobile node. Shantidev Mohanty [45] proposed a novel architecture using the Network Interoperating Agent called NIA and Integration Gateway (IG) to integrate the 3G systems and WiFi networks of various providers. The IG functions as a traffic monitoring unit and seamless roaming module.

In [42], mechanisms are presented for PMIPv6 (Proxy Mobile IPv6) by two multicast mobility listeners called LMA-MLM and MAG-MLM. Both listeners will be responsible for subscribing to the multicast group and receiving multicast packets on behalf of MNs in its

domain. The MAG-MLM uses a MAG (Mobile Access Gateway) to detect the detachment of the MN.

The LMA-MLM uses a LMA (Local Mobility Anchor), which is responsible for maintaining the reachability of MNs by updating the binding cache and maintaining the tunnel to the MAG for packet delivery. Their process can achieve the whole multicast handover process without the involvement of the MN. However, the LMA-MLM still has to deal with the problem of encapsulation.

Kim and Han [46] have proposed PMIP protocol with IP multimedia system. The protocol that they proposed naming PMIP-M protocol. The main idea of this protocol is the user can continue received multicast data stream even when they migrates in the new network which without IP multicast capability.

Mobile-controlled handover is where the mobile node must take its own signal strength measurements and make the handover decision on its own. In [47], an algorithm is proposed based on the Markov decision process (MDP) formulation, which tries to maximise the expected total reward of a connection.

Mobile-assisted handover is where the decision to handover is made by the mobile node and network in cooperation. In [43] a mobile QoS (Quality of Service) framework is proposed for heterogeneous IMS (Internetworking Management System) interworking by modifying SIP multicast. However, this method consumes network bandwidth and MNs need to reserve bandwidth. Park and Won [48] analyse about mobility management architecture such as MIP and PMIP protocol for avoiding any tunnels for multicast delivery in heterogeneous network.

C. Wen et al. [8] proposed an integrated framework for MNs and core wireless networks by using a context-aware handover scheme. They periodically collect the parameters of various available network status reports and information about host application services. For example, the best one from a list of available access points is selected based on network conditions and user defined policies. The multicast connection management is then performed efficiently by a multicast agent in WiFi and 3G networks. Hence, the MNs can make handover decisions to activate the proper network interface switching to avoid discontinuities in the delivery of multimedia application services.

One thing that is important when the handover occurs is considering the handover metrics used to make the decision. There are many handover metrics that are used to indicate when handover should be performed:

Connection cost

For users, connection cost is a key consideration, especially when different network operators may create different billing schemes. Hence, it might affect the user's choice of handover.

Network-related parameters

There are many network parameters used for making the decision such as bandwidth, load, network latency, traffic congestion, location information and so on. Moreover, that information is useful for load balancing across different networks and QoS [20].

Application types

For example, some multimedia applications require reliability in networks. Hence, different types of applications may require different levels of QoS determined by the percentage of lost packets or the delivered data rate.

Battery power

Battery power may be a significant factor for handover in some cases. Moreover, an MN with multiple interfaces must keep an interface active all the time but this consumes battery power even without receiving any data.

2.5.2.3 Handover Implementation

Several multicasting schemes have been proposed for mobile networks. In [44] they used the ISHO (inter-system handover) protocol and include the concept of a dynamic boundary area to support seamless roaming between different networks. However, this scheme requires the NIA and IG to be added into the network architecture.

In addition, when an MN moves across to the new network, it is important to consider the associated security mechanisms. The security mechanisms of 3G-WiFi do not address the security of multicasting. However, for supporting a secure link for multicasting a framework of multicast key agreement by a modified EAP-AKA protocol has been presented [49]. Here

EAP-AKA is the Extensible Authentication Protocol method for UMTS Authentication and Key Agreement used for authentication in wireless networks. They are divided into two phases: Initial phase and Key refresh phase. The benefit of this technique is that it saves communication overhead, computation overhead and does not need a huge change for existing protocols [50].

2.5.3 Handover within WiFi Networks

For the WiFi-3G handoff process, [51] a method is proposed to reduce latency by using the ISHO (inter-system handover) protocol which includes the concept of a dynamic boundary area to support seamless roaming between different networks. However, most of this research tries to solve the problem in the network layer with others seeking to solve it by modifying the transport layer. For instance, in [52] a mobile QoS framework is proposed for heterogeneous IMS interworking by modifying SIP multicast. However, this method consumes network bandwidth and MNs need to reserve bandwidth.

The handover procedure of Mobile IPv6 protocol can be expressed as 2 parts: L2 (link layer) handover latency and L3 (network layer) handover latency. The L2 handover consists of channel scanning process, authentication and association process. Generally, L2 handover latency is about 100 - 300 ms however it depends on the structure of network topology. For L3 handover latency in MIPv6 consists of two main parts: CoA configuration and Binding update. The process of CoA configuration is starting from Router discovery process until the MN obtained a new CoA. The Binding update procedure is about the MN inform HA and CA nodes about their new location which is new CoA address. Normally, handover latency of Mobile IPv6 is about 2000-3000 ms [53] this is why it is possible that MN can lose connection completely during handover process.

Tien-Thinh Nguyen [54] had applied DMM (Distributed Mobility Management) concept with multicast mobility in IPv6 network by enable IP Multicast with MLD proxy function. However, the result show that when the mobile receiver moves, the network have to build the tunnel between source and destination. This is a case of tunnel convergence problem. Also it has a problem about service disruption and delay which cannot acceptable in some delay-sensitive service. Moreover Nguyen and Bonnet [55] had been studying about load balancing

mechanism among LMAs (Local Mobility Anchor) to solve a bottleneck and single point of failure issues.

2.5.4 Multicast Handover in Wireless Networks

For streaming multimedia content in 3G network has been standardized under the 3GPP-PSS (3rd Generation Partnership Project – Packet Switched Streaming Standard) which is released in April 2001. The 3GPP-PSS are described presentation of information, the audio and video formats of that stream within complete protocol stack in IP layer [56]. In UMTS, the IMS was extended to include MBMS. The 3GPP MBMS has the following characteristics:

- There is no immediate Layer 2 source-to-destination transition, resulting in transit of all multicast traffic at the GGSN.
- As GGSNs commonly are regional, triangular routing, distant entities and encapsulation this may cause a significant degradation of efficiency.

In 3GPP2 (3rd Generation Partnership Project 2) [57], the MBMS has been extended to the Broadcast and Multicast Service (BCMCS) [58], which on the routing layer operates very similar to MBMS. In both 3GPP (3rd Generation Partnership Project) and 3GPP2, multicast can be sent using either point-to-point (PTP) or point-to-multipoint (PTM) tunnels, and there is support for switching between PTP and PTM.

A mobile multicast node may change its point of Layer 2 attachment within homogeneous access technologies (horizontal handover) or between heterogeneous links (vertical handover) [59]. In [60] has modified PIM-SM to support handover latency and keeping connection. By proposed multicast routing protocol named MC-PIM-SM by extended from PIM-SM protocol [61]. Mobility applications transport for MIH are required as an abstraction for Layer 2 multicast service transfer in an Internet context [45] and are specified in [62].

Functions required for MIH include:

Service context transfer.

- Service discovery.
- > Service invocation.
- > Service context transformation.

In [63] is shown the amount of multicast packet loss, when handover occur at the mobile node in equation. Suppose $\varphi_P^{(BASE)}$ is the amount of multicast packet loss for the base multicast handover procedure. Let δ_s denote the average multicast session arrival rate per second at the mobile node. $\varphi_P^{(BASE)}$ is obtained as

$$\varphi_P^{(BASE)} = \delta_S E(S) L_{HO}^{(BASE)} \tag{2.6}$$

Where E(S) is the average session length in packets

This research will focus on mobile receiver problems and so methods to solve these problems have been proposed. The problem of achieving seamless mobile receiver multicast handover can be addressed by one of the following:

> Home subscription-based solution:

- o *Mobile IP Home Subscription* or *bi-directional tunnelling:* this approach relies on the Mobile IP protocol and uses a local router in the home network as the multicast router for responses such as forwarding multicast group membership control messages to the mobile node even when it moves to a foreign network. However, tunnelling will create the process of encapsulation/ decapsulation and fragmentation problems.
- O Multicast encapsulation: that is encapsulation of multicast data packets to shield mobility and to enable access to remotely located data services such as from the home agent.
- Remote subscription-based solution: by forcing the mobile node to re-initiate multicast distribution following handover. However, this technique cannot support session persistence under multicast source mobility.

- o *Agent assistance*: there are many protocols that are proposed for agent-assisted handover for host-based mobility such as Fast MIPv6 (FMIP6) and Hierarchical MIPv6 (HMIPv6).
- O Network-based mobility management: Proxy MIPv6 (PMIPv6) [19] is multicast transparent in the sense that the MN experiences a point-to-point home link fixed at its LMA (Local Mobility Anchor). In [63] network based mobility management is deploying for the mobile nodes, also the tunnel between the LMA and itself for the MN. However, PMIPv6 still has a problem about MTU size from spanning tunnels at the receiver site.
- ➤ Hybrid architectures: that tries to find the methods, which avoid the complexity at the internet core network.
 - o *Hybrid shared tree:* [64] proposes the hybrid shared tree approach by introducing a mobility-agnostic multicast backbone on overlay.
 - O Hierarchical local registration: the network model has proposed hierarchical and local registration. The registration consists of having a root FA (Foreign Agent) and lower FAs. The MN registers its CoA with the root FA. All the FAs exchange summary reports that consist of the common multicast group of interest on the lower levels. However, this approach required an extra cost to select multicast service provider (MSP) [65].
- ➤ *MLD Extensions:* there are many methods by extended MLD message. Some of them modify an MN operating predictive handover such as FMIPv6.

2.6 Summary

The approach which is presented in this research aims to offer a smooth handover between the home network and the foreign network. The advantages of both home subscription-based and remote subscription-based solutions have been combined. A modification has been offered which responds to problems posed by mobility. However, this framework does not make use of bi-directional tunnelling, which means that the framework solves three problems: tunnel convergence, encapsulation / decapsulation overhead delay, and fragmentation problems.

The concept of a remote subscription-based solution has been applied to this framework by creating a reserve route to neighbouring networks but kept in standby. The FMIP6 protocol also uses the remote subscription but it still suffers a short lost-connection time of MN. Moreover, this approach does not require multicast tree reconstruction as many previous methods do. The design detail of this framework will be described in the next section.

Chapter 3 A New Framework for Multicast Mobility in WiFi Networks

3.1 Introduction

This chapter presents the designed framework that was produced during this research. The chapter describes about network architecture, protocol overview and the details of protocol process step by step.

3.2 Network Architecture

In this research, there is a requirement for a system that provides support for research scenarios in WiFi networks that are connected through the internet. According to the real world, the WiFi network will combine with an IP wired infrastructure network with a gateway router to route data through to the internet. There are a variety of clients in the network such as PCs, laptops and mobile devices [66].

Before starting to explain the detail in each process of the framework, it is first necessary to show the scope of the network architecture. Hence, in this chapter the network architecture being used in this research is described and shown in Figure 3-1.

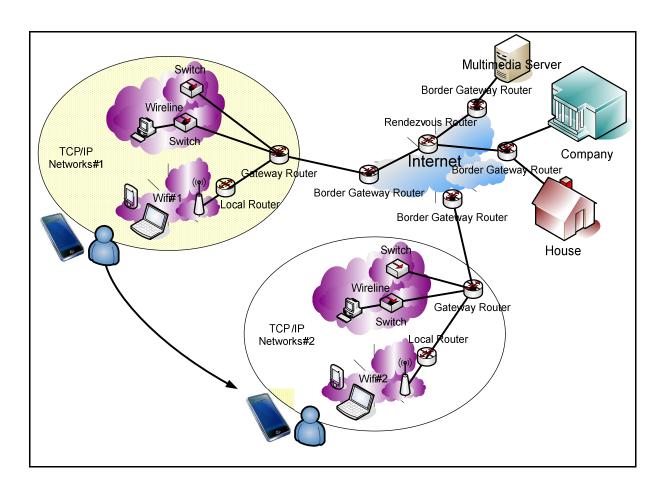


Figure 3-1 Network Architecture

Figure 3-1 is shown the stage which we consider WiFi to WiFi handover.

- Here, the architecture comprises only WiFi networks that comprise: a media server which represents the source server, which is sending the multicast packets to mobile devices in the network. The rendezvous router is a central router in a multicast tree.
- Border gateway router is a core router which is enabled for multicast delivery services and supports multicast routing protocols such as PIM (Protocol-Independent Multicast), MOSPF (Multicast Extensions to Open Shortest Path First), DVMRP (Distance Vector Multicast Routing Protocol) and so on.
- Gateway router is a core router within a company network which is connected to both WiFi and wireline networks, providing Internet access and supporting multicast routing protocols.
- Local router is a router that connects between a WiFi access point and other network devices in the local network.
- WiFi access point provides localized wireless coverage.

Mobile nodes with embedded WiFi interface roam within this network.

3.3 Protocol Overview

In this section we will describe the process by which the new proposed framework is designed to improve handover performance for multicast services. Procedures and protocols already exist for handling handover from one WiFi network to another. However, such protocols, of which mobile IP is a key example, achieve handover by firstly making contact with a new WiFi base station, obtaining a new IP address and then re-routing traffic to that new address through a modified multicast tree. Unfortunately this leads to a loss of connectivity and hence, service whilst this process is taking place. Similarly, IGMP manages the distribution of multicast services through the establishment of a multicast tree which is maintained by the routers. When a mobile node, moves, this tree needs to be modified to accommodate the mobile node's new location. This therefore leads to further delay which handover takes place. Our approach is to complete as much of the existing handover process as possible before the physical handover actually takes place. This therefore will minimise the actual handover delay at the expense of having to establish several connections to neighbouring networks, most of which may never be activated.

The overall concept of our new framework is that a mobile node will use mobile IP but modified in such a way as to allow connectivity to multiple foreign networks. In so doing a mobile node will receive a Care of Address from each foreign network that is within range. These addresses are then used to compute multiple extensions to the existing multicast tree, which are then held in a standby mode until required. The standby route represented as a dash line in Figure 3-2 below. Once a mobile node actually completes handover by moving to a new WiFi base station, the Care of Address that has already been obtained and the associated change to the multicast tree are then activated with traffic being routed to that mobile node via this new route.

Hence, establishing the Care of Address and determining the required modifications to the multicast tree in advance of needing them reduces the handover delay to one of switching between the route currently being used within the multicast tree and the new one. The goal is to achieve this through the minimal modification to existing protocols and procedures. In this

chapter it is shown how the new framework can be applied to a network that employs mobile IP and IGMP.

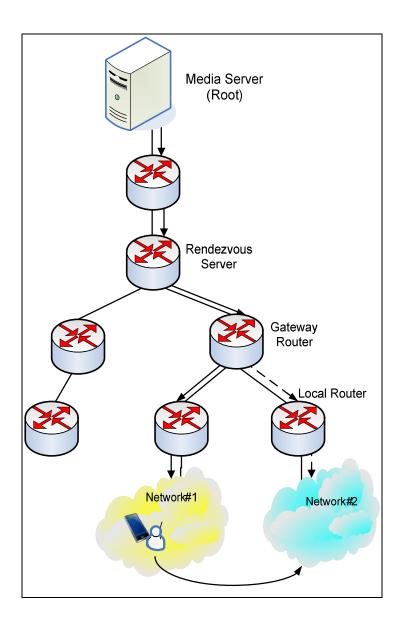


Figure 3-2 Multicast Route

3.4 Process Diagram

In this section we will describe a process diagram of the proposed framework in details. From the Figure 3-3 until 3-6, the modified messages or processes in the proposed framework will be represented by underlining the text, while standard protocol process will be represented by normal text.

The outline of our process is as follows [67]:

- The mobile device connects to the WiFi#1 as home network and receives a multicast data stream as usual.
- The mobile device sends the message to join WiFi#2 as a foreign network in advance by modified Mobile IP protocol.
- The mobile node uses its CoA address from WiFi#2 to establish a new multicast route with the same media server.
- After it receives the multicast message from the second route, it disconnects the second route at a point along the path between its local router for WiFi#1 and the rendezvous router.
- Local router of WiFi#2 network sends a modified PIM protocol message to keep its connection as a standby route.
- ➤ If handover occurs then the modified ICMP protocol message will be sent to reconnect the second route. This will minimize handover delay because the second route has already been configured and just needs to switch from standby to active.

Figure 3-3 is shown the initial steps of how a mobile node is able to register for receipt of a multicast stream being delivered by the media server.

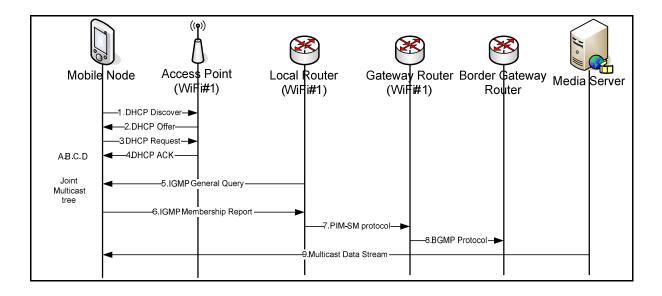


Figure 3-3 Starting connection process

Details of this process are as follows:

- 1. The mobile node sends a DHCP Discover to its local access point by broadcasting in order to discover the DHCP server that can supply it with an IP address.
- 2. In this case WiFi#1 will respond with a DHCP Offer including an IP address.
- 3. If the mobile node accepts that IP address, the mobile node will send a DHCP Request message including that IP address back to the access point for confirmation.
- 4. Access point WiFi#1 will send a DHCP ACK message back to the mobile and allow the mobile node access to the network. At this stage, the mobile node will has IP address for establish a connection to media server.
- 5. Usually, if a local router is enabled for multicast delivery it will send an IGMP General Query every 60 seconds within the network. Hence, at this stage WiFi#1 will send an IGMP General Query message out to the mobile node [68].
- 6. In this scenario, the mobile node wants to connect to the multicast tree and so it will send an IGMP Membership Report, including the IP multicast address to which it wants to connect
- 7. The local router will then use its multicast routing protocol to connect to the appropriate multicast tree. In this case it will create a PIM-SM protocol message and send it to the gateway router within the WiFi#1 network.
- 8. Since in this case the media server is outside the local network and on the internet, BGMP will be used by the gateway router to connect to the multicast tree.
- 9. Once connected to the multicast tree, the mobile node will start to receive the multicast data stream.

When a mobile node moves then in effect its position on the multicast tree moves or, in some cases, the multicast tree will need to be extended to accommodate the mobile nodes' new location. Moving within a network also requires the mobile node to be issued with a care of address. Therefore, in order to improve handover efficiency, in our scheme we seek to obtain the care of address and modify the multicast tree ahead of the time when it is actually needed.

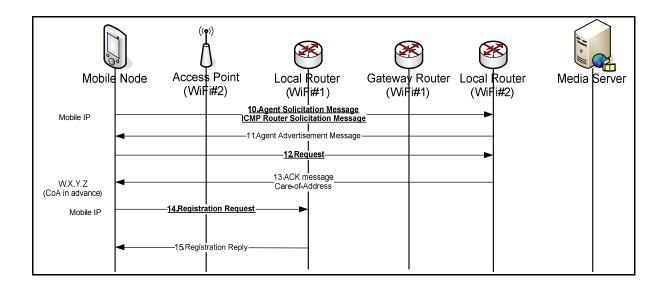


Figure 3-4 Request CoA address in advance process

Figure 3-4 is shown the process of obtaining a care of address in advance from access point WiFi#2 as the mobile node is preparing to move away from access point WiFi#1.

The procedure is described as follows:

- 1. When the mobile node received a signal from foreign agent which is Local Router on WiFi#2 network in Figure 3-4, that implies the mobile node is in range of foreign agent. In the new framework, we designed that the mobile node will broadcast a Mobile IP Agent Solicitation Message every 30 seconds for registering to the new foreign agent. This process Mobile IP does not use a new packet type for agent solicitation, it uses the router solicitation packet of ICMP.
- 2. After receiving Mobile IP Agent Solicitation Message, the local router (foreign agent) in the Wifi#2 network will send back an Agent Advertisement Message which includes the IP care of address (CoA) to the mobile node.
- 3. If the mobile node accepts this address, the mobile node will send a Request message to confirm to the local router in WiFi#2 network that it wants to use this IP address.
- 4. The local router will subsequently acknowledge this with an ACK message. At this stage, that means the mobile node had CoA address in advance before moving to WiFi#2 network. Also it means the mobile node can use this IP address if it handover into WiFi#2 network.
- 5. Normally, when a mobile node receives a CoA address it needs to register this with

its Home agent which, in this case, is the local router in WiFi#1 network.

6. After the home agent stores the information in their database, it will send a Registration Reply to confirm to the mobile node. At this stage, the mobile node has two addresses; one is an IP address from its home agent and the other is the CoA address from the foreign agent.

The next strategy is to establish a new multicast route from the foreign network to the multicast tree, the details of which are shown in Figure 3-5.

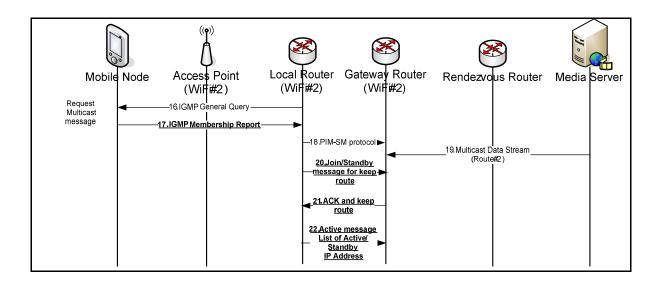


Figure 3-5 Request Multicast packets and keep route

The procedure is described as follows:

- 1. Normally, the local router will broadcast an IGMP General Query to the client in their network which in this case is the local router in the WiFi#2 network. For querying that there is any client would like to join any multicast tree in network.
- 2. If the mobile node wants to join the multicast tree it will send an IGMP Membership Report including the IP multicast address to the local router. In this framework the mobile node will use the CoA address to communicate.
- 3. The local router will use the PIM-SM protocol to communicate with the gateway router in the WiFi#2 network for creating the route to multicast tree by connecting to rendezvous router.
- 4. After joining the multicast tree, there is a multicast data stream from media server to

local router. However, at this stage, there are two routes connected to the same multicast tree but only one should be live. Therefore, we are proposing to modify the PIM-SM message to keep the route between the local router in the WiFi#2 network and the gateway router in a standby mode.

Figure 3-6 is shown details of the process when the handoff process occurs. The idea is to change the route at the foreign agent from standby to active. Here the secondary route which has been established using the CoA needs to be switched from standby to active and the currently active route which uses the home address needs to be switched from active to standby.

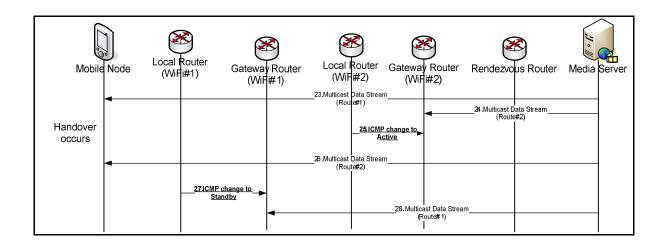


Figure 3-6 Handover process

The procedure for achieving this is as follows:

- 1. The local router in the WiFi#2 network will send a modified ICMP message to the gateway router to change mode from standby to action. This means that the new route for the multicast data stream is ready to connect. Moreover, this method should reduce handoff latency time.
- 2. After that the mobile node will start to receive multicast data stream via the new route.
- 3. The next step is to change the old route to standby mode by sending an ICMP change to Standby message to the gateway router in WiFi#1 network.
- 4. At this stage we still keep the old route from media server to the local router in WiFi#1 network in case the mobile node moves back to the old network. However, it will

have setting timeout for delete route.

3.5 Modified Protocol Message

In this research, there are some protocol messages that have been modified to support our designing framework, which are:

3.5.1 PIM Protocol Message

The message that we have been modified in PIM protocol is "Join/Prune message". The format of Join/Prune message is shown in Figure 3-7. This message modified for keeping status join/standby message information which sends between local router in WiFi#2 and rendezvous router in the process number 20 in the Figure 3-5.

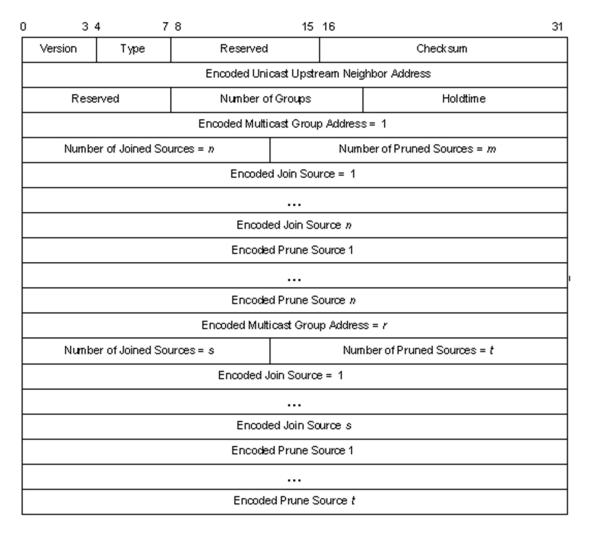


Figure 3-7 Join/Prune messages format [20]

3.5.2 ICMP Message

The ICMP message has been modified and adapted for controlling and changing the status of a second route when the handover happened. In our framework when handover occurs, local router in WiFi#2 will send the message to change the status to become active as in process number 25 in Figure 3-6. After that, the mobile node can continuously receive the multicast message from multicast tree without rebuilding the tree. We modified ICMP to support this strategy. The 3-8 is shown the standard ICMP message.

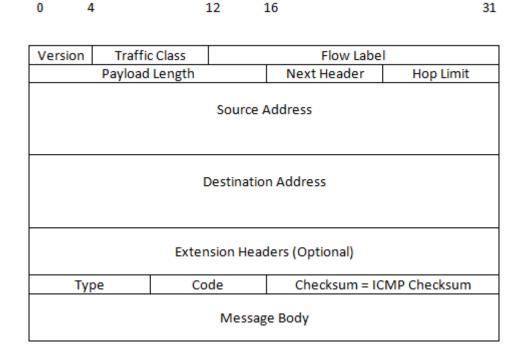


Figure 3-8 ICMP message format [35]

3.5.3 Mobile IP Message

The Mobile IP message is modified for sending to the local router in WiFi#2 network to ask for CoA address in advance. In process number 10 in Figure 3-4, we will modify an Agent solicitation Message of Mobile IP to ask for CoA address from local router in WiFi#2 network by changing the flag H to active. The format of Mobile IP message is shown in Figure 3-9.

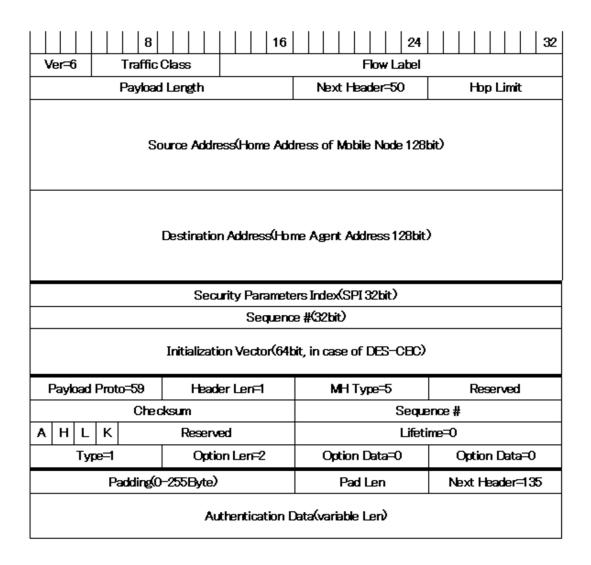


Figure 3-9 Mobile IP message format [32]

3.5.4 IGMP Message

In the designing framework, IGMP protocol has been modified to do a process of joining multicast tree for the second route before handover take place. This IGMP message will carry an "IGMP Membership Report" message as show in the process number 17 in Figure 3-5. The standard format of IGMP message is shown in Figure 3-10.

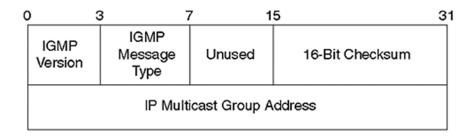


Figure 3-10 IGMP message format [6]

3.6 Summary

This chapter describes the network requirement, network architecture, the protocol process and modifying messages in detail. According to our goal we are trying to minimize the handover process in multicast mobile. The framework creates the reserve routes via neighbour zones, which are connected to the same multicast tree in advance. When handover occurs, the system only changes the status from standby to active and from active to standby mode.

We predict this strategy can reduce handover latency time because the network already has reserve paths, which are connected to neighbour zones. The delay will become only the time for changing the status of the reserve route. Hence, the handover latency time will become the time from local router in WiFi#2 to be detected by the mobile node including sending ICMP message to change the status until mobile node receives the multicast message. If it is compared with the previous methods that have been proposed, this strategy avoids many problems such as reconstruction of the multicast tree, network inactivity because it knows the new route in advance, multicast encapsulation/decapsulation because this method does not use tunnelling and so on.

Chapter 4 A Framework Simulation in OPNET Modeler

4.1 Introduction

In this chapter will present about a designed framework which simulating on OPNET Modeler software. The research is simulating those techniques and designed processes on OPNET Modeler software which is a licensing at University of Salford. OPNET Modeler software is a network simulation software and solution. This software provides for application and network management issues.

4.2 Network simulation

For doing research in the network field, network simulation software is a very useful and important tool. As researchers or protocol designers have to design and testing the system in simulation software before using it in a real network. There are many network simulations that widely used in networking research such as OMNET++ (Objective Modular Network Testbed in C++), NS-2 (Network Simulator version 2) and QualNet [69].

OPNET Modeler is generally used by researchers, developing protocol designers and so on. The OPNET software was funded in 1986 by Alain Cohen. OPNET stands for Optimizing Network Engineering Tools [70]. OPNET Modeler provides a comprehensive development environment which is powerful for instance simulation, data analysis, model design and etc. also it can support lot of technologies including local area network (LAN), mobile network, sensor network, wireless network and so on.

4.2.1 Basic Structure within OPNET Modeler

This is the workflow for OPNET Modeler. Normally, the researcher use these steps to build a network model, create the traffic, choose statistics and then run simulations.

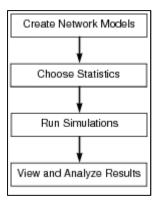


Figure 4-1 Basic step for creating network simulation

These 4 steps in Figure 4-1 consist of creating network environments which is including network devices and traffic, and then choose statistics that we want to study. Next step is run simulations. Finally, view and analyze the results. To complete these 4 steps, OPNET Modeler provides variety kinds of editor to support users as show below.

> The Project Editor

This is a main area of OPNET simulation. We use this area to create network topology, generate traffic within network and view the results via this editor. Moreover, this area still covers about choosing statistics and running simulations.

> The Node Editor

The user can define the behavior of each network object via "Node Editor". In Node Editor of each model, the behavior is defined using different modules for example data storage, data creation, etc. A network object in OPNET Modeler is typically building up from multiple modules which define that object. The user can add their modules into the network object via Node Editor.

➤ A Network Model in the Project Editor

The OPNET Modeler let user to design and create any elements of network as they wish. For instance, user can create node, link model, process models and build packet formats. Also, the user can create filters and parameters that they want to analyze.

➤ Node Model

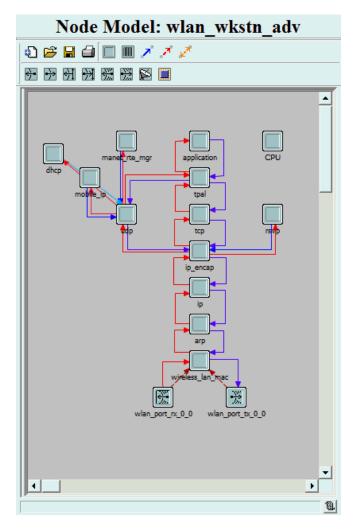


Figure 4-2 Node Mode example

> The Process Model Editor

The OPNET Modeler let user design and creates their process models via the "*Process Editor*". The user can start from create node model in Node Editor and then they can build process model, which control the functionality of that node mode.

Process Model

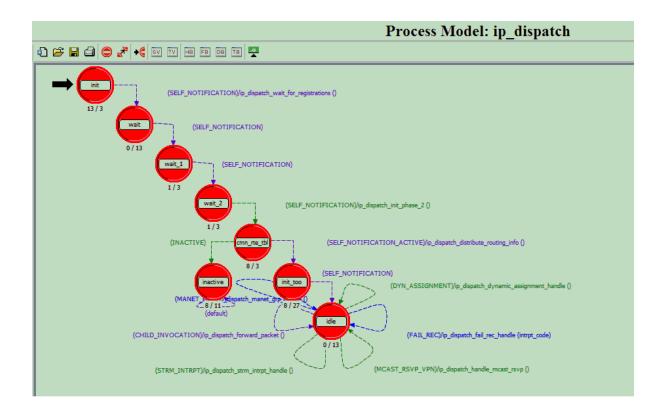


Figure 4-3 Process Mode example

4.3 Implementation of the Proposed Framework in OPNET Modeler

Due to the implementation of this research has been simulation environments and testing the performance of designing on OPNET Modeler software version 16.0 which is not supported multicast communication over IPv6 WiFi environment. Hence, the implementation and development of this thesis has been modifying based on IPv4 environment. However, the concept and designed of this framework can adapt to WiFi network both on IPv4 and IPv6 Networks.

4.3.1 Network Architecture

Normally, the first thing that has to start network simulation is OPNET Modeler is to create network architecture. The common start network topology that is used in this research is shown in Figure 4-4.

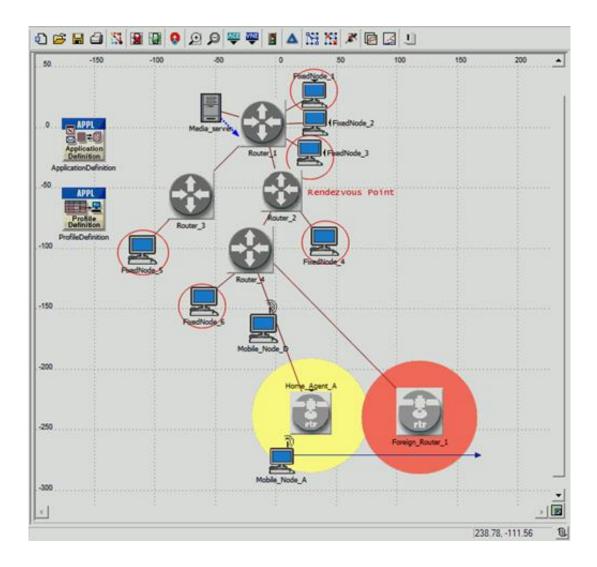


Figure 4-4 the common network topology that is used in this research.

4.3.2 Process Model

Therefore, we have to create a network environment and add some processes into standard protocol. So, we have to deal with process model many times.

4.3.2.1 Asking CoA in Advance process

In the designed framework, a mobile node will send an Agent Solicitation Message to foreign agent to ask for a CoA in advance. This process happens via the Mobile IP protocol. However in OPNET Modeler, this method will happen by creating an extra state and adding into the process model of mobile_ip_mn.

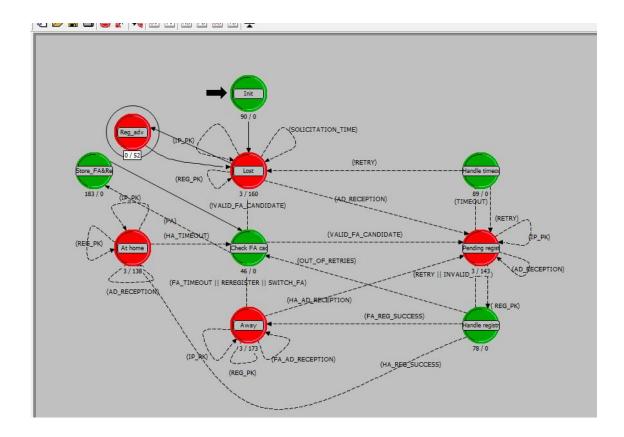


Figure 4-5 Process Model: mobile_ip_mn

Figure 4-5 is shown the Process Model of Mobile IP protocol within mobile node. To achieve the process of asking CoA in advance, state name "Reg_adv" have to be added. The coding of this state is shown in the Figure 4-6.

```
static void
mip_mn_agent_solicit_pk_send_adv (void)
               {
Packet*
double
                                      solicit_pkptr_adv;
solicit_interval;
6 7 8 9 10 111 12 13 14 15 16 17 18 19 20 21 22 22 24 25 26 27 28 29 30 40 41 42 43 44 54 64 7 48 9 50 ...
               /** PURPOSE: Send the ICMP agent solicitation packet.**/
/** REQUIRES: none. **/
/** EFFECTS: Packet will be given to IP to handle.**/
FIN (mip_mn_agent_solicit_pk_send (void));
               /* Time to send out the solicitation. */
               usleep(10000000); // 10 secs
solicit_pkptr_adv = op_pk_create_fmt ("mobile_ip_irdp_solicit");
                  * Send the packet out. */
odule_data->ip_ptc_mem.child_pkptr = mip_sup_irdp_pkt_encapsulate
   (solicit_pkptr_adv, home_address, subnet_bcast_addr, Icmpc_Type_IRDP_Sol);
                    Record some stats. */
                op_stat_write (irdp_sent_pkts_sh, 1.0);
op_stat_write (irdp_sent_bits_sh, op_pk_total_size_get (module_data->ip_ptc_mem.child_pkptr));
op_stat_write (g_irdp_sent_bits_sh, op_pk_total_size_get (module_data->ip_ptc_mem.child_pkptr));
op_stat_write (g_irdp_sent_bits_sh, op_ok_total_size_get (module_data->ip_ptc_mem.child_pkptr));
op_stat_write (g_irdp_sent_bits_sh, 0.0);
               /* Invoke IP to handle the packet. */
op_pro_invoke (proc_info_struct_ptr->ip_phndl, OPC_NIL);
                /* Schedule the next transmission. */
if (++solicit_count > 3)
                       solicit_interval = Mipc_MN_Solicit_Min_Interval * pow (2.0, (double) (solicit_count - 3));
                       if (solicit_interval > MipC_MN_Solicit_Max_Interval)
                               solicit_interval = MipC_MN_Solicit_Max_Interval;
                       solicit_interval = MipC_MN_Solicit_Min_Interval;
                solicit_timer_ehnd1 = op_intrpt_schedule_self (op_sim_time () + solicit_interval, MipC_MN_Timer_Solicit);
```

Figure 4-6 The coding of "Reg_adv" state

4.3.2.2 Joining Multicast using CoA Address

This process model has been called after the mobile node received CoA in advance and then tries to build another route to multicast tree by using CoA address. This stage has been extended from IGMP process model. To achieve this, we have to create a process state adding into IGMP process model name "JOIN_ADV" as is shown in Figure 4-7. The coding of process state is presented in Figure 4-8.

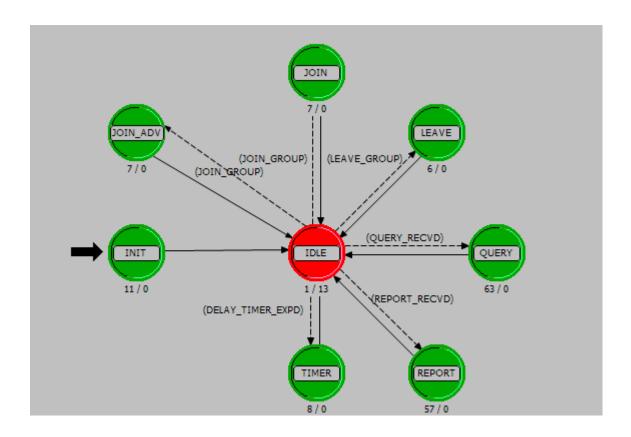


Figure 4-7 Process Model: IGMP host

Figure 4-8 The coding of "JOIN ADV" state

4.3.2.3 Re-join Multicast

When the mobile node has been realized that, now it is in a foreign agent. The process re-join will call multicast joining state in process model "ip_igmp_rte_grp" to send a message to join multicast group again.

```
/** time. Notify the router and start the group membership timer **/

/* Generate trace messages */

if (LTRACE_IGMP)

/* ip_address_print (ip_addr_str, ip_grp_addr);
sprintf (msg0, "IP Group Address : %s", ip_addr_str);
sprintf (msg0, "IP Group Address : %d", ip_interface);
op_prg_odb_print_major ("Received an IGMP Membership Report message for: ", msg0, msg1, OP op_prg_odb_print_major ("Received an IGMP Membership Report message for: ", msg0, msg1, OP op_prg_odb_print_major ("A local host has joined the above group. Notifying PIM-SM router

/* Notify the router */
ip_imp_rte_grp_notify_router (Ipc_Igmp_Rte_Pim_Sm_Notify_Plus);

/* Start the group membership timer only if simulation efficiency is disabled */
if (igmp_attrs_ptr->igmp_sim_efficiency == OPC_FALSE)

/* Generate trace messages */
if (LTRACE_IGMP)

sprintf (msg2, "Starting the Group Membership timer with value, %d for the above group op_prg_odb_print_major (msg2, OPC_NIL);

ip_imp_rte_grp_start_timer (IPC_IGMP_RTE_GRP_TIMER, igmp_attrs_ptr->grp_member_interval);

/* Tag this connection process for debugging purposes. */
if (op_sim_debug () == OPC_TRUE)

if (op_sim_debug () == OPC_TRUE)

if (op_sim_debug () == OPC_TRUE)

ip_address_print (ip_addr_str, ip_grp_addr);

sprintf (msg0, "Group address: %s\tinterface index: %d", ip_addr_str, ip_interface);

op_pro_tag_set (op_pro_self (), msg0);

/* Op_pro_tag_set (op_pro_self (), msg0);
```

Figure 4-9 the coding for re-join multicast

4.3.2.4 Keeping Multicast Route

After the mobile node created other multicast routes, the mobile node have to keep other routes become Standby mode, only one route at a time being Active mode. To achieve this goal, we modified Join/Prune message in PIM protocol to keep these multicast routes alive. The coding of this process state is shown in Figure 4-11.

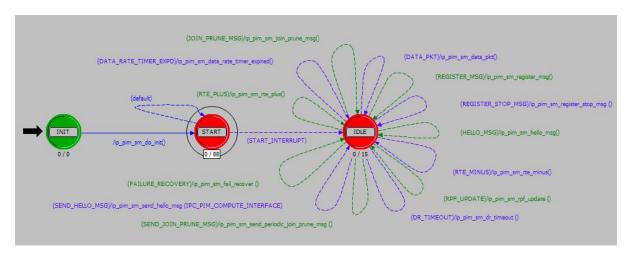


Figure 4-10 Process Model: PIM-SM protocol

Figure 4-11 The coding for keeping multicast route

4.3.2.5 Store CoA addresses

When the mobile node moves, the mobile node starts to receive CoA address along the path. However, in the large network which consists of many routers in different zones, the mobile node also receive multiple CoA address from foreign router. Hence, the mobile node needs to have a process of store and process multiple CoA addresses. To solve this issue, we had modified Mobile IP protocol to have the process of store and retrieve CoA address on the mobile node. We have created the state named "Store_FA&Re" state which is shown in Figure 4-5 to solve this issue.

```
De X PD ô
            /* Access the packet. */
reg_pkptr = op_pk_get (input_strm);
                                                                                                                                                                                                              •
            /* Create an ICI for communication to Mobile IP process. */
reg_ici_ptr = op_ici_create ("mobile_ip_reg_ici");
            /* First the packet format. */
op_pk_format (reg_pkptr, pk_format_name);
  10
11
12
13
14
15
16
17
18
19
20
21
22
            if (!strcmp (pk_format_name, "mobile_ip_reg_req"))
                    {
/* Request! */
reg_type = MipC_Reg_Type_Req;
                    /* Retrive information from the packet. */
op_pk_nfd_access (reg_pkptr, "Care-of Address", &care_of_address);
op_pk_nfd_access (reg_pkptr, "Lifetime", &lifetime_req);
op_pk_nfd_access (reg_pkptr, "S", &simultaneous_binding);
                    /* Set the values on the ici to mobile ip. */
op_ici_attr_set (reg_ici_ptr, "care_of_address", care_of_address);
op_ici_attr_set (reg_ici_ptr, "lifetime_req", lifetime_req);
op_ici_attr_set (reg_ici_ptr, "s", simultaneous_binding);
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             else
                     {
if (!strcmp (pk_format_name, "mobile_ip_reg_reply")) /* strcmp function is for compare
                              /* Reply! */
                             reg_type = MipC_Reg_Type_Reply;
                             /* Retrive information from the packet. */
op_pk_nfd_access (reg_pkptr, "Lifetime", &lifetime_grant);
op_pk_nfd_access (reg_pkptr, "Extension", &care_of_address); /* Use extension for F
op_pk_nfd_access (reg_pkptr, "Code", &reply_code);
                               /* Set the values on the ici to mobile ip.
                             /* Set the values on the ici to mobile ip. */
op_ici_attr_set (reg_ici_ptr, "lifetime_grant", lifetime_grant);
op_ici_attr_set (reg_ici_ptr, "extension", care_of_address);
op_ici_attr_set (reg_ici_ptr, "reply_code", reply_code);
                     else
  44
                              op_sim_end ("Unknown packet format received.","","","");
  46
                    }
```

Figure 4-12 the coding of "store FA&Re" state

State "store_FA&Re" has two processes of designed framework in there. That is keeping CoA from foreign router, and another process is retrieving CoA address. Some part of coding state is shown in Figure 4-12.

Figure 4-13 the coding of structure of CoA address list

Figure 4-13 presents the coding of the process in "Store_FA&Re" state. This part is shown the structure of each CoA information that has been stored in the mobile node responsible by Mobile IP message.

Figure 4-14 the coding of Binding Update multiple CoA address to home agent

Normally, the mobile node has to report every CoA address that received from foreign agent to home agent. Hence, in the framework the mobile node has to do the same but in advance. So, this process is part of reducing handover latency because the binding update process had been doing in advance before handover occur. The coding of this process is shown in Figure 4-14.

4.4 Summary

This chapter presents the fundamentals of OPNET Modeler software, the implementation that has been made to achieve our designed framework, the state diagram, some coding of the process. The full coding of this research has been attached in Appendix B.

However, some process of the designed framework, we do not need coding program. We can handle it by setting a value of attribute of the protocol such as IP parameter and wireless parameter. Also, the designed process setting a second multicast route to Standby/Active mode can be done by changing the parameter flag at gateway router.

Chapter 5 Simulation Scenarios,

Results and Evaluation

5.1 Introduction

The previous chapter described the implementation of a framework, which is simulated using OPNET Modeler software. In order to evaluate the performance of the designed network framework, we have to simulate a network environment within network simulation software, as it is impossible to test the designed network in the real network environment. This chapter will present the simulation scenarios, results and evaluation performance of the framework from the research project comparing it with the standard network. In the area of computer networking research, OPNET Modeler software is widely and reliably used for testing, debugging and performance evaluation of extended protocol and developing networks.

To prove that this framework can reduce handover latency and reduce packet delay within a wireless network, a variety of network scenarios have been produced. Some scenarios were created to measure robustness on the network, some of them for testing about scalability and so on. Every scenario will be compared with the standard network environment.

5.2 <u>Scenario 1</u>: The Performance of Unicast and Multicast Mechanism

5.2.1 <u>Scenario1</u>: Scenario Description

At first, we evaluated the performance of the unicast and multicast mechanism. The statistical parameters that we focused on are throughput in the links and load at wireless router.

Normally, the streaming video formats can range from 128*120 pixels to beyond 1920*1080 pixels for HD (High Definition) standards. The popular streaming video service such as

YouTube has 320*240 pixel resolution. For the video frame inter-arrival rates can be from 10 up to 30 frames per second [37]. However, the higher the video resolution, the higher the raw video content size, so that means it will affect the bandwidth on the network and packet delays.

In this simulation model, a media server that connects within the network exports the multicast streaming video traffic to the clients. The frame size of video is 128*120 pixels and the video frame inter-arrival rate is 10 frames/sec (fps). In the wireless network, there are four subnets and each subnet has only one wireless router. For the sake of simplicity, we consider that there is only one multicast group in the network. Also, there are fixed and mobile node clients in this scenario. During the simulation, clients can join or leave the multicast group at any time.

5.2.2 <u>Scenario1</u>: Simulation Topology

At the beginning, the network topology combines with wireline and wireless networks as presented in Figure 5-1.

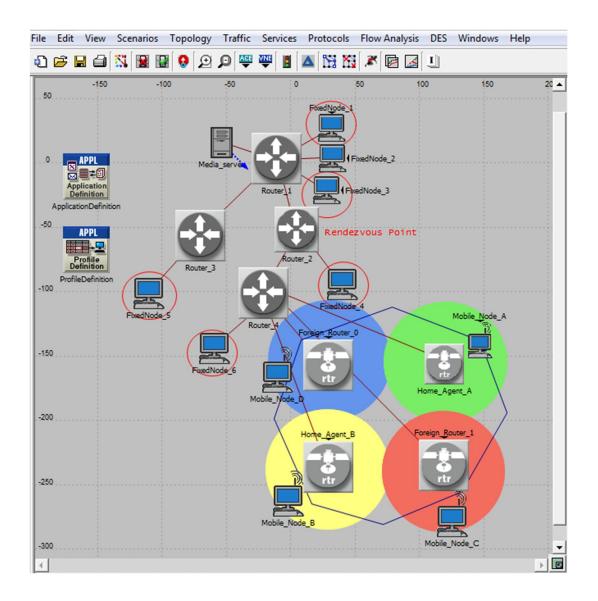


Figure 5-1 Scenario 1: Simulation Topology

In order to examine the performance, we have moved mobile node B and C to connect to the network via Home Agent_A access point. Hence the experiment at this stage is that the mobile nodes A, B and C are connected to Home Agent_A access point and are required to receive the data from the media server.

5.2.3 <u>Scenario1:</u> Simulation Results and Evaluation

In the simulation, all mobile nodes join the multicast group at 100 seconds and leave at the end of the simulation. During simulation time, we assume that all packets are delivered

correctly to all receivers without any disruptions to the service. The correlation between unicast and multicast are tested. The results are presented in Figure 5-2 and 5-3.

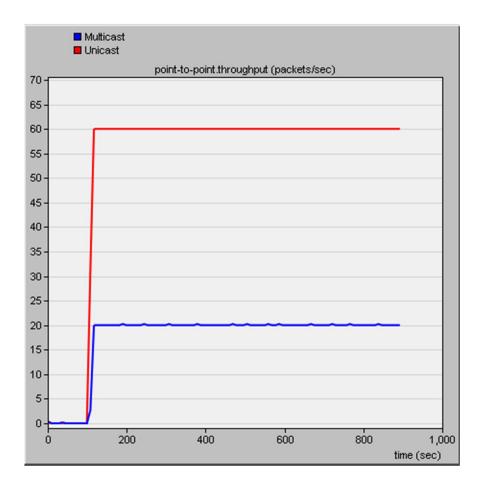


Figure 5-2 Throughputs between Router4 and Home Agent A access point

We can observe from Figure 5-2 that, the throughput of multicast transmission is only one third compared with the throughput in the same links when multiple unicast transmission is used. This means that the bandwidth consumption increases. It is the most important benefit of the multicast data delivery.

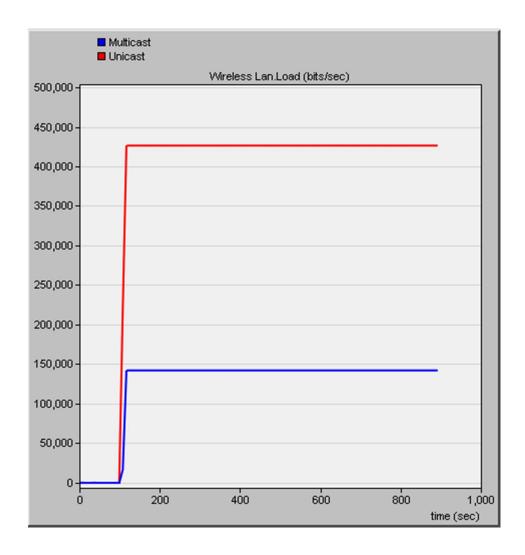


Figure 5-3 Load at Home Agent A access point

Figure 5-3 is shown the load at Home Agent_A access point compared between unicast and multicast transmission in units of bits/sec. It can be seen from the data that multicast mechanism reduces traffic and load at access point. Figure 5-4 illustrates CPU utilization at access point. It is apparent from this graph that CPU utilization of unicast is higher than the multicast mechanism. This result may be explained by the fact that the access point of unicast has to process 3 unicast copies while multicast processes is concerned with only one set of data. However, the other processes such as encapsulate/ decapsulate packet, routing process and forwarding packet are still the same method. This can explain why CPU utilization of unicast differs from multicast by only 0.04 %.

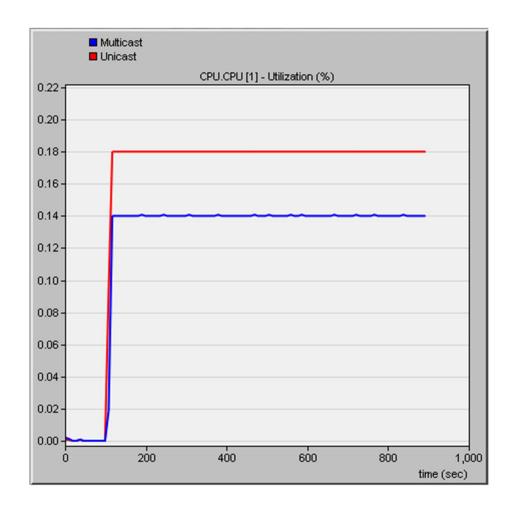


Figure 5-4 CPU Utilization at Home Agent A access point

Another aspect that we examined was when a mobile node moves while receiving multicast data. From Figure 5-1, mobile node A requests to receive video multicast from the media server at 100 seconds. During the simulation, mobile node A moves in a counter-clockwise trajectory roaming through all four access points in the network (the path represented in a blue line in Figure 5-1). At 800s, the mobile node A leaves the multicast group. The other mobile nodes did not move and also received multicast data through the access point as in the figure. Figure 5-5 compares the video traffic received at mobile node A and B.

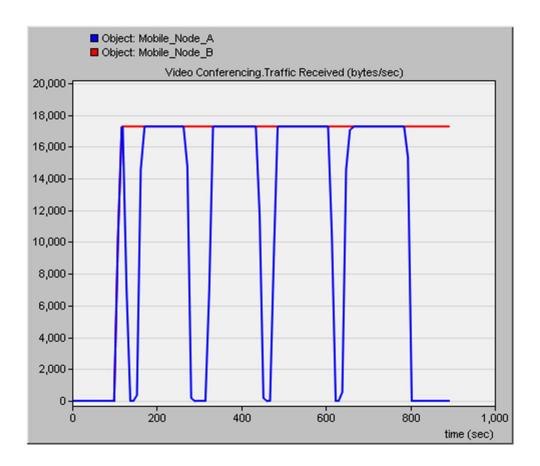


Figure 5-5 Video traffic received at mobile node A compared with B

Figure 5-5 depicts clearly that mobile node B, which is a stationary node, continuously received the multicast traffic until the end of the simulation. On the other hand, mobile node A moves through all four access points, and when handover occurred the traffic received dropped significantly.

In this case, it may completely disconnect the access point from the link layer. Thereafter, it needs to restart the process of performing an IP reconfiguration in the network layer and binding updates to home agent to its infrastructure. Until completion of all these operations the mobile node is likely to experience disruptions or disturbances of application, as the results of packet loss, jitter and delay increase. After the handover process is finished, the node will pick up multicast data again [47].

$$T_{\text{handoff}} = T_{L2} + T_{\text{local-IP}} + T_{\text{BU}}$$
 (5-1)

The T_{BU} is Binding update latency time. In order to examine the multicast service continuity

during the mobile node mobility procedures, therefore we simulated the following scenario. From Figure 5-1 we move the mobile node D to connect to Home Agent _B access point. Now there is no client in Foreign_Router_1 subnet. The mobile node A still moves in the same path. The result of this scenario is shown in Figure 5-6.

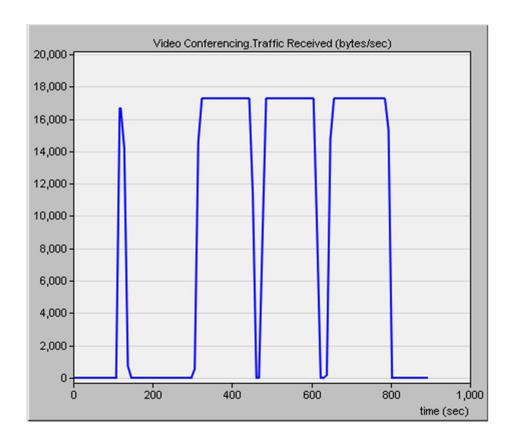


Figure 5-6 Video traffic received at mobile node A

Figure 5-6 is shown that mobile node A did not receive the multicast traffic whilst connected to Foreign_Router_1 zone. This means that mobile node A did not re-join the multicast group during that time. The mobile node A only joins the multicast group once the simulation starts, which is in Home_Agent A zone. Moreover, the mobile node A will receive multicast traffic if that zone already has a member in the same multicast group.

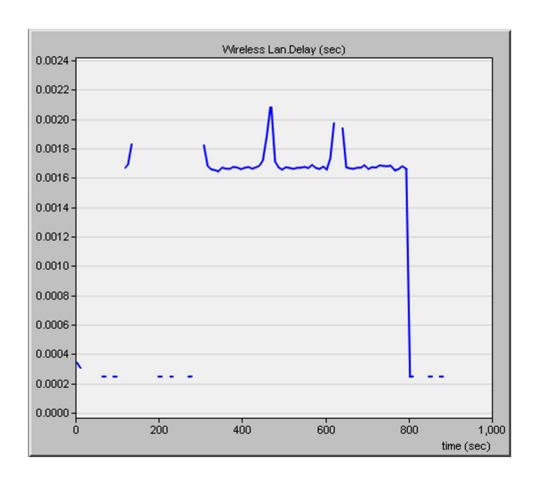


Figure 5-7 Delay at mobile node A

Figure 5-7 is shown the delay on mobile node A. It depicts clearly that the delay increased when the handover process happened. From this result it supports our idea that if we want to reduce the handover process time and also improve the performance, it should have the process of registering the multicast group in advance.

5.3 <u>Scenario 2</u>: Simple Network

5.3.1 Scenario 2: Scenario Description

In the second scenario, we focused on multicast mobility handover occurring between modified framework and comparing with standard network environments. In this scenario we have included the process of multicast re-join in the design framework. Therefore, the aim of this scenario is to measure how much difference of re-join process affects the performance of the network.

5.3.2 Scenario 2: Network Topology

The network topology of this scenario still combines the wireline and wireless network. There are two WiFi zones in this topology, which are Home Agent_A and Foreign_Router_1. In Home_Agent_A zone, Mobile_Node_A is a member in this zone. Mobile_Node_A receives multicast data from Media_Server, and then moves to the Foreign_Router_1 zone. During mobility, Mobile_Node_A still continuously receives multicast packets from the Media_Server. In this scenario, there is no function of Mobile IP protocol involved. Figure 5-8 is shown the network topology of this scenario.

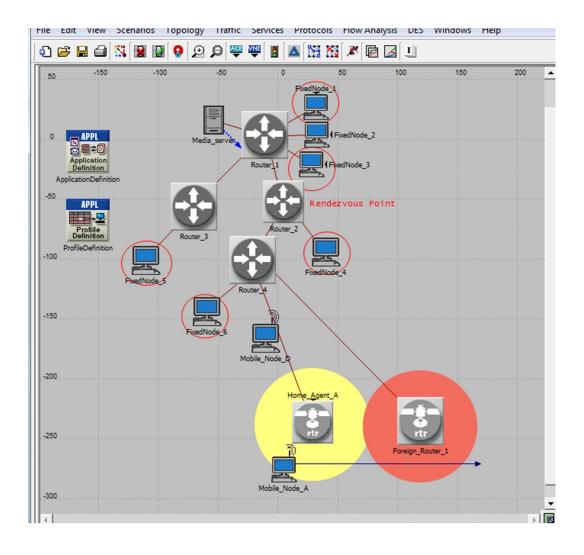


Figure 5-8 Scenario 2: Network Topology

5.3.3 Scenario2: Simulation Results and Evaluation

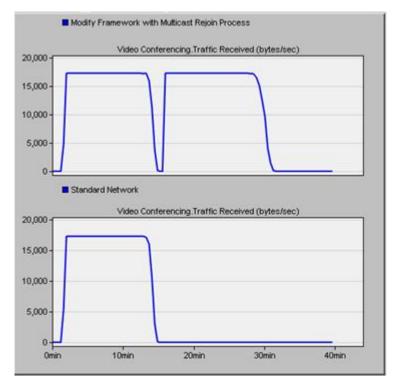


Figure 5-9 Scenario 2: Traffic received at Mobile Node A

The first graph in Figure 5-9 is multicast data that has been received at Mobile_Node_A during the simulation. A gap in the graph happens when Mobile_Node_A moves from Home Agent_A to Foreign Router_1 networks. The result is shown that Mobile_Node_A completely lost connection when handover occurred. However, the modified framework has the multicast re-join process, so a joining message has been sent to Foreign Router_1. Mobile_Node_A can get back to receive multicast data until it leaves multicast group. What is interesting in this data is that there is not the multicast re-join process in the standard network, hence Mobile_Node_A cannot receive traffic when leaving from the home agent network.

5.4 <u>Scenario 3</u>: Mobile IP and multicast Rejoin

5.4.1 Scenario 3: Scenario Description

This scenario intends to inspect the effect of Mobile IP protocol when handover occurs. The infrastructure network of scenario 3 is still similar to the previous scenario. However, some parameters might change. In the modified framework, the Mobile IP protocol has been enabled. Also, there it still has the process of multicast re-join in this scenario. The result will be compared with the standard network.

5.4.2 Scenario3: Simulation Results and Evaluation

The blue graph in Figure 5-10 is the result of the video conferencing traffic that Mobile_Node_A received. It can be seen that it still has a gap of handover. However, the gap is narrower than in scenario 2 and does not completely lose the connection. This is because Mobile IP protocol has a mobility function to support connection. In the red graph, there are no functions of multicast re-join and Mobile IP was not enabled. Hence, the connection was completely lost in this case.

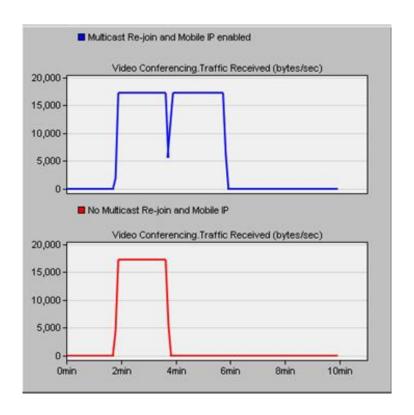


Figure 5-10 Scenario 3: Traffic received at Mobile Node A

5.5 Scenario 4: Care of Address in Advance

5.5.1 Scenario 4: Scenario Description

This scenario focused on the advantage of the process of registering a CoA address in a foreign network in advance. The method can reduce the waiting time in the process of assigning IP address to the network membership. The network topology in this scenario is still the same in scenario 2.

5.5.2 Scenario4: Simulation Results and Evaluation

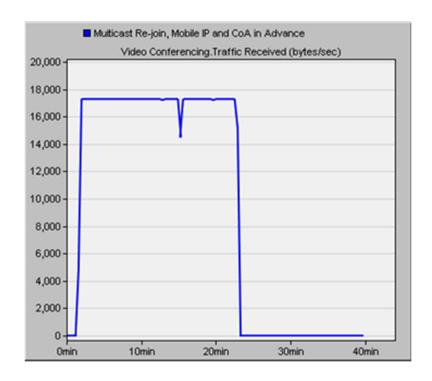


Figure 5-11 Scenario 4: Traffic received at Mobile Node A

The result is shown that when the foreign agent knows the new member node in advance, it helps to reduce the handover latency time on the network. Also, it increases the performance and throughput on the wireless network. The foreign router will establish the connection earlier, including the process of joining and distributing the multicast tree in advance. Furthermore, the method of binding update will happen before the handover occurs.

5.6 Scenario 5: Same Multicast Group

5.6.1 Scenario 5: Scenario Description

The purpose of the current scenario was to determine the performance of the network when the foreign network is already a member of the multicast group.

5.6.2 Scenario5: Network Topology

The network topology of this scenario is slightly different from the previous network. In this infrastructure, the foreign network already has a mobile node, which is a member of the multicast group that Mobile_Node_A wants to join. In Figure 5-12, Mobile_Node_B is already a member of the multicast group from the media server.

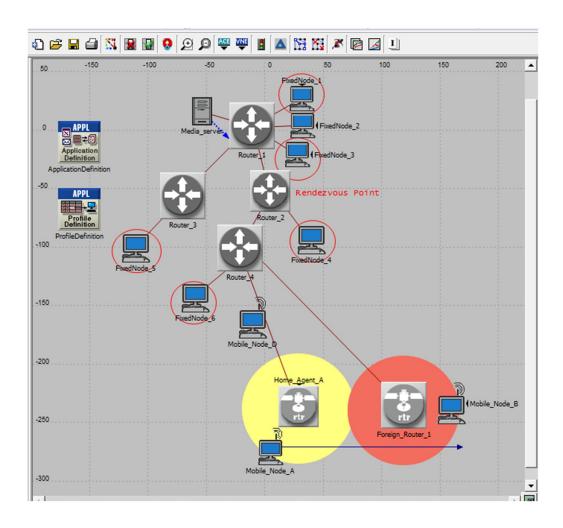


Figure 5-12 Scenario 5: Network Topology

5.6.3 Scenario5: Simulation Results and Evaluation

The graph below illustrates that if the foreign network is already a member of the multicast group the handover latency time will be reduced. This is because the new mobile node just deals with IP addressing and then picks up the multicast signal within the network. There are not any methods for multicast communication.

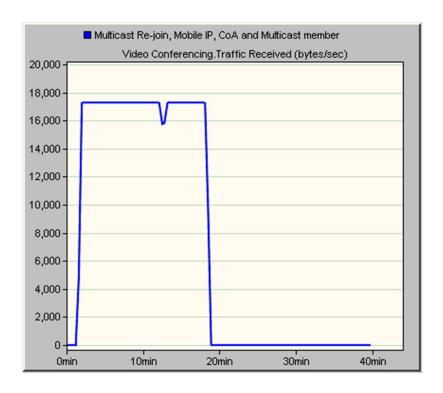


Figure 5-13 Scenario 5: Traffic received at Mobile Node A

5.7 <u>Scenario 6</u>: Multi-Hops

5.7.1 Scenario 6: Scenario Description

In this scenario, we increase a distance between home and foreign network. This study set out to determine the effects of multi-hops on the handover latency and packet delay.

5.7.2 Scenario 6: Network Topology

In the new topology, the foreign agent connects to Router_2. However, they are still connecting to the same multicast group.

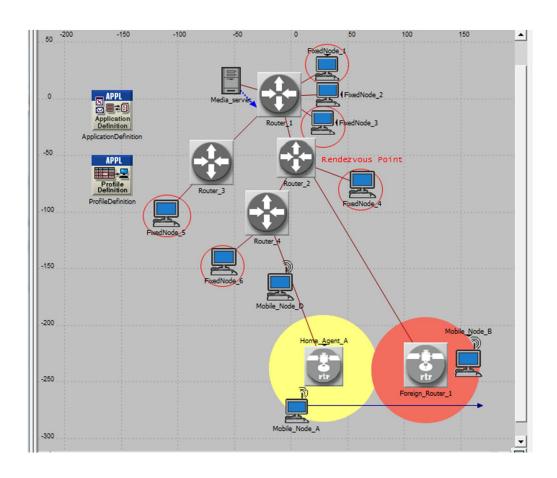


Figure 5-14 Scenario 6: Network Topology

5.7.3 Scenario 6: Simulation Results and Evaluation

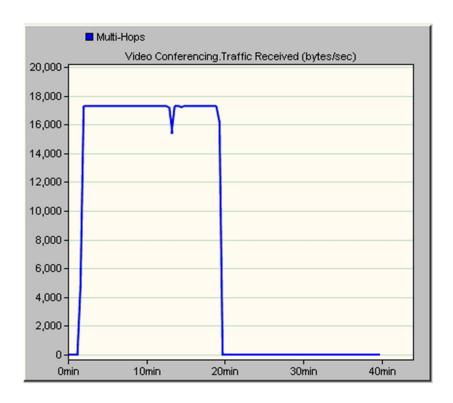


Figure 5-15 Scenario 6: Traffic received at Mobile Node A

The output result on Figure 5-15 is shown that the handover latency time slightly increases because the route path between both networks had changed. However, in this topology the foreign router connected to Router_2, which is one of the members of the multicast tree. Moreover, in this network, Router_2 is a Rendezvous point of the multicast tree. That is why the output result is only slightly different.

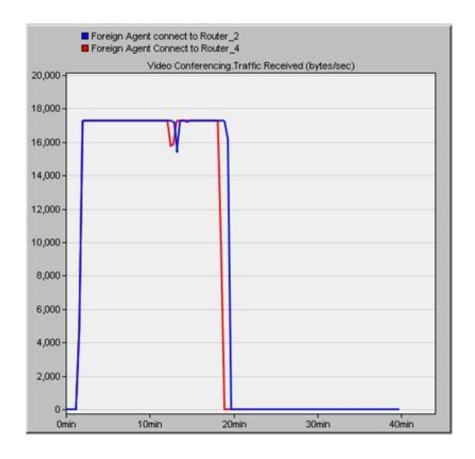


Figure 5-16 Scenario 6: Combined traffic received at Mobile_Node_A

Figure 5-16 combines the graph in Figure 5-13 and 5-15 in order to compare the results. The aim of this graph is to examine how hop distance affects handover latency in our designed framework. It can be seen from the data in the graph that the handover latency time slightly increases in this network topology.

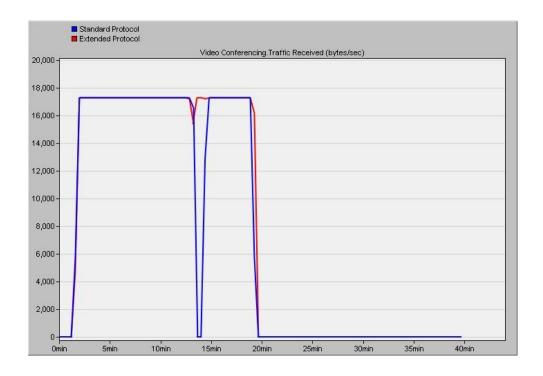


Figure 5-17 Is shown how much the extended protocols can reduce handover latency compared with standard protocols

The graph above illustrates the comparison between the graphs in Figure 5-15, which is the result from the modified framework, compared to the result from the standard framework. In the graph the handover latency of extended protocols is 48 seconds and 120 seconds for standard protocols. What is interesting in this data is that in a standard framework, mobile node completely disconnects until the process of IP readdressing and joining multicast are finished. The results of this study indicate that the modified framework can reduce the handover latency 60% when compared with the standard network.



Figure 5-18 Is shown how much the extended protocols can reduce packet delay compared with standard protocols

From the network topology in Figure 5-14, the purpose of the current study was to assess packet delay in both networks. The above Figure is shown comparison of the output result of packet end-to-end delay in terms of average. It is apparent from this graph that the packet delays in the modified framework are lower than the original framework. On average the modified framework can reduce packet delay by approximately 3.5 - 10 ms throughout the simulation.

5.8 Scenario 7: Handover

5.8.1 Scenario 7: Scenario Description

In this scenario, we increase a scale of network topology for testing performance of proposed framework when mobile node joins in many WiFi networks along the path. Also, this scenario has shown that the mobile node can discover and connect to the new network when the mobile node moves.

5.8.2 Scenario 7: Network Topology

In this topology when the simulation starts, the Mobile_Node_A is a member of Home_Agent_A and receives multicast data from Media_server. Then Mobile_Node_A starts to move in the direction of the blue arrow from Foreign_Router_1 throughout Foreign_Router_4, as presented in Figure 5-19.

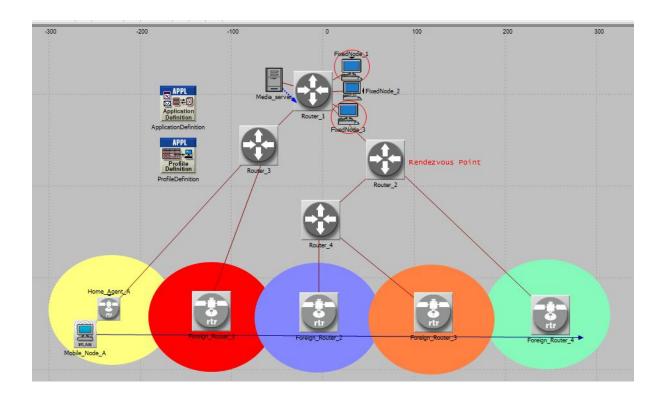


Figure 5-19 Scenario 7: Network Topology

5.8.3 Scenario 7: Simulation Results and Evaluation

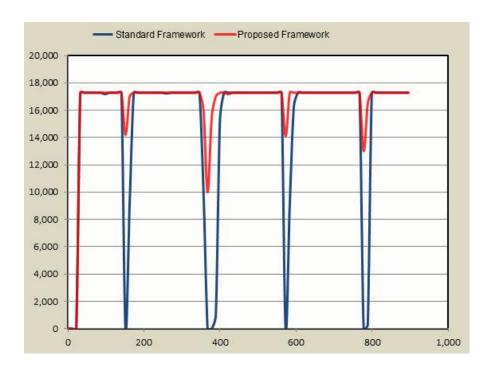


Figure 5-20 Scenario 7: Traffic received at Mobile_Node_A

As Figure 5-20 is shown, there is a significant difference between the two lines. The x-axes show the time in units of seconds. The y-axes represent traffic received in units of bytes/sec. From the graph above we can see that the proposed framework can significantly minimize handover latency more than the standard framework, especially when handover occurred between multi hops such as from Foreign_Router_1 to Foreign_Router_2 in Figure 5-19. In that time, the standard framework completely disconnects and has to restart every process from the beginning while Mobile_Node_A in the proposed framework can receive multicast traffic around 10,000 bytes/sec.

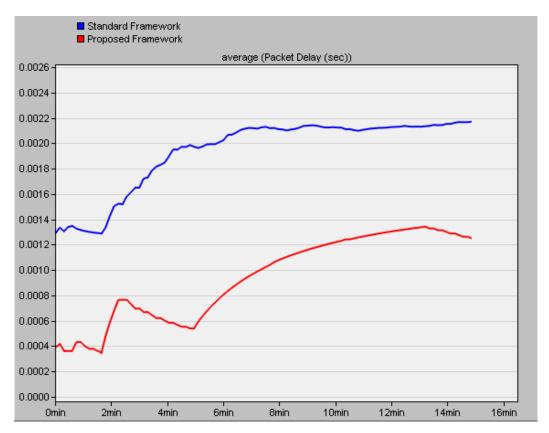


Figure 5-21 Scenario 7: Packet delay

Figure 5-21 represents the end-to-end delay of all the data packets that are successfully received by the mobile node. The graph is shown value in terms of average. There is a clear trend of lower packet delay in the design framework throughout the simulation. Further analysis showed that the designed framework reduced packet delay by approximately 0.7-1.5 ms in this scenario.

5.9 <u>Scenario 8</u>: Multiple Networks

5.9.1 Scenario 8: Scenario Description

In this scenario, we increase a group of WiFi zone to evaluate the performance of the proposed framework in the process of registering CoA address in the foreign network in advance. In this topology Mobile_Node_A will register the CoA address to all of the foreign agents in advance as soon as the mobile node moves into their coverage area. Moreover, this scenario will study how mobile nodes handle all CoA addresses after they are received.

5.9.2 Scenario 8: Network Topology

In this topology when the simulation starts, the Mobile_Node_A is a member of Home_Agent_A and receives multicast data from Media_server. Then Mobile_Node_A starts to move with constant speed in the direction of the blue arrow from Home_Agent_A to Foreign_Router_2 and then moves to Foreign_Router_4, as presented in Figure 5-22. There are 7 WiFi zones in this scenario.

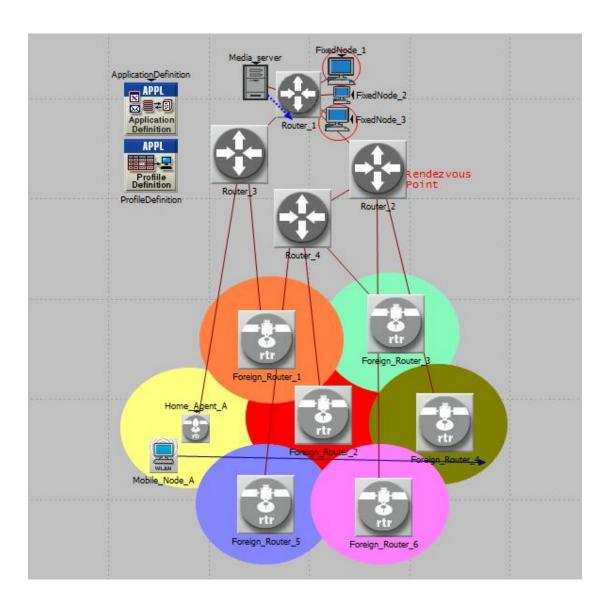


Figure 5-22 Scenario 8: Network Topology

5.9.3 Scenario 8: Simulation Results and Evaluation

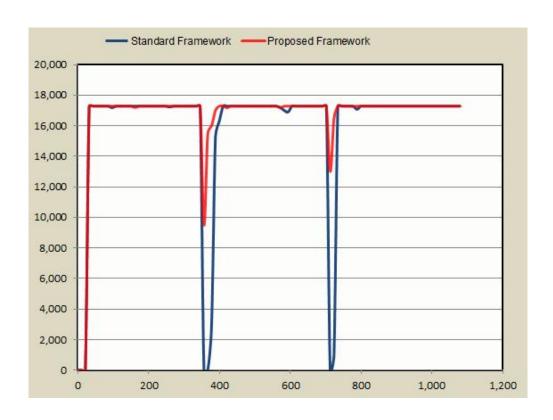


Figure 5-23 Scenario 8: Traffic received at Mobile Node A

The results, as shown in Figure 5-23, indicate that the proposed framework can reduce handover latency for the mobile node. Also, on the designed framework the mobile node continuously received multicast data and did not lose the connection as the standard framework does. In addition, this scenario confirms that the mobile node can store those CoA addresses from foreign agents and use them to connect to the multicast tree in advance. This method will help foreign routers establish the connection and distribute the multicast tree earlier.

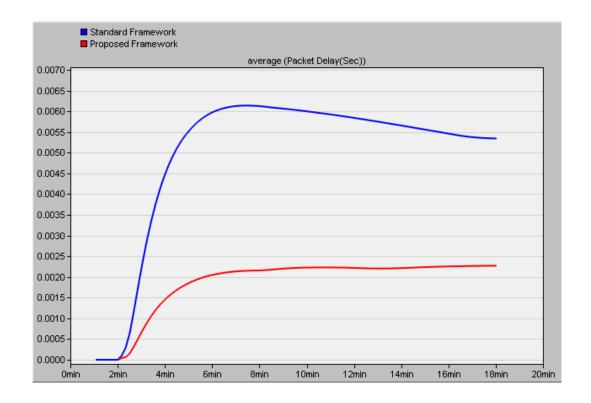


Figure 5-24 Scenario 8: Packet delay

The above figure represents the end-to-end delay on average at the mobile node. Further analysis showed that the designed framework reduced packet delay by approximately 1.5-40 ms in this scenario.

5.10 Scenario 9: Complex Networks

5.10.1 Scenario 9: Scenario Description

To increase the reliability of the proposed framework, we created more groups of WiFi zone to evaluate the performance. Also, in order to study the effectiveness of registering the CoA address in an advance process, the mobile node should register only the foreign network that the mobile node is in the coverage area of. Hence, in this scenario the direction of the mobile node was changed during simulation.

5.10.2 Scenario 9: Network Topology

In this topology when the simulation starts, the Mobile_Node_A is a member of Home_Agent_A and receives multicast data from Media_server. Then Mobile_Node_A starts to move with constant speed in the direction of the blue arrow from Home_Agent_A to Foreign_Router_5 and then moves to Foreign_Router_6, until it stops at Foreign_Router_8 at the end of the simulation, as presented in Figure 5-25.

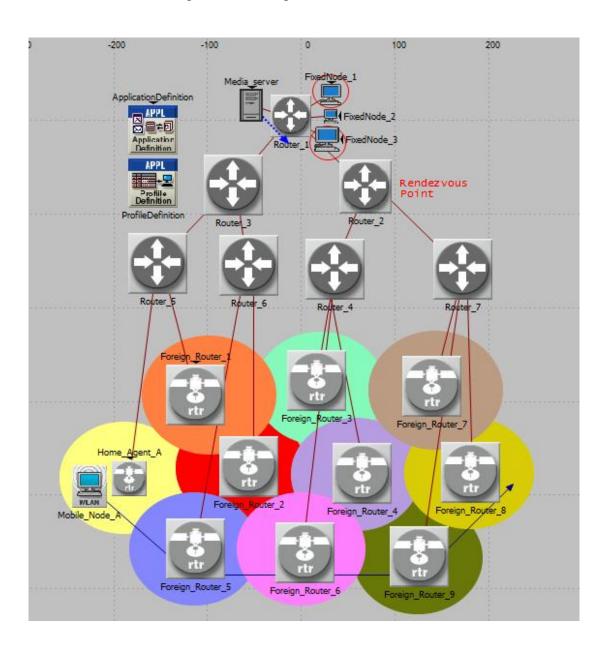


Figure 5-25 Scenario 9: Network Topology

5.10.3 Scenario 9: Simulation Results and Evaluation

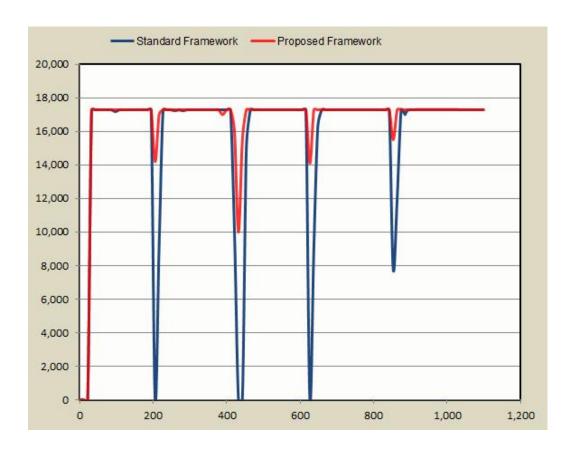


Figure 5-26 Scenario 9: Traffic received at Mobile Node A

The findings of the current network topology are consistent with the previous scenario that the proposed framework can minimize handover latency time. However, the performance of this framework depends on the structure and topology of the network. It can show that the framework can reduce handover latency time that happen in network and transport layer. But it is still create handover latency time that happen in data link layer.

5.11 Scenario 10: Internet

5.11.1 Scenario 10: Scenario Description

To increase the reliability of the proposed framework, we have applied our framework to the

internet. In this scenario, the network will be similar to scenario 9, however, Media server and all access points connect to each other and the multicast tree via the internet.

5.11.2 Scenario 10: Network Topology

The network topology in this scenario is similar to that of scenario 9. The mobile node moves in the same direction. However, in this scenario the mobile node received multicast data from Media_server via the internet. In Figure 5-27, the IP cloud represents the internet. In the simulation we have created some traffic such as email, www, and ftp traffic etc. within the internet as random traffic to make it more like reality. Moreover, the aim of this scenario is to evaluate the performance of the designed framework when it is a part of the internet.

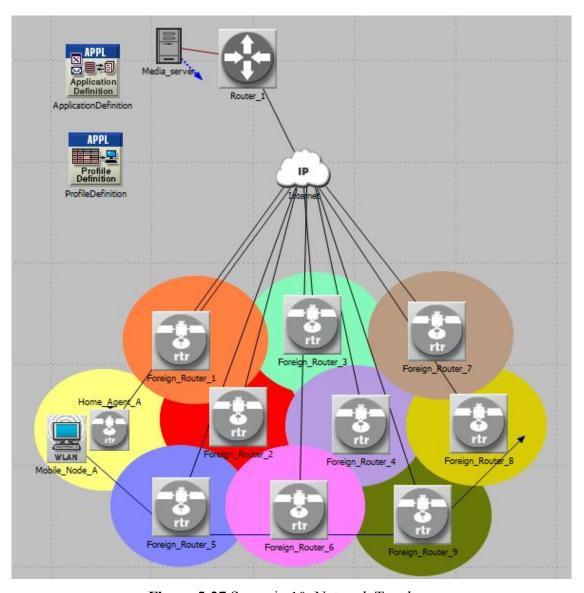


Figure 5-27 Scenario 10: Network Topology

5.11.3 Scenario 10: Simulation Results and Evaluation

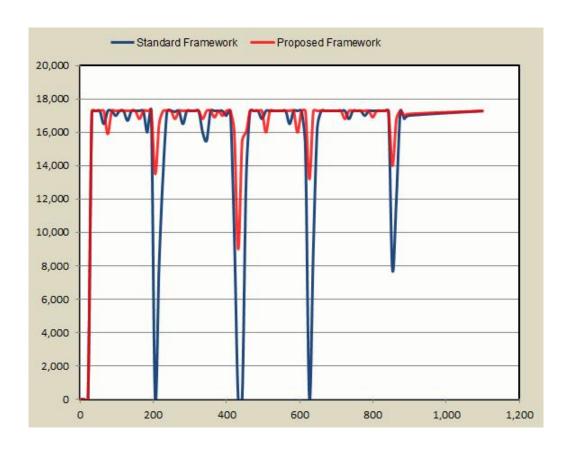


Figure 5-28 Scenario 10: Traffic received at Mobile_Node_A

The current result found that the traffic on the internet slightly affects handover latency time. However, the traffic on the internet directly affects multicast data received at the mobile node because there is traffic on the internet. The most interesting finding was that the proposed framework could be applied for using on the internet and it did not affect the performance.

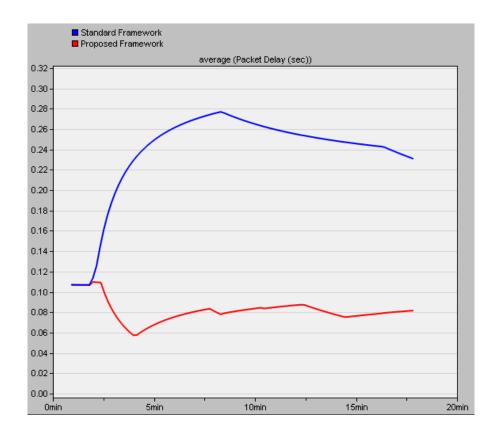


Figure 5-29 Scenario 10: Packet Delay

The above figure represents the end-to-end Packet delay of video conferencing on average at mobile node. The packet delay increased significantly when connected through the internet. Further analysis showed that the designed framework reduced packet delay approximately 80 - 150 ms in this scenario.

Chapter 6 Conclusion and Future Work

6.1 Conclusion

In this thesis, a framework for supporting multicast mobility has been designed. The concept of this framework can support WiFi network both on IPv4 and IPv6 networks. The contribution of this thesis is to find the way to reduce handover latency time, which means including handover delay packet loss and jitter.

However, from the output result on OPNET Medeler network simulation software in chapter 5, it can be confirmed that the designed framework of this research achieves the aim of the research in this thesis

Those methods and processes within the designed framework are a key factor that can produce an output result for achieving the research aim that can be analysed and classified into:

- Foreign agent arranges an IP address for a new mobile node before it becomes a member within the network. This way it helps the foreign agent to reduce the time for searching an available IP address on their database, including negotiation time between foreign router and new mobile node. It is affected from registering CoA in advance module
- As it knows a new mobile node member in advance the foreign router can establish a path to multicast group early.
- The results of this research support the idea that when mobile node connected to the multicast tree early, the process of rebuild multicast tree can happen early also. This saves time for connecting to multicast tree.
- Helps Mobile IP protocol to do the process of Binding update to home agent early.
- The results of this investigation show that mobile node can use the new IP address to join multicast tree in advance due to already having the new IP address.

Reduces time in the process of becoming a member in a new multicast group.

6.2 Recommendation for Future Work

From this research, there are many things that can be improved, modified and developed for making this framework more powerful, such as:

- Include the algorithm or process for predicting potential foreign agent in advance. This is to save the resource and bandwidth on the wireless network. Further work is required to establish this.
- Design a new algorithm or method for finding higher performance of the foreign agent.
- Find the way to store CoA address, in case of receiving lots of CoA in advance.
- Improve the performance of framework in terms of robustness such as when there are lots of mobile nodes within the wireless network.
- Increase the performance of framework in terms of scalability for supporting more mobile nodes. This is an important issue for future research.
- Improve framework for supporting large scale multicast tree.
- Improve framework for supporting a variety of applications and multimedia sizes.
- Extend framework for supporting many-to-many delivery applications. Future studies on this topic are therefore recommended.

Appendices

Appendix A: Essential Kernel Procedures

This appendix presents the most-used kernel procedures (KPs) and functions in OPNET Modeler. They are grouped by the following areas of functionality:

- ➤ Attribute Access
- Distributions
- Dynamic Processes
- > Events and Time
- > Identification and Discovery
- ➤ Interface Control Information (ICIs)
- > Interrupt Processing
- > Packet Generation and Processing
- > Statistic Recording

Attribute Access Get or set attribute values. (Simulation attributes are "global" to the simulation model.)

- op_ima_obj_attr_get_<type> -> completion code <type> = color, dbl, int32, objid, str, toggle
- op_ima_obj_attr_set_<type> -> completion code <type> = color, dbl, int32, objid, str, toggle
- op ima sim attr get $\langle type \rangle \rightarrow completion code \langle type \rangle = color, dbl, int32, str$

Distributions Load distributions by name; Obtain outcomes from loaded distributions.

- op dist load (dist name, dist arg0, dist arg1) -> distribution handle
- op dist outcome (dist ptr) -> outcome
- op dist uniform (limit) -> outcome (between 0.0 and limit)

Dynamic Processes Create a new "child" process of a given type; Destroy a process.

- op pro create (model name, ptc mem ptr) -> process handle
- *op_pro_destroy_options (pro_handle, options)* -> completion code Identify the current process.
- *op_pro_self* () -> handle for this process Invoke another process (cause it to execute now). As an invoked process, get optional state that is passed.
- op_pro_invoke (pro_handle, argmem_ptr) -> completion code
- op pro argmem access() -> argument pointer

Events and Time

Cancel an event.

• op ev cancel (evhandle) -> completion code

Obtain current simulation time.

• op sim time () -> current simulation time in seconds

Terminate simulation.

• op sim end (line0, line1, line2, line3) -> (no return value)

Identification and Discovery

Find the containing object.

• op id self () -> object ID of containing object

Find the parent of an object.

• op topo parent (child objid) -> object ID of parent

Find an object's descendants in the hierarchy.

• op topo child count (parent objid, child type) -> number of children of specified type

• op_topo_child (parent_objid, child_type, child_index) -> object ID of the i'th child meeting criteria

Find an object's peers. "objmtype" is one of an enumerated set; "direction" is IN or OUT. Possible use: how many links am I connected to; then, give me the i'th link.

- *op_topo_assoc_count (objid, direction, objmtype)* -> number of associations of given direction and type
- op_topo_assoc (objid, direction, objmtype, index) -> object ID of the i'th association meeting the direction and type criteria

Interface Control Information (ICIs)

Create or destroy an ICI.

- op ici create (fmt name) -> new ICI
- op ici destroy (iciptr) -> (no return value)

Get or set ICI attribute values.

- op ici attr get <type>, <type> = dbl, int32, int64, ptr -> completion code
- op ici attr set $\langle type \rangle$, $\langle type \rangle = dbl$, int32, int64, ptr \rightarrow completion code

Associate an ICI with a particular interrupt.

• op ici install (iciptr) -> previously installed ICI

Interrupt Processing Schedule an interrupt for this object or another at a given time. Optionally pass a "code".

- op intrpt schedule self (time, code) -> event handle for interrupt
- op intrpt schedule remote (time, code, mod objid) -> event handle for interrupt

Obtain various attributes of the current interrupt.

- op intrpt type () -> type (such as packet arrival, statistic change, self interrupt)
- op intrpt strm () -> stream for packet arrivals
- op intrpt ici () -> control information passed with an interrupt (arbitrary structure)

Packet Generation and Processing

Create, copy, or destroy a packet.

- op pk copy (pkptr) -> pointer to new copy of packet
- *op pk destroy (pkptr)* -> (no return value)

Get or send a packet. (with optional delay)

• op pk send delayed (pkptr, outstrm index, delay) -> (no return value)

Get and set named fields of a packet.

- op_pk_nfd_set_<type>, <type> = dbl, info, int32, int64, objid, pkid, pkt, ptr -> completion code
- op pk nfd get <type>, <type> = dbl, int32, int64, objid, pkid, pkt, ptr -> completion code

Get certain properties of a packet.

- op pk creation time get (pkptr) -> simulation time at which packet was created
- op_pk_total_size_get (pkptr) -> size of packet in bits (sum of field sizes) Insert or remove a packet from a specified subqueue.
- op_subq_pk_remove (subq_index, pos_index) -> pointer to packet removed from the specified subqueue

Statistic Recording Obtain a handle for a statistic, given its name. Type is Global or Local.

Optionally specify an index when a single statistic name encompasses multiple independent time series.

• op_stat_reg (stat_name, stat_index, type) -> statistic handle

Write a new value to a particular statistic. (new value is assumed to be recorded at the current time)

• op_stat_write (stat_handle, value) -> (no return value)

Appendix B: C++ Code

This appendix presents the source code of the research that had been developed in OPNET Modeler simulation software.

```
/* This variable carries the header into the object file */
const char mobile_ip_mn_pr_c [] = "MIL_3_Tfile_Hdr_ 30A op_runsim_dev 7
4F427647 4F427647 1 Khanista-PC Khanista 0 0 none none 0 0 none 0 0 0 0 0 0 0
0 4871 6
#include <string.h>
/* OPNET system definitions */
#include <opnet.h>
/* Header Block */
#include
                "ip rte support.h"
#include "mobile ip_support.h"
#include "ip addr v4.h"
#include "ip mcast_support.h"
#include
                "mobility support.h"
#include "mmobility_support.n"
#include "ip_igmp_support.h"
#include "ip_pim_sm_support.h"
#include "udp_api.h"
#include "ip_icmp_pk.h"
#include "ip_dgram_sup.h"
#include "math.h"
#define IP_PK #define REG_PK
                                                   is ip pk
                                                 is reg pk
#define AD_RECEPTION is_ad_reception
#define HA_AD_RECEPTION is_ha_ad_reception
#define FA_AD_RECEPTION is_fa_ad_reception
#define SOLICITATION_TIME is_solicitation_time
#define VALID_FA_CANDIDATE is_valid_fa_candidate
#define HA_TIMEOUT is_ha_timeout
#define
              FA TIMEOUT
                                                  is fa timeout
#define
              TIMEOUT
                                                  is timeout
#define
              RETRY
                                                  is retry
#define FA_REG_SUCCESS
#define HA_REG_SUCCESS
#define OUT_OF_RETRIES
#define INVALID_REPLY
#define REREGISTER
#define SWITCH_FA
                                                 is fa reg success //Me
                                                 is ha reg success //Me
                                                 is out of retries
                                                 is invalid reply
                                                  is reregister
                                                  is switch fa
#define
              TIMER
                                                   is timer //Me
#define
                SOLICIT
                                                   is solicit //Me
#define
                 MipC MN Rereg Buffer
#define
                 MipC MN Solicit Max Interval 60.0
#define
                 MipC MN Solicit Min Interval
```

```
IP DEFAULT TTL
#define
                                          32
                                          50
          IPC PIM SM RPF TIMER OFFSET
#define
          IPC PIM SM NOT BSR CAND
                                          \cap
#define
#define
          IPC PIM SM TOS
                                          0
           IPC PIM SM DATA RATE TIMER
#define
           IPC PIM SM START
#define
           IPC PIM SM RPF UPDATE
#define
           IPC PIM SM SEND JOIN PRUNE MSG
#define
            IPC PIM SM DR TIMEOUT OFFSET
#define
                                                 1000
#define
            IPC PIM SM SEND HELLO MSG OFFSET
                                                 2000
/***** Transition Macros *****/
#define
          HELLO MSG
                                     (transition code ==
IpC Pim Sm Hello Msg Recvd)
#define JOIN PRUNE MSG
                                     (transition code ==
IpC Pim Sm Join Prune Msg Recvd)
#define DATA PKT
                                     (transition code ==
IpC Pim Sm Data Pkt Recvd)
#define RTE PLUS
                                     (transition code == IpC Pim Sm Rte Plus)
#define RTE MINUS
                                     (transition code == IpC Pim Sm Rte Minus)
#define REGISTER MSG
                                     (transition code ==
IpC Pim Sm Register Msg Recvd)
#define REGISTER STOP MSG
                                     (transition code ==
IpC Pim Sm Register Stop_Msg_Recvd)
#define DATA RATE TIMER EXPD (transition code ==
IpC Pim Sm Data Rate Timer Expd)
#define SEND HELLO MSG
                                     (transition code ==
IpC Pim Sm Send Hello Msg)
#define SEND JOIN PRUNE MSG
                                     (transition code ==
IpC Pim Sm Send Join Prune Msg)
#define DR TIMEOUT
                                     (transition code ==
IpC Pim Sm Dr Timeout)
#define RPF UPDATE
                                     (transition code ==
IpC_Pim_Sm RPF Update)
#define FAILURE RECOVERY
                                     (transition code ==
IpC Pim Sm Failure Recovery)
#define START INTERRUPT
                                     (invmode == OPC PROINV DIRECT)
#define DELAY TIMER EXPD
                                     (transition code ==
IpC Igmp Host Delay Timer Expd)
#define REPORT RECVD
                                     (transition code ==
IpC_Igmp_Host_Report_Recvd)
#define QUERY RECVD
                                     (transition code ==
IpC Igmp Host Query Recvd)
#define LEAVE GROUP
                                     (transition code ==
IpC Igmp Host Leave Grp)
#define JOIN GROUP
                                     (transition code ==
IpC Igmp Host Join Grp)
#define ICMP IP PROCESS INVOKE
                                     (intrpt_type == OPC_INTRPT SELF) &&
#define ECHO REQUEST GEN
(pkt from ip == OPC FALSE)
#define ECHO REQUEST RCVD
                              (icmp message type == IpC Icmp Echo Request)
#define ECHO REPLY RCVD
                             (icmp message type == IpC Icmp Echo Reply)
                                          "0.0.0.0"
#define IP ICMP DEST ADDR UNSPECIFIED
```

```
/***** Macro for Traces ******/
#define LTRACE PIM SM
                                   (op prg odb ltrace active ("pim-sm") ||
op prg odb trace active ())
#define LTRACE PIM SM ALL BUT DATA (op prg odb ltrace active ("pim-
sm all but data"))
                                  (LTRACE PIM SM ||
#define LTRACE PIM SM JOIN PRUNE
LTRACE PIM SM ALL BUT DATA | | op prg odb ltrace active ("pim-sm join prune"))
#define LTRACE PIM SM HELLO (LTRACE PIM SM | |
LTRACE_PIM_SM_ALL_BUT_DATA || op_prg_odb_ltrace_active ("pim-sm_hello"))
#define LTRACE PIM SM TIMERS (LTRACE PIM SM ||
LTRACE PIM_SM_ALL_BUT_DATA || op_prg_odb_ltrace_active ("pim-sm_timers"))
#define LTRACE PIM SM DATA
                                  (LTRACE PIM SM ||
op prg odb ltrace active ("pim-sm data"))
#define LTRACE PIM SM FAIL RECOVER (LTRACE PIM SM ||
op prg odb ltrace active ("pim-sm fail recover"))
                             (op prg odb ltrace active ("igmp") ||
#define LTRACE IGMP
op prg odb trace active ())
static Boolean
                             log call scheduled = OPC FALSE;
static int
                             ip pim sm efficiency mode = -1;
/* Global RP lists.
                     * /
List* bootstrap rp lptr = OPC NIL;
          auto rp lptr = OPC NIL;
          boostrap_support = OPC FALSE;
Boolean
          auto rp agent found = OPC FALSE;
                   ip_igmp_host_config_warn loghndl;
static Log Handle
                      ip igmp host lowlevel error loghndl;
static Log Handle
                      ip igmp host loghndls init = OPC FALSE;
static Boolean
typedef struct
     {
     int
                interval;
     int
                retry;
                req lifetime;
MipT MN Reg Info;
typedef enum
     MipC MN Timer Rereg,
     MipC MN Timer Agent,
     MipC MN Timer Solicit,
     MipC MN Timer Retry
MipC MN Timer;
typedef enum IpT_Icmp_Echo_Message_Type
     IpC Icmp Unspec = -1,
         Icmp Echo Reply = 0,
     IpC Icmp Echo Request = 8
      } IpT Icmp Echo Message Type;
```

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#define IPC ICMP ECHO PKSIZE BITS

```
typedef struct
     Stathandle pkts sent stathandle;
     Stathandle pkts_rcvd_stathandle;
Stathandle resp_time_stathandle;
      } IpT Icmp Stats;
typedef struct
     IpT_Address ip_grp_addr;
                       interface;
      int
                     delay_timer_evh;
delay_timer_id;
     Evhandle
     int
                     timer_on_flag;
unsolicit_msg_count;
     Boolean
     int
     Boolean report sent flag;
     } IpT Igmp Host Grp Elem;
static Log Handle
                       ip igmp host config warn loghndl;
                      ip igmp host lowlevel error loghndl;
static Log Handle
                       ip igmp host loghndls init = OPC_FALSE;
static Boolean
typedef struct
     {
                       address;
     InetT Address
     double
                       lifetime;
                      pref level;
      IpT Interface Info *incoming intf ptr;
MipT MN Agent Info;
typedef struct {
     Evhandle expire_time;
                DUID;
      Objid
            IAID;
      Objid
      unsigned char assign type;
     InetT Address Range assignment;
                                                    /* Holds the v4 addr or
v6 addr/prefix */
                                   link local addr; /* The address used to
     InetT Address
unicast back to the client */
} DhcpT Srv Assignment;
/* Function declarations. */
static void mip mn register (int, MipT MN Agent Info, Boolean);
static void mip_mn_agent_timer_update (double);
static void mip mn tunneled pk stat write (Packet*);
static void mip mn agent solicit pk send (void);
static void mip mn agent cache update (MipT MN Agent Info*, InetT Address,
double, int, MipT Invocation Info*);
static void mip mn ip pk handle (MipT Invocation Info*);
static void mip mn ad packet parse (Packet*, int*, int*, InetT Address*,
int*, int*);
void dhcp parse msg();
     dhcp msg server duid match(PrgT List* rcv pkt opts);
int
```

```
int
     dhcp get free addr(int iface index);
int
      dhcp get free prefix(int iface index);
      dhcp_get_iface_info from index(int index);
int
static void mip mn agent solicit pk send adv (void);//me
                        ip icmp sv init (void);
static void
static void
                        ip icmp ping specs parse (IpT Icmp Temp Ping Specs*
ip temp_ping_specs_ptr);
static void
                        ip icmp initial echo requests schedule (void);
                       ip_icmp_echo_request_packet create (int
static Packet*
req index);//Me
                       ip icmp pkt encapsulate (Packet* icmp req pkptr, int
static Packet*
req index);
                       ip icmp echo reply create (Packet* ip dgram pkptr,
static void
Packet* ip icmp_pkptr, Packet* icmp_reply_pkptr);
static void
                       ip icmp ip process invoke (Packet* ip dgram pkptr);
EXTERN C BEGIN
static void
                       ip icmp ip process invoker (void* state ptr, int
code);//Me
EXTERN C END
static void
                       ip icmp ip process invoke schedule (Packet*
ip dgram pkptr);
                       ip icmp reply stats update (Packet* icmp reply pkptr,
static double
int req index);
                        ip icmp next echo request schedule (int req index);
static void
static void
                        ip_icmp_ping_stats_register (int stat_index, char*
dest host name);
static void
                        ip icmp sim log init ();
static void
                       ip icmp ip dgram discard (Packet* ip dgram pkptr);
EXTERN C BEGIN
static void
                       ip icmp request timeout (void *state ptr, int
index);//Me
EXTERN C END
static IpT Icmp Temp Ping Specs *
                        ip icmp ping traffic list generate (Objid
node objid);
                        ip icmp dgram fdstruct update for reply
static void
(IpT Dgram Fields* ip dgram fd ptr);
                        ip pim sm do init (void);
static void
static void
                        ip pim sm rte plus (void);
static void
                       ip pim sm data pkt (void);
                       ip_pim_sm_join_prune_msg (void);
static void
static void
                       ip pim sm data rate timer expired (void);
                       ip_pim_sm_register_msg (void);
static void
                       ip pim sm register stop msg (void);
static void
static void
                       ip pim sm hello msg (void);
static void
                       ip pim sm rte minus (void);
static void
                       ip pim sm rpf update (void);
                       ip pim sm fail recover (void);
static void
/***** Procedures ********/
static void ip igmp host sv init (void);
static IpT Igmp Host Transition ip igmp host get transition code (Packet**
igmp_pkt_pptr, Boolean* is_gs_query_msg_ptr);
```

```
static IpT Igmp Host Grp Elem*
                                    ip igmp host get grp elem (IpT Address
ip grp addr, int interface);
static void ip igmp host join grp (IpT Address ip grp addr, int interface);
//Me: Joinning multicast group by using CoA address.
static void ip igmp host leave grp (IpT Address ip grp addr, int interface);
static void
                 ip_igmp_host_start_timer_for_grp (IpT_Igmp_Host_Grp_Elem*
grp elem ptr, double max resp time);
                ip igmp host start timer for grps on intf (int interface,
static void
double max resp time);
static void
                 ip igmp host cancel timer for grp (IpT Igmp Host Grp Elem*
grp elem ptr);
static void
                 ip igmp host timer expired (int timer code);
static IpT Igmp Host Grp Elem* ip igmp host grp elem alloc (void);
                 ip igmp host grp elem dealloc (IpT Igmp Host Grp Elem*
static void
grp elem ptr);
                 ip igmp host error (const char* msq1, const char* msq2,
static void
const char* msq3);
static void
                  ip igmp host log handles init (void);
                  ip igmp host log found no grp info (const char*
static void
ip addr str, int ip intf num);
static void
                 ip igmp host grp info print (List* grp lptr, Boolean
short version);
static void
                  ip igmp host igmp msgs sent stat update (OpT Packet Size
size); //Me
EXTERN C BEGIN
static void
      ip igmp host ip process invoke (void *state ptr, int
interface to send);
EXTERN C END
/* End of Header Block */
#if !defined (VOSD NO FIN)
#undef
          BIN
           BOUT
#undef
#define
          BIN
                        FIN LOCAL FIELD( op last line passed) = LINE -
op block origin;
          BOUT BIN
           BINIT FIN LOCAL FIELD( op last line passed) = 0; op block origin
#define
= LINE__;
#else
#define
           BINIT
#endif /* #if !defined (VOSD NO FIN) */
/* State variable definitions */
typedef struct
      /* Internal state tracking for FSM */
      FSM SYS STATE
      /* State Variables */
     MipC Node Type
                                          mip node type
     /* Type of MN I am (MN or MR) */
     MipT Proc Info*
                                          proc info struct ptr
      /\star General info shared between parent and me \star/
     MipT MN Reg Info
                                          reg info
      /* registration information for this MN or MR */
```

```
InetT Address
                                         subnet bcast addr
     /* brodcast address of the mobile interface */
                                         ha address
     InetT Address
     /* IP address of the home agent interface */
     InetT Address
                                         home address
     /* local interface ip address */
     double
                                         time to reregister
     /* sim time for next registration */
     Evhandle
                                         reregister timer ehndl
     /* eventhandle for reregistration */
     int
                                         reg id
     /* identification for pending registration */
     int
                                         retry counter
     /* current number of registration retries */
     InetT Address
                                         agent address
     /* Current agent address serving me */
                                         direct rea
     /* direct with HA or indirect through FA */
     MipT MN Agent Info
                                         latest fa info
     /\star Last FA information on FA other than the currnt one \star/
     Evhandle
                                         agent timer ehndl
     /* event handle for agent timeouts */
     MipT MN Agent Info
                                         latest ha info
     /* Latest info on the HA from advertisements */
     IpT Rte Module Data*
                                        module data
     /* node wide IP module info */
     Stathandle
                                         tunneled pk rcvd sec sh
     Stathandle
                       tunneled bit rcvd sec sh
                       solicit timer ehndl
     Evhandle
     /* Event handle for sending solicitation packet. */
                      solicitation
     Boolean
     /* Whether or not the MN/MR solicit when lost */
                       solicit count
      /* number of times solicitation was sent out before getting an ad */
     Stathandle
                  irdp sent pkts sh
     Stathandle
                       irdp sent bits sh
     Objid
                       node objid
                       simultaneous binding
     /* Flag indicating if this MN/MR will be asking for HA to keep */
     /* simulatneous binding.
                       reg retry timer ehndl
     Evhandle
     /* Event handle for registration retry in case of failure. */
                       g irdp sent bits sh
                       loopback intf
      /* flag indicating if this MN/MR is configured on the loopback
interface */
     int
                       current agent pref level
                                                                        ;
      /* preference level of the current FA */
     IpT Interface Info* current roaming intf
                                                                        ;
     Boolean
                       default gateway
                                                                        ;
      /* Flag to indicate, have a default gateway setup */
     IpT Address last default addr
                                                                        ;
      } mobile ip mn state;
```

```
/* Internal state tracking for FSM */
      FSM SYS STATE
      /* State Variables */
      Objid
                           module objid
                                                                              ;
                           node objid
      Objid
                                                                              ;
      OmsT Pr Handle
                          my proc handle
                                                                              ;
                            udp objid
      Objid
      IpT Rte Module Data*
                            ip support module ptr
      Ici*
                            command ici ptr
      int
                            num dhcp interfaces
      /* Number of interfaces DHCP will be listening on. */
      DhcpT_Srv_Interface* interfaces_config
      int
                             input strm
      int
                             output strm
      InetT_Address
     InetT_Address server_multicast_addr
DhcpT_Stathandles* global_stats
DhcpT_Stathandles* local_stats
                             intrpt code
                             logging
      Boolean
      /* Boolean to indicate if global DHCP logging is turned on/off */
     Log_Handle new_log_handle renew_log_handle
     Log_Handle
Log_Handle
                           error log handle
                                                                              ;
                            expire log handle
      } dhcp server state;
static enum
      {IpC Pim Grp Tbl Grp Addr = 0, IpC Pim Grp Tbl RP Addr,
      IpC_Pim_Grp_Tbl_Src_Addr, IpC_Pim_Grp_Type,
      IpC_Pim_Grp_Tbl_In_Iface, IpC_Pim_Grp_Tbl_Out_Iface,
      IpC Pim Grp Tbl Num Columns
      } IpC Pim Grp Tbl Table Column Index;
typedef struct
                      pim sm status;
      OpT Int8
                      hello period;
               hello_holdtime;
      IpT Pim Intf Stat Handle* stat handle ptr;
                                   hello evhandle;
      Evhandle
      int priority;
      } IpT Pim Intf;
#define mip node type
                                          op sv ptr->mip node type
#define proc info struct ptr
                                         op sv ptr->proc info struct ptr
#define reg info
                                         op sv ptr->reg info
                                         op_sv_ptr->subnet_bcast addr
#define subnet bcast addr
#define ha address
                                         op sv ptr->ha address
#define home address
                                         op sv ptr->home address
                                        op_sv_ptr->time_to_reregister //Me
op_sv_ptr->reregister_timer_ehndl
#define time to reregister
#define reregister timer ehndl
#define reg id
                                         op sv ptr->reg id
#define retry_counter
                                          op_sv_ptr->retry_counter
#define agent address
                                           op sv ptr->agent address
#define direct reg
                                           op sv ptr->direct reg
#define latest fa info
                                           op sv ptr->latest fa info
```

```
#define agent timer ehndl
                                          op sv ptr->agent timer ehndl
                                           op sv ptr->latest ha info
#define latest ha info
#define module data
                                           op sv ptr->module data
#define tunneled pk rcvd sec sh
                                           op sv ptr->tunneled pk rcvd sec sh
#define tunneled bit rcvd sec sh
                                           op sv ptr->tunneled bit rcvd sec sh
#define solicit timer ehndl
                                           op sv ptr->solicit timer ehndl
#define solicitation
                                           op sv ptr->solicitation
#define solicit count
                                           op sv ptr->solicit count
                                           op_sv_ptr->irdp_sent_pkts_sh
#define irdp_sent_pkts_sh
                                          op_sv_ptr->irdp_sent_bits_sh
#define irdp_sent_bits_sh
#define node_objid
                                           op_sv_ptr->node_objid
#define simultaneous binding
                                           op sv ptr->simultaneous binding
#define reg retry timer ehndl
                                           op sv ptr->reg retry timer ehndl
#define g irdp sent bits sh
                                           op sv ptr->g irdp sent bits sh
#define loopback intf
                                           op sv ptr->loopback intf
#define current_agent_pref_level
                                           op sv ptr->current agent pref level
#define current roaming intf
                                           op sv ptr->current roaming intf
                                           op_sv_ptr->default gateway
#define default gateway
#define last default addr
                                          op sv ptr->last default addr
#define module objid
                                          op sv ptr->module objid
#define node objid
                                          op sv ptr->node objid
#define my proc handle
                                           op sv ptr->my proc handle
#define udp objid
                                           op sv ptr->udp objid
#define ip support module ptr
                                           op sv ptr->ip support module ptr
#define command ici ptr
                                           op sv ptr->command ici ptr
#define max sol tmout
                                           op sv ptr->max sol tmout
#define max req tmout
                                           op sv ptr->max req tmout
#define max_req_retries
                                           op_sv_ptr->max_req_retries
#define init renew tmout
                                           op sv ptr->init renew tmout
                                           op sv ptr->max renew tmout
#define max renew tmout
#define init rebind tmout
                                           op sv ptr->init rebind tmout
#define max_rebind_tmout
                                           op_sv_ptr->max_rebind_tmout
#define send iface
                                           op sv ptr->send iface
                                           op sv ptr->last_trans_id
#define last trans id
#define rapid commit
                                           op sv ptr->rapid commit
#define server rapid commit
                                           op sv ptr->server rapid commit
#define server id
                                           op sv ptr->server id
#define srv str
                                           op sv ptr->srv str
#define interfaces_config
                                          op sv ptr->interfaces config
                                          op sv ptr->gateway node
#define gateway node
#define input strm
                                          op sv ptr->input strm
#define output strm
                                          op sv ptr->output strm
#define snd_pkt_ptr
                                          op_sv_ptr->snd_pkt_ptr // Me
#define snd pkt opts
                                          op sv ptr->snd pkt opts // Me
                                          op sv_ptr->snd_msg_type // Me
#define snd msg type
                                           op sv ptr->rcv pkt ptr // Me
#define rcv pkt ptr
                                           op sv ptr->rcv_pkt_opts // Me
#define rcv pkt opts
#define rcv msg_type
                                           op sv ptr->rcv msg type // Me
                                           op sv_ptr->rcv_msg_trans // Me
#define rcv_msg_trans
#define RTprev
                                           op sv ptr->RTprev
#define expire time
                                           op sv ptr->expire time
#define rebind time
                                           op_sv_ptr->rebind_time
#define retrans count
                                           op sv ptr->retrans count
#define intrpt_type
                                           op_sv_ptr->intrpt_type
#define intrpt code
                                           op_sv_ptr->intrpt_code
```

```
#define num interfaces
                                         op sv ptr->num interfaces
#define next evh
                                          op sv ptr->next evh
#define server multicast addr
                                          op sv ptr->server multicast addr
#define node ll addr
                                          op sv ptr->node ll addr
#define global stats
                                         op sv ptr->global stats
#define local stats
                                         op sv ptr->local stats
#define transaction time
                                         op sv ptr->transaction time
#define logging
                                          op sv ptr->logging
#define new log handle
                                         op_sv_ptr->new_log_handle
                                      op_sv_ptr->renew_log_handle
op_sv_ptr->error_log_handle
op_sv_ptr->expire_log_handle
#define renew_log_handle
#define error log handle
#define expire log handle
#define server inet addr
                                         op sv ptr->server inet addr
/* These macro definitions will define a local variable called
/* "op sv ptr" in each function containing a FIN statement. */
/* This variable points to the state variable data structure,
/* and can be used from a C debugger to display their values.
#undef FIN PREAMBLE DEC
#undef FIN PREAMBLE CODE
#define FIN PREAMBLE DEC
                              mobile ip mn state *op sv ptr;
#define FIN PREAMBLE CODE
            op sv ptr = ((mobile ip mn state *)(OP SIM CONTEXT PTR-
> op mod state ptr));
#undef FIN PREAMBLE DEC
#undef FIN PREAMBLE CODE
#define FIN PREAMBLE DEC
                              dhcp client state *op sv ptr;
#define FIN_PREAMBLE CODE
            op sv ptr = ((dhcp client state *)(OP SIM CONTEXT PTR-
> op mod state ptr));
/* Function Block */
#if !defined (VOSD NO FIN)
enum { op block origin = LINE + 2};
/* Transitional Executives */
void SendMsq()
      /* Build a DHCP message with
      * appropriate options as determined by the nodes
       * configuration, and send the packet to the well known
       * DHCP multicast address.
      */
      DhcpT Opt*
                              tmp dhcp opt;
      DhcpT_Cli_Assignment* iface config;
      double
                             retrans time;
                             int code;
      int
                             list cell ptr;
      PrgT List Cell*
                             oro length
                                               = 0;
      Boolean
                              retrans msg
                                               = OPC FALSE;
                              log_str[1000];
      char
```

```
FIN(SendMsq());
      /* Double check if a result of receiving a message from potentially
      many servers, if we've already sent a request to a particular server,
      then ignore this recent message if it's not from the server we sent the
      request to. */
      if (intrpt type == OPC INTRPT STRM)
            tmp dhcp opt = dhcp optlist get opt(DHCPC OPT SERVERID,
rcv pkt opts);
            if(tmp_dhcp_opt->simple_data != server id)
                  Discard();
                  FOUT;
                  }
            }
      /* If retransmitting, determine if this is a normal retransmission,
       * or if we are switching from a Renew phase to a Rebind phase. If
       ^{\star} it is a normal retransmission, we only increment the retransmission
       * count.
      if ((intrpt type == OPC INTRPT SELF) && (intrpt code ==
DHCPC MSG TIMEOUT))
            if((rebind time == 0) \mid | ((rebind time > 0) && (op sim time() <
rebind time)))
                  retrans count++;
                  retrans msg = OPC TRUE;
                  if(LTRACE ACTIVE)
                        op prg odb print minor ("Client timed out waiting for
response. Retransmitting...", OPC NIL);
            else
                  snd msg type = DHCPC MSG REBIND;
                  sprintf(log str, "Current server %s not responding to DHCP
messages. "
                        "Sending Rebind message instead of Renew message to
obtain service "
                        "from any available server.", srv str);
                  LOG(renew log handle, log str, PrgC Log Severity Warning);
                  }
            }
      /* If we received an Advertise, increment the transaction ID and send a
Request: */
      if ((intrpt type == OPC INTRPT STRM) && (rcv msg type ==
DHCPC MSG ADVERTISE))
            snd msg type = DHCPC MSG REQUEST;
            last trans id++;
```

```
/* If our assignment is close to expiring, increment the transaction ID
and send a Renew: */
      if ((intrpt type == OPC INTRPT SELF) && (intrpt code ==
DHCPC ASSIGN TIMEOUT))
            snd msg type = DHCPC MSG RENEW;
            last trans id++;
            sprintf(log str, "Current assignment(s) getting stale. Initiating
Renew "
                  "message exchange with current server %s", srv str);
            LOG(renew log handle, log str, PrgC Log Severity Information);
      /* Now, determine the interval for the next retransmission attempt
        Set the time :: Me*/
      if((snd msg type == DHCPC MSG SOLICIT) || (snd msg type ==
DHCPC MSG SOLICIT RAPID))
            RTprev = dhcp get retrans time(RTprev, DHCPC TRP SOL TIMEOUT,
max sol tmout);
            retrans time = op sim time() + RTprev;
            int code = DHCPC MSG TIMEOUT;
      if(snd msg type == DHCPC MSG REQUEST)
            if (retrans count < max req retries)</pre>
                  RTprev = dhcp get retrans time(RTprev,
DHCPC TRP REQ TIMEOUT, max req tmout);
                  retrans time = op sim time() + RTprev;
                  int code = DHCPC MSG TIMEOUT;
            else
                  /* Immediately fallback to Solicit message by going back to
"Begin" FSM state */
                  op intrpt schedule self(op sim time(), DHCPC MSG FALLBACK);
                  FOUT;
                  }
            }
      if(snd_msg_type == DHCPC MSG RENEW)
            RTprev = dhcp get retrans time(RTprev, init renew tmout,
max renew tmout);
            if((op_sim_time() + RTprev) > rebind time)
                  /* Another scheduled retransmission would exceeded the
Renew phase, so we must switch to the Rebind phase at the rebind time */
                  retrans time = rebind time;
```

```
else
                  retrans time = op sim time() + RTprev;
            int code = DHCPC MSG TIMEOUT;
      if(snd msg type == DHCPC MSG REBIND)
            RTprev = dhcp get retrans time(RTprev, init rebind tmout,
max rebind tmout);
            if((op sim time() + RTprev) > expire time)
                  /* Another scheduled retransmission would exceed the valid
lifetime of the assignment. */
                  retrans time = expire time;
                  int code = DHCPC MSG FALLBACK;
            else
                  retrans time = op sim time() + RTprev;
                  int code = DHCPC MSG TIMEOUT;
            }
      /* Schedule the next timer self-interrupt: */
      if (op ev pending(next evh))
            op ev cancel (next evh);
      next evh = op intrpt schedule self(retrans time, int code);
      /* Now that we know we'll be sending a message (and not falling back,
      * create the message):
      snd pkt ptr = dhcp msg create(snd msg type, last trans id);
      snd pkt opts = dhcp optlist create();
      prg list insert (snd pkt opts,
                        dhcp opt create id(DHCPC OPT CLIENTID, node objid),
                              PRGC LISTPOS TAIL);
      /* If the message have reached a server, include its Server ID, as long
as this message is NOT a Rebind message: */
      if((server id != 0) && (snd msg type != DHCPC MSG REBIND))
            prg list insert (snd pkt opts,
                                    dhcp opt create id(DHCPC OPT SERVERID,
server id),
                                    PRGC LISTPOS TAIL);
      /* If this message is a rapid commit Solicit, add the option: */
      if(snd msg type == DHCPC MSG SOLICIT RAPID)
            prg list insert(snd pkt opts,
                                    dhcp opt create (DHCPC OPT RAPID COMMIT),
                                    PRGC LISTPOS TAIL);
      /* Add an option for every interface we are requesting configuration
info for: */
      if((snd msg type == DHCPC MSG REQUEST) || (snd msg type ==
DHCPC MSG RENEW)
            || (snd msg type == DHCPC MSG REBIND)
```

```
|| (snd msg type == DHCPC MSG SOLICIT RAPID))
            for(list cell ptr = prg list head cell get(&interfaces config);
                  list cell ptr != PRGC NIL;
                  list cell ptr = prg list cell next get(list cell ptr))
                  iface_config = (DhcpT Cli Assignment
*)prg list cell data get(list cell ptr);
                  if(iface config->assign type == DHCPC ASSIGN ADDR)
                        /* We're adding an option to request an address: */
                        tmp dhcp opt = dhcp opt create(DHCPC OPT IA NA);
                  else
                        /* We're adding an option to request a prefix: */
                        tmp dhcp opt = dhcp opt create(DHCPC OPT IA PD);
                  tmp dhcp opt->simple data = iface config->iface index;
                  if(iface config->configured == OPC FALSE)
                        tmp dhcp opt->length = 12;
                  else
                        if(iface config->assign type == DHCPC ASSIGN ADDR)
                              tmp dhcp opt->length =
DHCPC OPT IA NA DATA LEN;
                        else
                              tmp dhcp opt->length =
DHCPC OPT IA PD DATA LEN;
                  /* Add this option to the total length count of the Option
Request option: */
                  oro length += 2;
                  prg list insert(snd pkt opts, tmp dhcp opt,
PRGC LISTPOS TAIL);
            tmp dhcp opt = dhcp opt create(DHCPC OPT ORO);
            tmp dhcp opt->length = oro length;
            prg list insert(snd pkt opts, tmp dhcp opt,
      PRGC LISTPOS TAIL);
      /* Finalize the packet and send it to the UDP module for delivery.*/
      dhcp msg finalize(snd pkt ptr, snd pkt opts);
      op ici install (command ici ptr);
      op pk send(snd pkt ptr, output strm);
      op ici install(OPC NIL);
      /* Update the DHCP stats for the message type sent: */
      if(retrans_msg == OPC TRUE)
            dhcp update cli stat(DHCPC MSG RETRANSMIT, 1);
      else
            /* If this is not a retransmission, reset the transaction time:
* /
```

```
transaction time = op sim time();
      dhcp update cli stat(snd msg type, 1);
      /* Free the received packet memory: */
      if(intrpt_type == OPC INTRPT STRM)
           {
            op pk destroy(rcv pkt ptr);
            dhcp optlist destroy(rcv pkt opts);
      FOUT;
      }
void Config()
      /* Me:Once a mobile node has received a Reply message from access point
       which includes address: CoA address, this function is called to
      configure the assignments on the nodes interfaces on mobile node.
      * /
      DhcpT Opt*
                                    dhcp opt ptr;
                              inet range ptr;
      InetT Address Range*
      DhcpT Cli Assignment*
                              assign ptr;
                                    timers_configured = OPC FALSE;
      int
                                               = 0;
      int
                                    new addr
     int
                                    new prefix
                                                     = 0;
                                    renew_addr = 0;
renew_prefix = 0;
      int
      int
                  log str[1000], tmp str[100], val str[100], val2 str[100];
      PrgT List Cell*
                                    list cell opt ptr;
      PrgT List_Cell*
                                    list cell assign ptr;
      double
                                    trans delay;
      PrgT List
                                    *dyn assignment lptr = OPC NIL;
                                    dyn assignment count = 0;
      int.
      IpT Dynamic Assignment *dyn assignment ptr;
      IpT Dynamic Assignment Array *dyn array ptr;
      IpT Dynamic Assignment Type
                                    dyn assignment type;
                                          default route action =
DHCPC DEFAULT ROUTE UNCHANGED;
                                    *rcvd ici ptr;
      InetT Address
                                    *rcvd inet address;
      FIN(Config());
      /* Me: Calculate the transaction delay for this message exchange:- some
      factor of assigning CoA address delay */
      trans delay = op sim time() - transaction time;
      /* Me: Get the server address from access point:- the address of
      foreign agent */
      rcvd ici ptr = op ev ici (op ev current ());
      op ici attr get (rcvd ici ptr, "rem addr", &rcvd inet address);
      /* Cycle through the options to find the IA NA and IA PD options: */
      for(list cell opt ptr = prg list head cell get(rcv pkt opts);
```

```
list cell opt ptr != PRGC NIL;
            list cell opt ptr = prg list cell next get(list cell opt ptr))
            dyn assignment type = IpC Dynamic Assignment None;
            dhcp opt ptr = (DhcpT Opt
*)prg list cell data get(list cell opt ptr);
            if((dhcp opt ptr->code != DHCPC OPT IA NA) && (dhcp opt ptr->code
!= DHCPC OPT IA PD))
                  continue;
            if(dhcp opt ptr->simple data2 == DHCPC NO ADDR AVAIL)
                  /* If this option indicated no address available, skip it
to go to the next option: */
                  sprintf(log str, "Option received from server %s indicated
no address available.", srv str);
                  LOG(error log handle, log str, PrgC Log Severity Notice);
                  continue;
            else if(dhcp opt ptr->simple data2 == DHCPC NO PREFIX AVAIL)
                  /* If this option indicated no prefix available, skip it to
go to the next option: */
                  sprintf(log str, "Option received from server %s indicated
no prefix available.", srv str);
                  LOG(error log handle, log str, PrgC Log Severity Notice);
                  continue;
            /* Me: In this point, the mobile node received the valid address
in the complex data */
            inet range ptr = (InetT Address Range *) (dhcp opt ptr-
>complex data);
            /* Me: Store all info for this particular assignment, have to
check interface by cycling through the array of interfaces and comparing the
interface index with the value we received as the IAID.*/
            for(list cell assign ptr =
prg list head cell get(&interfaces config);
                  list cell assign ptr != PRGC NIL;
                  list cell assign ptr =
prg list cell next get(list cell assign ptr))
                  assign_ptr = (DhcpT Cli Assignment
*)prg list cell data get(list cell assign ptr);
                  if(assign ptr->iface index == dhcp opt ptr->simple data)
                        break;
            if(list cell assign ptr == PRGC NIL)
                  op sim end ("Error while attempting to configure client:",
                                    "Unable to determine interface for
assignment.", OPC NIL, OPC NIL);
```

```
/* Me: If mobile node already have an assignment on this
interface, determine if it's a renewal, or a different assignment (potentially
from a different server or access point). */
            if(assign ptr->configured == OPC TRUE)
                  if (inet address range equal (& (assign ptr->assignment),
inet range ptr) != OPC TRUE)
                        /* Me: Set the IP notification to update */
                        dyn assignment type =
IpC Dynamic Assignment Addr Update;
                        /* Me: Just for logging: */
                        inet address range print(val str, &(assign_ptr-
>assignment));
                        inet address range print (val2 str, inet range ptr);
                        if(assign ptr->assign type == DHCPC ASSIGN ADDR)
                              sprintf(tmp str, "address");
                        else
                              sprintf(tmp str, "prefix");
                        sprintf(log str, "Received new %s assignment for %s:
%s\n Old assignment: %s\n "
                              "New server: %s\n Lifetime: %d\n Transaction
time: %f sec",
                              tmp str, assign ptr->iface name, val2 str,
val str, srv str,
                              dhcp opt ptr->simple data2, trans delay);
                        LOG(new_log handle, log str,
PrgC Log Severity Notice);
                        /* This is a different assignment from what we
already have destroy the existing assignment before installing the new one:*/
                        inet address range destroy(&(assign ptr-
>assignment));
                        assign ptr->assignment = inet address range ptr copy(
      (InetT Address Range *) dhcp opt ptr->complex data);
                        /* Me: Increment the count of configured
addresses/prefixes CoAs: */
                        if (dhcp opt ptr->code == DHCPC OPT IA NA)
                              new addr++;
                        else if (dhcp opt ptr->code == DHCPC OPT IA PD)
                              new prefix++;
                              /* Also inform IP that the default route will
need to be updated */
                              default route action =
DHCPC DEFAULT ROUTE UPDATE;
                  else
                  /* This is a successful renew of an existing assignment. */
```

```
/* No IP notification needed */
                        /* For logging: */
                        inet address range print (val str, & (assign ptr-
>assignment));
                        if(assign ptr->assign type == DHCPC ASSIGN ADDR)
                              sprintf(tmp str, "address");
                        else
                              sprintf(tmp str, "prefix");
                        sprintf(log str, "Renewing existing %s assignment for
%s: %s\n "
                              "From server: %s\n Lifetime: %d\n Transaction
time: %f sec",
                              tmp str, assign ptr->iface name, val str,
srv str, dhcp opt ptr->simple data2, trans delay);
                        LOG(renew log handle, log str,
PrgC Log Severity Information);
                        /* Write appropriate stats: */
                        if(dhcp_opt_ptr->code == DHCPC OPT IA NA)
                              renew addr++;
                        else if(dhcp opt ptr->code == DHCPC OPT IA PD)
                              renew prefix++;
                        }
                  }
            /* If this interface has not yet been configured, then copy the
option data containing the assignment into our local assignment array.
            if(assign_ptr->configured != OPC TRUE)
                  /* Set IP notification to create a new assignment */
                  if (assign ptr->assign type == DHCPC ASSIGN ADDR)
                        dyn assignment type =
IpC Dynamic Assignment Addr Create;
                  else
                        dyn assignment type =
IpC Dynamic Assignment Prefix Create;
                  assign ptr->assignment = inet address range ptr copy(
      (InetT Address Range *) dhcp opt ptr->complex data);
                  /* For logging: */
                  inet address range print(val str, &(assign ptr-
>assignment));
                  if(assign ptr->assign type == DHCPC ASSIGN ADDR)
                        sprintf(tmp str, "address");
                  else
                        sprintf(tmp str, "prefix");
            sprintf(log str, "Obtained initial %s assignment for %s: %s\n "
            "From server: %s\n Lifetime: %d\n Transaction time: %f sec",
            tmp str, assign ptr->iface name, val str, srv str, dhcp opt ptr-
>simple data2, trans delay);
            LOG(new log handle, log str, PrgC Log Severity Information);
```

```
/* Write appropriate stats: */
                  if (dhcp opt ptr->code == DHCPC OPT IA NA)
                        new addr++;
                  else if (dhcp opt ptr->code == DHCPC OPT IA PD)
                        new_prefix++;
                        /* Inform IP to set up a new default route */
                        default route action = DHCPC DEFAULT ROUTE ADD;
                  }
            /* This interface is now configured: */
            assign ptr->configured = OPC TRUE;
            /\star Me : Set DHCP timers, the timers are only set once per
message. The protocol allows for each individual assignment to have it's
             * own timers/expiration times. */
            if(timers configured == OPC FALSE)
                  /* Cancel any pending events: */
                  if(op ev pending(next evh)) { op ev cancel(next evh); }
                  /* Store times needed to properly implement protocol
timers: */
                  expire time = (op sim time()) + dhcp opt ptr->simple data2;
                  rebind time = (op sim time()) + (.8 * dhcp opt ptr-
>simple data2);
                  /* Schedule the interrupt to start sending renews: */
                  next evh = op intrpt schedule self(
                        op_sim_time() + (.5 * dhcp_opt_ptr->simple_data2),
DHCPC ASSIGN TIMEOUT);
                  /* Reset the retransmission timer: */
                  RTprev = 0;
                  timers configured = OPC TRUE;
            /* Notify IP process of interface assignments: */
            if (dyn assignment type != IpC Dynamic Assignment None)
                  /* - Build the IpT Dynamic Assignment structure */
                  dyn assignment ptr = (IpT Dynamic Assignment *)
op prg mem alloc (sizeof (IpT Dynamic Assignment));
                  dyn assignment ptr->assignment type
                                                             = (assign ptr-
>assign type == DHCPC ASSIGN ADDR) ? IpC Dynamic Assignment Addr Create :
IpC_Dynamic_Assignment Prefix Create;
                  dyn assignment ptr->intf index
assign ptr->iface index;
                  dyn assignment ptr->dynamic addr range
                        inet address range ptr copy (inet range ptr);
                  /* - Append to list of dynamic assignments */
```

```
if (dyn assignment lptr == OPC NIL)
                        dyn assignment lptr = op prg list create ();
                  op prg list insert (dyn assignment lptr,
dyn assignment ptr, OPC LISTPOS TAIL);
            }
      /* If we were unable to configure ANY interfaces, fallback to
       * the begin state, after a waiting period. Otherwise, calculate the
       * total time of this transaction and record it in the stats.
      if(new addr + new prefix + renew addr + renew prefix == 0)
            next evh = op intrpt schedule self(op sim time() + max sol tmout,
DHCPC MSG FALLBACK);
            sprintf(log str, "The message from server %s did not contain any
configuration information. "
                  "Soliciting service from any available server", srv str);
            LOG(error log handle, log str, PrgC Log Severity Error);
      else
            /* Record the stats for the number of addr/prefixes new or
renewed,
             * and the total transaction time:
             */
            if (new addr)
                  dhcp update cli stat(DHCPC COUNT NEW ADDR, new addr);
            if(new prefix)
                  dhcp update cli stat (DHCPC COUNT NEW PREFIX, new prefix);
            if(renew addr)
                  dhcp update cli stat(DHCPC COUNT RENEW ADDR, renew addr);
            if(renew prefix)
                  dhcp update cli stat (DHCPC COUNT RENEW PREFIX,
renew prefix);
            dhcp update cli stat(DHCPC TRANSACTION DELAY, trans delay);
      /* Notify IP by sending a remote interrupt with all assignments: */
      if (dyn assignment lptr != OPC NIL)
            void *prev state;
            int i;
            /* Remove an old default route */
            if (default route action == DHCPC DEFAULT ROUTE UPDATE)
                  /* - Build the IpT Dynamic Assignment structure */
                  dyn assignment ptr = (IpT Dynamic Assignment *)
op prg mem alloc (sizeof (IpT Dynamic Assignment));
                  dyn_assignment ptr->assignment type
IpC Dynamic Assignment Default Route Add;
                  dyn assignment ptr->intf index
send iface;
```

```
dyn assignment ptr->dynamic addr range
inet address range create (server inet addr, \overline{0});
                  /* - Append to list of dynamic assignments */
                  op prg list insert (dyn assignment lptr,
dyn assignment ptr, OPC LISTPOS TAIL);
            /* Add the default route if necessary */
            if (default route action >= DHCPC DEFAULT ROUTE ADD)
                  /* Make the received address the new server address */
                  server inet addr = inet address copy (*rcvd inet address);
                  /* - Build the IpT Dynamic Assignment structure */
                  dyn assignment ptr = (IpT Dynamic Assignment *)
op prg mem alloc (sizeof (IpT Dynamic Assignment));
                  dyn assignment ptr->assignment type
IpC Dynamic Assignment Default Route Add;
                  dyn assignment ptr->intf index
send iface;
                  dyn assignment ptr->dynamic addr range
inet_address_range_create (server inet addr, \overline{0});
                  /* - Append to list of dynamic assignments */
                  op prg list insert (dyn assignment lptr,
dyn assignment ptr, OPC LISTPOS TAIL);
            /* - Convert list of dynamic assignments to
IpT Dynamic Assignment Array */
            dyn array ptr = (IpT Dynamic Assignment Array*) op prg mem alloc
(sizeof (IpT Dynamic Assignment Array));
            dyn_array_ptr->assignment_count = op_prg_list_size
(dyn assignment lptr);
            dyn array ptr->assignments = (IpT Dynamic Assignment **)
                  op prg mem alloc (dyn array ptr->assignment count * sizeof
(IpT Dynamic Assignment *));
            for (i = 0; i < dyn array ptr->assignment count; i++)
                  dyn array ptr->assignments [i] = (IpT Dynamic Assignment *)
                        op prg list remove (dyn assignment lptr,
OPC LISTPOS HEAD);
            /* - Set assignment array as event state */
            prev state = op ev state install (dyn array ptr, OPC NIL);
            /* - Schedule a remote interrupt for ip dispatch */
            op intrpt schedule remote (op_sim_time (),
IPC DYNAMIC ASSIGNMENTS INTRPT CODE,
                  ip support module ptr->module id);
            op ev state install (prev state, OPC NIL);
            /* Destroy the list structure */
```

```
prg list destroy (dyn assignment lptr, OPC FALSE);
      /* Free all memory associated with this packet: */
      op pk destroy(rcv pkt ptr);
      dhcp optlist destroy(rcv pkt opts);
      FOUT;
void
dhcp get packet(void) /* Me: When mobile node received DHCP message*/
      FIN(dhcp get packet());
      /* Consume the packet from the input stream: */
      rcv pkt ptr = op pk get(input strm);
      /* Parse the message to obtain values for the state variables: */
      dhcp msg parse(rcv pkt ptr, &rcv msg type, &rcv msg trans);
      /* Get the options from the received packet: */
      op pk fd get ptr (rcv pkt ptr, DHCPC PK FIELD OPTIONS,
(void**) &rcv pkt opts);
      /* If this is a Reply message with a Rapid Commit option
       * change the received message type:*/
      if((rcv msg type == DHCPC MSG REPLY)
            && (dhcp optlist get opt(DHCPC OPT RAPID COMMIT, rcv pkt opts))
!= OPC NIL)
            rcv msg type = DHCPC MSG REPLY RAPID;
      /* Update the local stats for the message type received: */
      dhcp update cli stat(rcv msg type, 1);
      FOUT;
      }
/* End of Function Block */
void
dhcp client (OP SIM CONTEXT ARG OPT)
#if !defined (VOSD NO FIN)
     int op block origin = 0;
#endif
      FIN MT (dhcp client ());
            /* Temporary Variables */
            Objid cli params objid, timers objid, timers row objid, objid1;
            char tmp str[TMP STR SIZE];
            char val str[TMP STR SIZE];
                        tmp int;
            int
                        initial solicit delay;
            double
            /* IP address manipulations: */
```

```
InetT Addr Family addr fam;
            DhcpT Opt* dhcp opt ptr;
           /* End of Temporary Variables */
           FSM ENTER ("dhcp client")
           FSM BLOCK SWITCH
                  /** state (Begin) enter executives **/
                  FSM STATE ENTER FORCED (0, "Begin", state0 enter exec,
"dhcp client [Begin enter execs]")
                       FSM PROFILE SECTION IN ("dhcp client [Begin enter
execs]", state0 enter exec)
                        /* Me: Initialize protocol related info that is reset
every time we restart a Solicit message exchange:*/
                       server rapid commit = OPC TRUE;
                        retrans count
                                              = 0;
                        server id
                                               = 0;
                       RTprev
                                                      = 0;
                       rebind time
                                               = 0;
                       expire time
                                               = 0;
                                               = OPC NIL;
                        rcv pkt ptr
                                               = OPC NIL;
                        rcv pkt opts
                        /* Create the DHCP message to be sent: */
                        if(Rapid Commit)
                             snd msg type = DHCPC MSG SOLICIT RAPID;
                        else
                              snd msg type = DHCPC MSG SOLICIT;
            snd pkt ptr = dhcp msg create(snd msg type, last trans id++);
                        /* Initialize the DHCP options list to be sent: */
                        snd pkt opts = dhcp optlist create();
                        /* Always insert our Client ID in every message: */
                        prg list insert(snd pkt opts,
      dhcp opt create id(DHCPC OPT CLIENTID, node objid),
                                                PRGC LISTPOS TAIL);
                        if (Rapid Commit)
                              DhcpT Opt*
                                                           tmp dhcp opt
      = OPC NIL;
                              DhcpT Cli Assignment* iface config =
OPC NIL;
                              int
                                                                  oro length
     = 0;
                              int
                                                                  rep;
                              /* Insert the Rapid Commit option: */
                              prg list insert(snd pkt opts,
```

```
dhcp opt create (DHCPC OPT RAPID COMMIT),
                                                       PRGC LISTPOS TAIL);
                              /* Insert an option for every interface we're
requesting configuration for: */
                              for(rep = 0; rep < num interfaces; rep++)</pre>
                                     iface config = (DhcpT Cli Assignment
*)prg list access(&interfaces config, rep);
                                     if(iface config->assign type ==
DHCPC ASSIGN ADDR)
                                           /* Requesting an address: */
                                           tmp dhcp opt =
dhcp opt create (DHCPC OPT IA NA);
                                     else
                                           /* Requesting a prefix: */
                                           tmp dhcp opt =
dhcp opt create(DHCPC OPT IA PD);
      tmp dhcp opt->simple data = iface config->iface index;
                                     /* Set the length of this option: */
                                     tmp dhcp opt->length = 12;
                                     /* Add this option to the total length
count of the Option Request option: */
                                     oro length += 2;
                                     prg list insert(snd pkt opts,
tmp dhcp opt, PRGC LISTPOS TAIL);
                              tmp dhcp opt = dhcp opt create(DHCPC OPT ORO);
                              tmp dhcp opt->length = oro length;
                              prg list insert (snd pkt opts, tmp dhcp opt,
      PRGC LISTPOS TAIL);
                        dhcp msg finalize(snd pkt ptr, snd pkt opts);
      /* Schedule a timer self-interrupt in case we need to retransmit: */
                        RTprev = dhcp get retrans time(RTprev,
DHCPC TRP SOL TIMEOUT, max sol tmout);
                        next_evh = op_intrpt_schedule_self(op sim time() +
RTprev, DHCPC MSG TIMEOUT);
                        /* Me: Send the message to the multicast address,
with an initial delay for this first Solicit message: */
                        initial solicit delay =
op dist uniform(DHCPC TRP SOL MAX DELAY);
                        op ici install(command ici ptr);
                        op pk send delayed(snd pkt ptr, output strm,
initial solicit delay);
                        op ici install (OPC NIL);
```

```
/* Log this initial transmission: */
                       if(logging == OPC TRUE)
                             op prg log handle severity set(&new log handle,
PrgC Log Severity Information);
                             op prg log entry write t(new log handle,
op sim time() + initial solicit delay,
                                  "Soliciting any server for new
configuration information");
                       /* Note this delayed transmission in the stats: */
                       if(Rapid Commit)
                             op stat write t(global stats-
>msg count solicit rapid, 1.0, op sim time() + initial solicit delay);
                             op stat write t(local stats-
>msg count solicit rapid, 1.0, op sim time() + initial solicit delay);
                       else
                             op stat write t(global stats-
>msg count solicit, 1.0, op sim time() + initial solicit delay);
                             op stat write t(local stats->msg count solicit,
1.0, op sim time() + initial solicit delay);
                       /* Record the start time of this transaction: */
                       transaction time = op sim time();
                       FSM PROFILE SECTION OUT (state0 enter exec)
                 /** state (Begin) exit executives **/
                 FSM STATE EXIT FORCED (0, "Begin", "dhcp client [Begin exit
execs]")
                 /** state (Begin) transition processing **/
                 FSM PROFILE SECTION IN ("dhcp client [Begin trans
conditions]", state0 trans conds)
                 FSM INIT COND (Rapid_Commit)
                 FSM TEST COND (!Rapid Commit)
                 FSM TEST LOGIC ("Begin")
                 FSM PROFILE SECTION OUT (state0 trans conds)
                 FSM_TRANSIT_SWITCH
                       FSM CASE TRANSIT (0, 1, state1 enter exec, ;,
"Rapid_Commit", "", "Begin", "Wait_Reply", "tr_12", "dhcp_client [Begin ->
Wait Reply : Rapid Commit / ]")
                       FSM CASE TRANSIT (1, 2, state2 enter exec, ;,
"!Rapid Commit", "", "Begin", "Wait Advertise", "tr 13", "dhcp client [Begin
-> Wait Advertise : !Rapid Commit / ]")
            /*----*/
```

```
/** state (Wait Reply) enter executives **/
                  FSM STATE ENTER UNFORCED (1, "Wait Reply",
state1 enter exec, "dhcp client [Wait Reply enter execs]")
                  /** blocking after enter executives of unforced state. **/
                  FSM EXIT (3,"dhcp client")
                  /** state (Wait Reply) exit executives **/
                  FSM STATE EXIT UNFORCED (1, "Wait Reply", "dhcp client
[Wait Reply exit execs]")
                        FSM PROFILE SECTION IN ("dhcp client [Wait Reply exit
execs]", state1 exit exec)
                        intrpt type = op intrpt type();
                        if (intrpt type == OPC INTRPT STRM)
                              dhcp get packet();
                              if(server\ id == 0)
                                    /* If Mobile nodes have not yet
discovered a server, remember this new server's identification:*/
                                    dhcp opt ptr =
dhcp optlist get opt(DHCPC OPT SERVERID, rcv pkt opts);
                                    server id = dhcp opt ptr->simple data;
                                    op ima obj hname get(server id, srv str,
200);
      /* Turn off rapid commit if the server didn't indicate support: */
      if(!(dhcp optlist get opt(DHCPC OPT RAPID COMMIT, rcv pkt opts)))
                                          server rapid commit = OPC FALSE;
                                    }
                        else
                              intrpt code = op intrpt code();
                        FSM PROFILE SECTION OUT (state1 exit exec)
                  /** state (Wait Reply) transition processing **/
                  FSM PROFILE SECTION IN ("dhcp client [Wait_Reply trans
conditions]", state1 trans conds)
                  FSM INIT COND (Rcv Reply)
                  FSM TEST COND (Rcv NonReply & !Rcv Advertise)
                  FSM TEST COND (Msg Tmout)
                  FSM TEST_COND (Msg_Fail)
                  FSM TEST COND (Rcv Advertise)
                  FSM TEST LOGIC ("Wait Reply")
                  FSM PROFILE SECTION OUT (state1 trans conds)
                  FSM TRANSIT SWITCH
                        FSM CASE TRANSIT (0, 3, state3 enter exec, Config();,
"Rcv Reply", "Config()", "Wait Reply", "Idle", "tr 21", "dhcp client
[Wait Reply -> Idle : Rcv Reply / Config()]")
```

```
FSM_CASE_TRANSIT (1, 1, state1_enter_exec,
Discard();, "Rcv NonReply & !Rcv Advertise", "Discard()", "Wait Reply",
"Wait Reply", "tr 24", "dhcp client [Wait Reply -> Wait Reply : Rcv NonReply
& !Rcv Advertise / Discard()]")
                        FSM CASE TRANSIT (2, 1, state1 enter exec,
SendMsg();, "Msg Tmout", "SendMsg()", "Wait Reply", "Wait Reply", "tr 26",
"dhcp client [Wait Reply -> Wait Reply : Msg Tmout / SendMsg()]")
                        FSM CASE TRANSIT (3, 0, state0 enter exec, ;,
"Msg Fail", "", "Wait Reply", "Begin", "tr 46", "dhcp client [Wait Reply ->
Begin : Msg Fail / ]")
FSM_CASE_TRANSIT (4, 1, state1_enter_exec, SendMsg();, "Rcv_Advertise", "SendMsg()", "Wait_Reply", "Wait_Reply",
"tr 51", "dhcp client [Wait Reply -> Wait Reply : Rcv Advertise /
SendMsq()]")
                  /** state (Wait Advertise) enter executives **/
                  FSM STATE ENTER UNFORCED (2, "Wait Advertise",
state2 enter exec, "dhcp client [Wait Advertise enter execs]")
                   /** blocking after enter executives of unforced state. **/
                  FSM EXIT (5, "dhcp client")
                  /** state (Wait_Advertise) exit executives **/
                  FSM STATE EXIT UNFORCED (2, "Wait Advertise", "dhcp client
[Wait Advertise exit execs]")
                         FSM PROFILE SECTION IN ("dhcp client [Wait Advertise
exit execs]", state2 exit exec)
                         intrpt type = op intrpt type();
                         if (intrpt type == OPC INTRPT STRM)
                               dhcp get packet();
                               if(server\ id == 0)
      /* If the mobile nodes have not yet discovered a server, remember this
new server's identification:*/
                                     dhcp opt ptr =
dhcp optlist get opt(DHCPC OPT SERVERID, rcv pkt opts);
                                     server id = dhcp opt ptr->simple data;
                                     op ima obj hname get (server id, srv str,
200);
      /* Turn off rapid commit if the server didn't indicate support: */
      if(!(dhcp optlist get opt(DHCPC OPT RAPID COMMIT, rcv pkt opts)))
                                           server rapid commit = OPC FALSE;
                         else
                               intrpt code = op intrpt code();
                         FSM PROFILE SECTION OUT (state2 exit exec)
```

```
/** state (Wait Advertise) transition processing **/
                 FSM PROFILE SECTION IN ("dhcp client [Wait Advertise trans
conditions]", state2 trans conds)
                 FSM INIT COND (Rcv Advertise)
                 FSM TEST COND (Rcv NonAdvertise)
                 FSM TEST COND (Msg Tmout)
                 FSM TEST LOGIC ("Wait Advertise")
                 FSM PROFILE SECTION OUT (state2 trans conds)
                 FSM TRANSIT SWITCH
                       FSM CASE TRANSIT (0, 1, state1 enter exec,
SendMsg();, "Rcv Advertise", "SendMsg()", "Wait Advertise", "Wait Reply",
"tr 16", "dhcp client [Wait Advertise -> Wait Reply : Rcv Advertise /
SendMsq() ] ")
                       FSM CASE TRANSIT (1, 2, state2 enter exec,
Discard();, "Rcv NonAdvertise", "Discard()", "Wait Advertise",
"Wait Advertise", "tr 18", "dhcp client [Wait Advertise -> Wait Advertise :
Rcv NonAdvertise / Discard()]")
                       FSM CASE TRANSIT (2, 2, state2 enter exec,
SendMsg();, "Msg Tmout", "SendMsg()", "Wait Advertise", "Wait Advertise",
"tr 19", "dhcp client [Wait Advertise -> Wait Advertise : Msg Tmout /
SendMsg()]")
      /*----*/
                 /** state (Idle) enter executives **/
                 FSM STATE ENTER UNFORCED (3, "Idle", state3 enter exec,
"dhcp client [Idle enter execs]")
                 /** blocking after enter executives of unforced state. **/
                 FSM EXIT (7,"dhcp client")
                  /** state (Idle) exit executives **/
                 FSM STATE EXIT UNFORCED (3, "Idle", "dhcp client [Idle exit
execs]")
                       FSM PROFILE SECTION IN ("dhcp client [Idle exit
execs]", state3 exit exec)
                       intrpt type = op intrpt type();
                       if (intrpt type == OPC INTRPT STRM)
                             dhcp get packet();
                       else
                             intrpt_code = op intrpt code();
                       FSM PROFILE SECTION OUT (state3 exit exec)
                  /** state (Idle) transition processing **/
                 FSM PROFILE SECTION IN ("dhcp client [Idle trans
conditions]", state3 trans conds)
                 FSM INIT COND (Msg Rcv)
                 FSM TEST COND (Assign Tmout)
                 FSM TEST COND (Msg Fail)
```

```
FSM DFLT COND
                  FSM TEST LOGIC ("Idle")
                  FSM PROFILE SECTION OUT (state3 trans conds)
                  FSM TRANSIT SWITCH
                        FSM_CASE_TRANSIT (0, 3, state3_enter_exec,
Discard();, "Msg Rcv", "Discard()", "Idle", "Idle", "tr 27", "dhcp client
[Idle -> Idle : Msg Rcv / Discard()]")
FSM_CASE_TRANSIT (1, 4, state4_enter_exec, SendMsg();, "Assign_Tmout", "SendMsg()", "Idle", "Wait_Reply_Conf", "tr_28",
"dhcp_client [Idle -> Wait_Reply_Conf : Assign_Tmout / SendMsg()]")
                        FSM_CASE_TRANSIT (2, 0, state0_enter_exec, ;,
"Msg Fail", "", "Idle", "Begin", "tr 55", "dhcp client [Idle -> Begin :
Msg Fail / ]")
                       FSM CASE TRANSIT (3, 3, state3 enter exec, ;,
"default", "", "Idle", "Idle", "tr 56", "dhcp client [Idle -> Idle : default
/ 1")
                     _____*/
                  /** state (Wait Reply Conf) enter executives **/
                  FSM STATE ENTER UNFORCED (4, "Wait Reply Conf",
state4 enter exec, "dhcp client [Wait Reply Conf enter execs]")
                  /** blocking after enter executives of unforced state. **/
                  FSM EXIT (9, "dhcp client")
                  /** state (Wait Reply Conf) exit executives **/
                  FSM STATE EXIT UNFORCED (4, "Wait Reply Conf", "dhcp client
[Wait Reply Conf exit execs]")
                        FSM PROFILE SECTION IN ("dhcp client [Wait Reply Conf
exit execs]", state4 exit exec)
                        intrpt type = op intrpt type();
                        if (intrpt type == OPC INTRPT STRM)
                              dhcp get packet();
                              /\star If mobile nodes have not yet discovered a
server, or if we are in the Rebind phase, remember this new server's
identification:*/
                              if((server_id == 0) || (snd_msg_type ==
DHCPC MSG REBIND))
                                    dhcp opt ptr =
dhcp optlist get_opt(DHCPC_OPT_SERVERID, rcv_pkt_opts);
                                    server id = dhcp opt ptr->simple data;
                                    op ima obj hname get (server id, srv str,
200);
      /* Turn off rapid commit if the server didn't indicate support: */
      if(!(dhcp optlist get opt(DHCPC OPT RAPID COMMIT, rcv pkt opts)))
                                          server rapid commit = OPC FALSE;
                                    }
```

```
else
                              intrpt code = op intrpt code();
                        FSM PROFILE SECTION OUT (state4 exit exec)
                  /** state (Wait Reply Conf) transition processing **/
                  FSM PROFILE SECTION IN ("dhcp client [Wait Reply Conf trans
conditions]", state4 trans conds)
                  FSM INIT COND (Rcv Reply)
                  FSM_TEST_COND (Rcv_NonReply)
                  FSM TEST COND (Msg Tmout)
                  FSM TEST COND (Msg Fail)
                  FSM TEST LOGIC ("Wait Reply Conf")
                  FSM PROFILE SECTION OUT (state4 trans conds)
                  FSM TRANSIT_SWITCH
                        FSM CASE TRANSIT (0, 3, state3 enter exec, Config();,
"Rcv Reply", "Config()", "Wait Reply Conf", "Idle", "tr 29", "dhcp client
[Wait Reply Conf -> Idle : Rcv Reply / Config()]")
                        FSM CASE TRANSIT (1, 4, state4 enter exec,
Discard();, "Rcv_NonReply", "Discard()", "Wait_Reply_Conf",
"Wait_Reply_Conf", "tr_32", "dhcp_client [Wait_Reply_Conf -> Wait_Reply_Conf
: Rcv NonReply / Discard()]")
                        FSM CASE TRANSIT (2, 4, state4 enter exec,
SendMsg();, "Msg_Tmout", "SendMsg()", "Wait_Reply Conf", "Wait Reply Conf",
"tr 34", "dhcp client [Wait Reply Conf -> Wait Reply Conf : Msg Tmout /
SendMsg()]")
                       FSM CASE TRANSIT (3, 0, state0 enter exec,
Unconfig();, "Msg Fail", "Unconfig()", "Wait Reply Conf", "Begin", "tr 50",
"dhcp_client [Wait_Reply_Conf -> Begin : Msg_Fail / Unconfig()]")
                            -----*/
                  /** state (Init) enter executives **/
                  FSM STATE ENTER UNFORCED NOLABEL (5, "Init", "dhcp client
[Init enter execs]")
                        FSM PROFILE SECTION IN ("dhcp client [Init enter
execs]", state5 enter exec)
                        if (LTRACE ACTIVE)
                              op prg odb print major ("DHCP Client: Begin
simulation" , OPC NIL);
                        /* Initialize some variables: */
                        module objid
                                                      = op id self();
                        node objid
op topo parent (module objid);
                        ip support module ptr = ip support module data get
(node objid);
                        gateway node
ip rte node is gateway(ip support module ptr);
                        last trans id
                                                      = 1;
                        server inet addr = INETC ADDRESS INVALID;
                        prg list init(&interfaces config);
```

```
/* Determine if the nodes are logging DHCP: */
            op ima sim attr get(OPC IMA TOGGLE, "DHCP Logging", &logging);
                        if (LOGGING ACTIVE)
                              new log handle =
op prg log handle create(OpC Log Category Protocol, "DHCP", "New
Configuration", 100);
                              renew log handle =
op prg log handle create(OpC Log Category Protocol, "DHCP", "Renewal", 100);
                              error log handle =
op prg log handle create(OpC Log Category Protocol, "DHCP", "Protocol Error",
10\overline{0});
                              expire log handle =
op prg log handle create(OpC Log Category Protocol, "DHCP", "Expiration",
100);
                              }
                        /* Register process in the Process Registry: */
            op ima obj attr get (module objid, "process model", tmp str);
      my proc handle = oms pr process register(node objid, module objid,
                              op pro self(), tmp str);
oms pr attr set (my proc handle,
                              "protocol", OMSC PR STRING, "dhcp",
                              OPC NIL);
                  /* Schedule a self interrupt to wait for lower layers */
                        op intrpt schedule self(op sim time(), 0);
                        FSM PROFILE SECTION OUT (state5 enter exec)
                  /** blocking after enter executives of unforced state. **/
                  FSM EXIT (11,"dhcp client")
                  /** state (Init) exit executives **/
                  FSM STATE EXIT UNFORCED (5, "Init", "dhcp client [Init exit
execs]")
                  /** state (Init) transition processing **/
                  FSM TRANSIT FORCE (8, state8 enter exec, ;, "default", "",
"Init", "Wait_0", "tr_19_0", "dhcp_client [Init -> Wait 0 : default / ]")
                  /** state (Wait 1) enter executives **/
                  FSM STATE ENTER UNFORCED (6, "Wait 1", state6 enter exec,
"dhcp client [Wait_1 enter execs]")
                        FSM PROFILE SECTION IN ("dhcp client [Wait 1 enter
execs]", state6 enter exec)
                        op intrpt schedule self(op sim time(), 0);
                        FSM PROFILE SECTION OUT (state6 enter exec)
                  /** blocking after enter executives of unforced state. **/
                  FSM_EXIT (13,"dhcp_client")
```

```
/** state (Wait 1) exit executives **/
                 FSM STATE EXIT UNFORCED (6, "Wait 1", "dhcp client [Wait 1
exit execs]")
                 /** state (Wait 1) transition processing **/
                 FSM TRANSIT FORCE (9, state9 enter exec, ;, "default", "",
"Wait 1", "Wait 2", "tr 19 4", "dhcp client [Wait 1 -> Wait 2 : default / ]")
           /*----*/
                 /** state (Wait_3) enter executives **/
                 FSM STATE ENTER UNFORCED (7, "Wait 3", state7 enter exec,
"dhcp client [Wait_3 enter execs]")
                       FSM PROFILE SECTION IN ("dhcp client [Wait 3 enter
execs]", state7 enter exec)
                       op intrpt schedule self(op sim time(), 0);
                       FSM PROFILE SECTION OUT (state7 enter exec)
                 /** blocking after enter executives of unforced state. **/
                 FSM EXIT (15, "dhcp client")
                 /** state (Wait 3) exit executives **/
                 FSM STATE EXIT UNFORCED (7, "Wait 3", "dhcp client [Wait 3
exit execs]")
                       FSM PROFILE SECTION IN ("dhcp client [Wait 3 exit
execs]", state7 exit exec)
                       /* Read in all of our attribute values: */
                       op ima obj attr get (module objid, "DHCPv6 Client
Parameters", &cli params_objid);
                       objid1 = op topo child(cli params objid,
OPC OBJTYPE GENERIC, 0);
                       op ima obj attr get str(objid1, "Sending Interface",
TMP STR SIZE, val_str);
                       send iface =
dhcp get ipv6 table index from name(val str, ip support module ptr);
                       if(send iface == -1)
                             sprintf(tmp str, "Unable to get interface index
for interface '%s'", val str);
                             op sim end ("Error resolving interface index:",
                                   tmp str, OPC NIL, OPC NIL);
                       else
                 /* Remember the link layer address of this interface: */
                             IpT Interface Info*
                                                   ip iface elem ptr =
inet_rte_intf_tbl_access(ip_support_module_ptr, send_iface);
                             node ll addr =
ip rte intf link local addr get(ip iface elem ptr);
```

```
&rapid commit);
                        op ima obj attr get(objid1, "Timers", &timers objid);
                        timers row objid = op topo child(timers objid,
OPC OBJTYPE GENERIC, 0);
                        op ima obj attr get int32(timers row objid, "Max
Solicit Timeout", &max_sol_tmout);
                        op ima obj attr get int32(timers row objid, "Max
Request Timeout", &max req tmout);
                        op ima obj attr get int32(timers row objid, "Max
Request Retries", &max req retries);
                        op ima obj attr get int32(timers row objid, "Initial
Renew Timeout", &init_renew tmout);
                        op ima obj attr get int32(timers row objid, "Max
Renew Timeout", &max renew tmout);
                        op ima obj attr get int32(timers row objid, "Initial
Rebind Timeout", &init rebind tmout);
                        op ima obj attr get int32(timers row objid, "Max
Rebind Timeout", &max rebind tmout);
                        /* Create the ICI to be used with UDP: */
                        command ici ptr = op ici create ("udp command inet");
                        /* Create a receive port for this application: */
                        tmp int = dhcp connect to udp(DHCPC PORT CLIENT,
module objid,
                              &udp objid, &input strm, &output strm,
command ici ptr);
                        if (tmp int != UDPC IND SUCCESS)
                              sprintf (tmp str, "%d in response to
CREATE PORT command", tmp int);
                              op sim end ("Error: process model dhcp received
error", tmp str, "", "");
            /* Me: Since we are a client, mobile node will always send to the
well known multicast address for DHCP servers. Set this address on the ICI
here: */
                        addr fam = InetC Addr Family v6;
                        server multicast addr =
inet address create(DHCPC ADDR ALL AGENTS AND SERVERS, addr fam);
                        op ici attr set ptr(command ici ptr, "rem addr",
&server multicast addr);
            /* Me: For multicast to work, mobile node need a setting into the
"strm index" field of the UDP ici - this setting will eventually be written
into the "multicast major port" of the IP layer ICI.*/
                        op ici attr set (command ici ptr, "strm index",
IPC MCAST ALL MAJOR PORTS);
                        server multicast addr =
inet_address_create(DHCPC_ADDR_ALL_AGENTS_AND_SERVERS, addr_fam);
```

op ima obj attr get (objid1, "Rapid Commit",

```
/* Mobile node always send to the DHCP server port: */
                        op ici attr set (command ici ptr, "rem port",
DHCPC PORT SERVER);
                        /* Also, since we are a client, mobile node will
always send using our link local address: */
                        op ici attr set (command ici ptr, "src addr",
&node ll addr);
/* Build the list to contain configuration information for all existing
interfaces. For both routers and hosts.*/
                        if (gateway node)
                              int
                                          rep;
                                          num total interfaces = 0;
                              Objid ip group objid, iface info objid,
iface row objid;
                              Objid gaddr info objid, gaddr row objid;
                              DhcpT Cli Assignment * iface config;
                              /* First get the total number of interfaces on
this node. We're only dealing with physical interfaces for now.*/
                              op ima obj attr get(node objid, "IPv6
Parameters", &ip group objid);
                              objid1 = op topo child(ip group objid,
OPC OBJTYPE GENERIC, 0);
                              op_ima_obj_attr_get(objid1, "Interface
Information", &iface info objid);
                              num total interfaces =
op topo child count(iface info objid, OPC OBJTYPE GENERIC);
                              for (rep = 0; rep < num total interfaces;</pre>
rep++)
                                     {
                                     iface row objid =
op topo child(iface info objid, OPC OBJTYPE GENERIC, rep);
                                     op ima obj attr get(iface row objid,
"Global Address(es)", &gaddr info objid);
                                     gaddr row objid =
op topo child(gaddr info objid, OPC OBJTYPE GENERIC, 0);
                                     /* Get the "Address": */
                                     op_ima_obj_attr_get_str(gaddr_row_objid,
"Address", 100, val str);
                                     if(strstr(val str, "DHCP"))
                                           int intf table index;
                                           char iface str [8];
      /\star First confirm that the interface has been added into the interface
table */
                                           op ima obj attr get str
(iface row objid, "Name", 7, iface str);
```

```
if ((intf table index =
dhcp get ipv6 table index from name (iface str, ip support module ptr)) == -
                                                 char log str [500];
                  printf (log str, "Interface %s has DHCP specified, but is
not being used in the simulation. Confirm that it is connected, has a link-
local address, and no other global addresses are specified.", iface str);
                                                LOG(error log handle,
log str, PrgC Log Severity Error);
                                                 continue;
                                          else
            /* This interface is using DHCP for either address or prefix
assignment. */
                                                 iface config =
(DhcpT Cli Assignment *) op prg mem alloc(sizeof(DhcpT Cli Assignment));
                                                 iface config->configured =
OPC FALSE;
                                                 iface config->iface index =
rep;
                                                 if(strstr(val str, "Prefix
Delegation"))
                                                       iface config-
>assign type = DHCPC ASSIGN PREFIX;
                                                 else
                                                       iface config-
>assign type = DHCPC ASSIGN ADDR;
      op ima obj attr get str(iface row objid, "Name", TMP STR SIZE,
tmp str);
                                                 iface config->iface name =
prg string copy(tmp str);
      prg list insert(&interfaces config, iface config, PRGC LISTPOS TAIL);
                        else
      /* Initialize the configuration data for only a single interface: */
                              DhcpT Cli Assignment * iface config =
                                    (DhcpT Cli Assignment
*)op prg mem alloc(sizeof(DhcpT Cli Assignment));
                              iface config->configured = OPC FALSE;
                              iface config->iface index = send iface;
                              iface config->assign type = DHCPC ASSIGN ADDR;
                              iface config->iface name = prg string copy("N/A
(single interface host)");
```

```
prg list insert (&interfaces config,
iface config, PRGC LISTPOS TAIL);
                      num interfaces = prg list size(&interfaces config);
                      if(num interfaces < 1)</pre>
                            LOG(error log handle, "No interfaces are
requesting configuration information. "
                                 "Exiting DHCP client.",
PrgC Log Severity Error);
                            op pro destroy(op pro self());
                      /* Get the handles to the DHCP statistics: */
                      global stats = dhcp get global stathandles();
                      local stats =
dhcp get local stathandles(DHCPC CLIENT);
                      FSM PROFILE SECTION OUT (state7 exit exec)
                 /* Me: Register this address for multicast: */
     Inet Address Multicast Register(server multicast addr, tmp int,
     DHCPC PORT SERVER, ip support module ptr);
                /** state (Wait 3) transition processing **/
                FSM TRANSIT FORCE (0, state0 enter exec, ;, "default", "",
/** state (Wait 0) enter executives **/
                FSM_STATE_ENTER_UNFORCED (8, "Wait_0", state8_enter_exec,
"dhcp client [Wait 0 enter execs]")
                      FSM PROFILE SECTION IN ("dhcp client [Wait 0 enter
execs]", state8 enter exec)
                      op intrpt schedule self(op sim time(), 0);
                      FSM PROFILE SECTION OUT (state8 enter exec)
                 /** blocking after enter executives of unforced state. **/
                FSM EXIT (17, "dhcp client")
                 /** state (Wait 0) exit executives **/
                FSM STATE EXIT UNFORCED (8, "Wait 0", "dhcp client [Wait 0
exit execs]")
                /** state (Wait 0) transition processing **/
                FSM TRANSIT FORCE (6, state6 enter exec, ;, "default", "",
"Wait_0", "Wait_1", "tr_19_3", "dhcp_client [Wait_0 -> Wait_1 : default / ]")
     /*----*/
```

```
/** state (Wait 2) enter executives **/
                  FSM STATE ENTER UNFORCED (9, "Wait 2", state9 enter exec,
"dhcp client [Wait 2 enter execs]")
                        FSM PROFILE SECTION IN ("dhcp client [Wait 2 enter
execs]", state9 enter exec)
                        op intrpt schedule self(op sim time(), 0);
                        FSM PROFILE SECTION OUT (state9 enter exec)
                  /** blocking after enter executives of unforced state. **/
                  FSM EXIT (19, "dhcp client")
                  /** state (Wait 2) exit executives **/
                  FSM STATE EXIT UNFORCED (9, "Wait 2", "dhcp client [Wait 2
exit execs]")
                  /** state (Wait 2) transition processing **/
                  FSM TRANSIT FORCE (7, state7 enter exec, ;, "default", "",
"Wait_2", "Wait_3", "tr_19_1", "dhcp_client [Wait_2 -> Wait_3 : default / ]")
            FSM EXIT (5, "dhcp client")
      }
/* Undefine shortcuts to state variables to avoid */
#undef module objid
#undef node objid
#undef my proc handle
#undef udp objid
#undef ip support module ptr
#undef command_ici_ptr
#undef max sol tmout
#undef max req tmout
#undef max req retries
#undef init renew tmout
#undef max renew tmout
#undef init rebind tmout
#undef max rebind tmout
#undef send iface
#undef last trans id
#undef rapid commit
#undef server_rapid_commit
#undef server id
#undef srv str
#undef interfaces config
#undef gateway node
#undef input strm
#undef output_strm
#undef snd pkt ptr
#undef snd pkt opts
#undef snd msg type
#undef rcv_pkt_ptr
#undef rcv pkt opts
#undef rcv_msg_type
```

```
#undef rcv msg trans
#undef RTprev
#undef expire time
#undef rebind time
#undef retrans count
#undef intrpt type
#undef intrpt code
#undef num interfaces
#undef next evh
#undef server multicast addr
#undef node ll addr
#undef global stats
#undef local stats
#undef transaction time
#undef logging
#undef new log handle
#undef renew log_handle
#undef error log handle
#undef expire log handle
#undef server inet addr
#undef FIN PREAMBLE DEC
#undef FIN PREAMBLE CODE
#define FIN PREAMBLE DEC
#define FIN PREAMBLE CODE
VosT Obtype
op dhcp client init (int * init block ptr)
      VosT Obtype obtype = OPC NIL;
      FIN MT ( op dhcp client init (init block ptr))
      obtype = Vos_Define_Object_Prstate ("proc state vars (dhcp_client)",
            sizeof (dhcp_client_state));
      *init block ptr = 10;
      FRET (obtype)
      }
VosT Address
_op_dhcp_client_alloc (VosT_Obtype obtype, int init block)
#if !defined (VOSD NO FIN)
      int op block origin = 0;
#endif
      dhcp client state * ptr;
      FIN MT ( op dhcp client alloc (obtype))
      ptr = (dhcp client state *)Vos Alloc Object (obtype);
      if (ptr != OPC NIL)
            ptr-> op current block = init block;
#if defined (OPD \overline{\text{ALLOW}} ODB)
            ptr-> op current state = "dhcp client [Init enter execs]";
#endif
```

```
FRET ((VosT Address)ptr)
void
op dhcp client svar (void * gen ptr, const char * var name, void **
var_p_ptr)
      dhcp_client state
                              *prs ptr;
      FIN MT ( op dhcp_client_svar (gen_ptr, var_name, var_p_ptr))
      if (var name == OPC NIL)
            *var p ptr = (void *)OPC NIL;
      prs ptr = (dhcp client state *)gen ptr;
      if (strcmp ("module objid" , var name) == 0)
            *var_p_ptr = (void *) (&prs_ptr->module objid);
            FOUT
      if (strcmp ("node objid" , var name) == 0)
            *var p ptr = (void *) (&prs ptr->node objid);
            FOUT
      if (strcmp ("my proc handle" , var name) == 0)
            *var p_ptr = (void *) (&prs_ptr->my_proc_handle);
            FOUT
      if (strcmp ("udp_objid" , var_name) == 0)
            *var p ptr = (void *) (&prs ptr->udp objid);
            FOUT
      if (strcmp ("ip support module ptr" , var name) == 0)
            *var_p_ptr = (void *) (&prs_ptr->ip_support_module_ptr);
            FOUT
      if (strcmp ("command_ici_ptr" , var_name) == 0)
            *var_p_ptr = (void *) (&prs_ptr->command_ici_ptr);
            FOUT
      if (strcmp ("max sol tmout" , var name) == 0)
            *var_p_ptr = (void *) (&prs_ptr->max_sol_tmout);
            FOUT
      if (strcmp ("max_req_tmout" , var_name) == 0)
            *var p ptr = (void *) (&prs ptr->max req tmout);
            FOUT
```

```
if (strcmp ("max req retries" , var name) == 0)
      *var p ptr = (void *) (&prs ptr->max req retries);
      FOUT
if (strcmp ("init renew tmout" , var name) == 0)
      *var p ptr = (void *) (&prs ptr->init renew tmout);
      FOUT
if (strcmp ("max renew tmout" , var name) == 0)
      *var p ptr = (void *) (&prs ptr->max renew tmout);
      FOUT
if (strcmp ("init rebind tmout" , var name) == 0)
      *var p ptr = (void *) (&prs ptr->init rebind tmout);
      FOUT
if (strcmp ("max rebind tmout" , var name) == 0)
      *var_p_ptr = (void *) (&prs_ptr->max_rebind tmout);
      FOUT
if (strcmp ("send iface" , var name) == 0)
      *var p ptr = (void *) (&prs ptr->send iface);
      FOUT
if (strcmp ("last trans id" , var name) == 0)
      *var_p_ptr = (void *) (&prs_ptr->last trans id);
      FOUT
if (strcmp ("rapid commit" , var name) == 0)
      *var p ptr = (void *) (&prs ptr->rapid commit);
      FOUT
if (strcmp ("server rapid commit" , var name) == 0)
      *var_p_ptr = (void *) (&prs_ptr->server_rapid_commit);
      FOUT
if (strcmp ("server id" , var name) == 0)
      *var p ptr = (void *) (&prs ptr->server id);
      FOUT
if (strcmp ("srv str" , var name) == 0)
      *var_p_ptr = (void *) (prs_ptr->srv str);
      FOUT
if (strcmp ("interfaces_config" , var_name) == 0)
```

```
*var p ptr = (void *) (&prs ptr->interfaces config);
      FOUT
if (strcmp ("gateway node" , var name) == 0)
      *var_p_ptr = (void *) (&prs_ptr->gateway_node);
      FOUT
if (strcmp ("input_strm" , var_name) == 0)
      *var_p_ptr = (void *) (&prs_ptr->input_strm);
      FOUT
if (strcmp ("output_strm" , var_name) == 0)
      *var_p_ptr = (void *) (&prs_ptr->output_strm);
      FOUT
if (strcmp ("snd_pkt_ptr" , var_name) == 0)
      *var_p_ptr = (void *) (&prs_ptr->snd_pkt_ptr);
      FOUT
      }
if (strcmp ("snd pkt opts" , var name) == 0)
      *var_p_ptr = (void *) (&prs_ptr->snd_pkt_opts);
      FOUT
if (strcmp ("snd_msg_type" , var_name) == 0)
      *var p ptr = (void *) (&prs ptr->snd msg type);
      FOUT
if (strcmp ("rcv_pkt_ptr" , var_name) == 0)
      *var p_ptr = (void *) (&prs_ptr->rcv_pkt_ptr);
      FOUT
if (strcmp ("rcv_pkt_opts" , var_name) == 0)
      *var_p_ptr = (void *) (&prs_ptr->rcv_pkt_opts);
      FOUT
if (strcmp ("rcv_msg_type" , var_name) == 0)
      *var_p_ptr = (void *) (&prs_ptr->rcv_msg_type);
      FOUT
if (strcmp ("rcv_msg_trans" , var_name) == 0)
      *var_p_ptr = (void *) (&prs_ptr->rcv_msg_trans);
      FOUT
if (strcmp ("RTprev" , var_name) == 0)
      *var_p_ptr = (void *) (&prs_ptr->RTprev);
```

```
FOUT
if (strcmp ("expire time" , var name) == 0)
      *var p ptr = (void *) (&prs ptr->expire time);
      FOUT
if (strcmp ("rebind time" , var name) == 0)
      *var p ptr = (void *) (&prs ptr->rebind time);
      FOUT
if (strcmp ("retrans count" , var name) == 0)
      *var p ptr = (void *) (&prs ptr->retrans count);
      FOUT
if (strcmp ("intrpt_type" , var_name) == 0)
      *var p ptr = (void *) (&prs ptr->intrpt type);
      FOUT
if (strcmp ("intrpt code" , var name) == 0)
      *var p ptr = (void *) (&prs ptr->intrpt code);
      FOUT
if (strcmp ("num interfaces" , var name) == 0)
      *var p ptr = (void *) (&prs ptr->num interfaces);
      FOUT
if (strcmp ("next evh" , var name) == 0)
      *var p ptr = (void *) (&prs ptr->next evh);
      FOUT
if (strcmp ("server multicast addr" , var name) == 0)
      *var p ptr = (void *) (&prs ptr->server multicast addr);
      FOUT
if (strcmp ("node ll addr" , var name) == 0)
      *var_p_ptr = (void *) (&prs_ptr->node ll addr);
      FOUT
      }
if (strcmp ("global stats" , var name) == 0)
      *var p ptr = (void *) (&prs ptr->global stats);
if (strcmp ("local stats" , var name) == 0)
      *var p ptr = (void *) (&prs ptr->local stats);
      FOUT
```

```
if (strcmp ("transaction time" , var name) == 0)
            *var p ptr = (void *) (&prs ptr->transaction time);
            FOUT
      if (strcmp ("logging", var name) == 0)
            *var p ptr = (void *) (&prs ptr->logging);
            FOUT
      if (strcmp ("new log handle" , var name) == 0)
            *var p ptr = (void *) (&prs ptr->new log handle);
            FOUT
      if (strcmp ("renew log handle" , var name) == 0)
            *var p ptr = (void *) (&prs ptr->renew log handle);
            FOUT
      if (strcmp ("error log handle" , var name) == 0)
            *var_p_ptr = (void *) (&prs_ptr->error log handle);
            FOUT
      if (strcmp ("expire log handle" , var name) == 0)
            *var_p_ptr = (void *) (&prs ptr->expire log handle);
            FOUT
      if (strcmp ("server inet addr" , var name) == 0)
            *var p ptr = (void *) (&prs ptr->server inet addr);
            FOUT
      *var p ptr = (void *)OPC NIL;
      FOUT
      }
static void /* Me: Joinning multicast tree by using CoA address*/
ip igmp host join grp (IpT Address ip grp addr, int interface)
                                    grp elem ptr;
      IpT Igmp Host Grp Elem*
      List*
                                           ip_grp_intf_list_ptr;
     Packet*
                                                 igmp msg pkptr;
     Packet*
                                                 ip dgram pkptr;
                                           msg0 [256], msg1 [256], msg2 [256];
      char
      char
                                           ip addr str [IPC ADDR STR LEN];
      FIN (ip_igmp_host_join_grp (ip_grp_addr, interface));
      /* Generate a trace message
      if (LTRACE IGMP)
            ip address print (ip addr str, ip grp addr);
```

```
sprintf (msg0, "IP Group Address :
                                                             %s",
ip addr str);
           sprintf (msgl, "IP Interface
                                                                %d",
interface);
           op prg odb print major ("Received a Join request from an
application for: ", msg0, msg1, OPC NIL);
      /* Check if the group membership information for the given IP group
address already exists */
      grp_elem_ptr = ip_igmp_host_get_grp_elem (ip_grp_addr, interface);
      if (grp elem ptr == OPC NIL)
           grp elem ptr = ip igmp host grp elem alloc ();
           /* Set the fields of the data structure */
           grp elem ptr->ip grp addr = ip address copy (ip grp addr);
           grp elem ptr->interface = interface;
           /* Until now only one application on this node has joined this
group */
           grp elem ptr->num of apps = 1;
           /* Assign a timer ID to this group's delay timer */
           grp elem ptr->delay timer id = timer id++;
            /* Initialize this field to the value of Robustness Variable
attribute minus 1 */
           grp elem ptr->unsolicit msg count = robustness variable - 1;
           /* Since we are going to send an Unsolicited Membership Report
message, set this field to OPC TRUE*/
           grp elem ptr->report sent flag = OPC TRUE;
           /* Start the delay timer to send the next Unsolicited Membership
Report message. */
           if (igmp sim efficiency == OPC FALSE)
                  /* Set the delay timer to OPC TRUE, as we are going to
start */
                 /* the timer for sending the next Unsolicited Report
message */
                 grp elem ptr->timer on flag = OPC TRUE;
                  /* Start the delay timer
                 grp elem ptr->delay timer evh = op intrpt schedule self
(op sim time () +
                             unsolicited report interval, grp elem ptr-
>delay_timer_id);
           else
                  /* Set the unsolicited messages count to zero */
                 grp_elem_ptr->unsolicit msg count = 0;
```

```
}
            /** Add this group membership information to the corresponding
list
            /* Generate a trace message
            if (LTRACE IGMP)
                  ip address print (ip addr str, ip grp addr);
                  sprintf (msq0, "IP Group Address
                                                                         %s",
ip addr str);
                  sprintf (msql, "IP Interface
                                                                         %d",
interface);
                  strcpy (msq2, "to the table, which is maintained by IGMP
Host process.");
                  op prg odb print major ("Adding group membership
information for: ", msg0, msg1, msg2, OPC NIL);
            /* Get the list for the interface on which this node has joined
the group */
            ip grp intf list ptr = (List *) op prg list access
(ip grp list ptr, interface);
            /* Add the group membership information data structure to the
list
            op_prg_list_insert (ip grp intf list ptr, grp elem ptr,
OPC LISTPOS TAIL);
            /** Create and send an Unsolicited Membership Report message **/
            /* Generate a trace message
            if (LTRACE IGMP)
                  ip address print (ip addr str, ip grp addr);
                  sprintf (msg0, "IP Group Address
                                                                         %s",
ip addr str);
                  sprintf (msq1, "IP Interface
                                                                         %d",
interface);
                  op prg odb print major ("Sending an IGMP Unsolicited
Membership Report message for: ", msg0, msg1, OPC NIL);
                  op prg odb print major ("Starting Delay timer for the above
group.", OPC NIL);
            /* Create an IGMP Report message */
            igmp_msg_pkptr = ip_igmp_create_igmp_msg
(IpC Igmp Membership Report Msg, 0, ip grp addr);
            /* Update IGMP messages sent statistics */
            ip igmp host igmp_msgs_sent_stat_update (op_pk_total_size_get
(igmp msg pkptr));
            /* Create an IP datagram for transmitting the IGMP message */
            ip_dgram_pkptr = ip_igmp_create_ipdgram (igmp msg pkptr,
ip grp addr);
```

```
/* Schedule a procedure interrupt for the current simulation time
to invoke the IP process */
            op intrpt schedule call (op sim time (), interface,
ip igmp host ip process invoke, ip_dgram_pkptr);
      else
            /* Group membership information already exists. Another
application on this node has joined */
      /* this group. Increment the num of apps field's value by one */
            grp elem ptr->num of apps++;
      FOUT;
      }
static void /* Using IGMP protocol to leave multicast group*/
ip igmp host leave grp (IpT Address ip grp addr, int interface)
      IpT Igmp Host Grp Elem*
                                    grp elem ptr;
                                    ip grp intf list ptr;
   List*
                                    num grp elems, i;
      int
      Packet*
                                    igmp msg pkptr;
                                    ip dgram pkptr;
      Packet*
                                    ip addr str [IPC ADDR STR LEN];
      char
                                    msg0 [256], msg1 [256], msg2 [256];
      char
      FIN (ip igmp host leave grp (ip grp addr, interface));
      /* Generate a trace message */
      if (LTRACE IGMP)
            ip address print (ip addr str, ip grp addr);
            sprintf (msq0, "IP Group Address
                                                                   %s",
ip addr str);
            sprintf (msgl, "IP Interface
                                                                   %d",
interface);
            op prg odb print major ("Received a Leave request from an
application for: ", msg0, msg1, OPC NIL);
      /* Get the list corresponding to the given interface */
      ip grp intf list ptr = (List *) op prg list access (ip grp list ptr,
interface);
    /* Determine the number of elements in the list */
      num grp elems = op prg list size (ip grp intf list ptr);
/* Access the group membership information data structure for the given */
                              from the list
      /* IP group address
      for (i=0; i<num grp elems; i++)</pre>
            /* Access ith element from the list */
            grp elem ptr = (IpT Igmp Host Grp Elem *)
                  op prg list access (ip grp intf list ptr, i);
/* If this is the list element we are looking for, break from the loop */
            if (ip address equal (grp elem ptr->ip grp addr, ip grp addr))
```

```
{
                  break;
      /* Sanity check */
      if (i == num grp elems)
            /** The group membership information for the given group address
and interface doesn't exist in the list **/
            /* Report a log message */
            ip address print (ip_addr_str, ip_grp_addr);
            ip igmp host log found no grp info (ip addr str, interface);
            /* Generate a trace message */
            if (LTRACE IGMP)
sprintf (msg0, "IP Group Address
                                                      %s", ip addr str);
sprintf (msgl, "IP Interface
                                                       %d", interface);
strcpy (msq2, "doesn't exist in the list, which is maintained by IGMP Host
process.");
op prg odb print major ("Group membership information for: ", msq0, msq1,
msg2, OPC \overline{\text{NIL}});
      else
            /* An application on the node has left this group. Decrement
num_of_apps field */
            grp elem ptr->num of apps--;
            if (grp elem ptr->num of apps == 0)
                  /* Generate a trace message */
                  if (LTRACE IGMP)
                        ip address print (ip addr str, ip grp addr);
                        sprintf (msq0, "IP Group Address :
      %s", ip addr str);
                        sprintf (msg1, "IP Interface
      %d", interface);
                        strcpy (msg2, "from the table, which is maintained by
IGMP Host process.");
                        op prg odb print major ("Removing the group
membership information for: ", msg0, msg1, msg2, OPC NIL);
                  /* Remove the group membership information from the list */
                  op prg list remove (ip grp intf list ptr, i);
                  /* Cancel the delay timer for this group membership, if its
running */
                  if (grp elem ptr->timer on flag == OPC TRUE)
                        op ev cancel (grp elem ptr->delay timer evh);
```

```
/* If this is the last node to send a Report message, send
a Leave message */
                  if (grp elem ptr->report sent flag == OPC TRUE)
                        /* Generate a trace message */
                        if (LTRACE IGMP)
                              ip address print (ip addr str, ip grp addr);
                              sprintf (msg0, "IP Group Address"
      %s", ip addr str);
                              sprintf (msgl, "IP Interface
      %d", interface);
                              op prg odb print major ("Sending an IGMP Leave
Group message for: ", msg0, msg1, OPC_NIL);
                        /* Create an IGMP Leave Group message */
                        igmp msg pkptr = ip igmp create igmp msg
(IpC Igmp Leave Group Msg, 0, grp elem ptr->ip grp addr);
                        /* Update IGMP messages sent statistics
                        ip igmp host igmp msgs sent stat update
(op pk total size get (igmp msg pkptr));
                        /* Create an IP datagram to transmit the Leave group
message */
                        ip dgram pkptr = ip igmp create ipdgram
(igmp msg pkptr, ip address create (IPC ALL ROUTERS MULTICAST ADDR));
                        /* Schedule a procedure interrupt for the current
simulation time to invoke the IP process ^{\star}/
                        op intrpt schedule call (op sim time (),
grp elem ptr->interface, &ip igmp host ip process invoke, ip dgram pkptr);
                  ip_igmp_host_grp_elem_dealloc (grp_elem_ptr);
            }
      FOUT;
      }
static void
mip mn register (int lifetime, MipT MN Agent Info agent info, Boolean direct)
      IpT Port Info
                        port info;
      double
                                    retry timer;
      Ici*
                        reg ici ptr;
      /** PURPOSE: Registers with home agent.**/
      /** REQUIRES: lifetime it is requesting and care of address to use.
      /** EFFECTS: An interrupt sent to the reg manager.**/
      FIN (mip mn register (lifetime, agent info, direct));
      /* Create Ici for mobile IP module. */
      reg ici ptr = op ici create ("mobile ip reg ici");
```

```
/* Set the foreign agent address for later reference. */
      agent address = agent info.address;
      current agent pref level = agent info.pref level;
      current roaming intf = agent info.incoming intf ptr;
      /* Set the appropriate values on the ici. */
      op ici attr set (reg ici ptr, "reg type", MipC Reg Type Req);
      op ici attr set (reg ici ptr, "home address", inet ipv4 address get
(home address));
      op ici attr set (reg ici ptr, "home agent", inet ipv4 address get
(ha address));
      op ici attr set (reg ici ptr, "lifetime req", lifetime);
      op ici attr set (reg ici ptr, "dest address", inet ipv4 address get
(agent address));
      op ici attr set (reg ici ptr, "identification", ++reg id);
      op ici attr set (reg ici ptr, "s", simultaneous binding);
      /* What is the type of the current registration? */
      if (direct)
            /* Diect to the home agent. */
            direct reg = OPC TRUE;
            op ici attr set (reg ici ptr, "care of address", OPC NIL);
      else
            direct reg = OPC FALSE;
            op ici attr set (reg ici ptr, "care of address",
inet ipv4 address get (agent address));
      /* Send the registration packet out. */
      op ici install (reg ici ptr);
      op intrpt schedule process (proc info struct ptr->mip reg mgr phndl,
op sim time (), 0);
      op ici install (OPC NIL);
      /* need to reset the default gateway to the agent, that are registering
with. */
      port_info = ip_rte_port_info from tbl index (module data,
ip rte intf index get (agent info.incoming intf ptr));
      if (!default gateway)
            Ip Cmn Rte Table Entry Add (module data->ip route table, OPC NIL,
                  IpI Default Addr, IpI Default Addr,
agent info.incoming intf ptr->addr range ptr->address, port info,
                  O, IP CMN RTE TABLE UNIQUE ROUTE PROTO ID
(IPC DYN RTE MOBILE IP, IPC NO MULTIPLE PROC), 0, OPC NIL);
            /* Cache info. */
            default gateway = OPC TRUE;
      else
            Ip Cmn Rte Table Entry Update (module data->ip route table,
IpI Default Addr, IpI Default Addr,
```

```
last default addr, IP CMN RTE TABLE UNIQUE ROUTE PROTO ID
(IPC DYN RTE MOBILE IP, IPC NO MULTIPLE PROC),
                  agent info.incoming intf ptr->addr range ptr->address,
port info, 0, OPC NIL);
      /* Cache info. */
      last default addr = ip rte intf addr get
(agent info.incoming intf ptr);
      /* need to schedule retry in case that do not get the answer back. */
      retry timer = reg info.interval * pow (2.0, (double) retry counter);
      if (retry timer > (double) (reg info.req lifetime))
            retry_timer = (double) (reg info.req lifetime);
      reg retry timer ehndl = op intrpt schedule self (op sim time () +
retry timer, MipC MN Timer Retry);
      retry counter++;
      if (MIP TRACE)
           {
            op prg odb print major ("Trying registering with HA.", OPC NIL);
      FOUT;
      }
static void
mip mn agent timer update (double new lifetime)
      /** PURPOSE: Update the timer for agent timeouts.**/
      /** REQUIRES: new lifetime value from the last agent ad.
      /** EFFECTS: timer event handle gets updated.**/
      FIN (mip mn agent timer update (new lifetime));
      /* Cancel the current handle. */
      op ev cancel if pending (agent timer ehndl);
      /* Schedule new one. */
      agent timer ehndl = op intrpt schedule self (new lifetime,
MipC MN Timer Agent);
      FOUT;
      }
static Packet* /* Switch the Active/Standby mode. */
ip_icmp_echo_request_packet_create (int req index)
      {
      Packet*
                              req pkptr;
      FIN (ip icmp echo request packet create (req index));
      /* Create an ICMP packet to switch Active/Standby mode.
      req pkptr = op pk create fmt ("ip icmp echo");
```

```
/* Set the "type" field to indicate that this is a */
                                                                          */
      /* request packet.
      op pk nfd set (req pkptr, "type", IpC Icmp Echo Request);
      op pk nfd set (req pkptr, "identifier", req index);
      op pk nfd set (req pkptr, "sequence number", ping specs ptr
[req index].seq number);
                                                                          */
      /* set the source node object id in the packet.
      op pk nfd set (req pkptr, "source module objid", (double) my objid);
^{\prime \star} Increment the sequence number for the next message send operation. ^{\star \prime}
      ping specs ptr [req index].seq number++;
      if (ping specs ptr [req index].ping pattern ptr->pkt size > 0)
            {
            op pk bulk size set (req pkptr, ping specs ptr
[req index].ping pattern ptr->pkt size * 8);
      FRET (req pkptr);
      }
static void
ip icmp request timeout (void *state ptr, int index)/* Me:Check timeout for
the connection*/
      IpT Icmp_Ping_Specs*
                              ping spec ptr;
                                    dest host name [OMSC HNAME MAX LEN];
      char
      FIN (ip icmp request timeout (void *state ptr, int index));
      /* Get the ping spec ptr from the state ptr.
      ping spec ptr = (IpT Icmp Ping Specs *)state ptr;
      oms tan hname get (ping spec ptr->dest objid, dest host name);
      op prg log entry write (ip icmp timeout log handle,
            "ERROR(S):\n"
            " The echo request for destination (%s) \n"
            " sent at (\%.2f) seconds failed to \n"
            " receive a response before the timeout \n"
            " interval of (\%.5f) seconds. \n"
            "\n",
            dest host name, ping spec ptr [index].start time,
            ping spec ptr [index].ping pattern ptr->timeout, ping spec ptr
[index].ping pattern ptr->timeout);
      FOUT;
static void
```

```
ip pim sm join prune msg (void) /* Me:Function Join/Standby message for
keeping route */
      IpT Pim Sm Rte Entry*
                                          rpt rte entry ptr;
      IpT Pim Sm Rte Entry*
                                          rte entry ptr;
                                          rp_ip_addr;
      IpT Address
      IpT_Pim_Sm_Msg*
                                          pim sm msg ptr;
      IpT Pim Sm Join Prune Msg*
                                          join prune msg ptr;
      IpT Pim Sm Join Prune List Elem*
                                           join prune lelem ptr;
      IpT Pim Sm Join Prune Src*
                                          join prune src ptr;
                                          i, j;
      Boolean
                                          send join prune = OPC FALSE;
      FIN (ip pim sm join prune msg ());
      /* Fist check that PIM-SM is supported on the interface.
      if (intf_array [pkt_recvd intf].pim sm status == OPC FALSE)
            /* The interface does not support PIM-SM. */
            op pk destroy (pimsm pkptr);
            op pk destroy (ip dgram pkptr);
            FOUT;
      /* Obtain the Join/Prune message from the PIM-SM packet
      op pk nfd access (pimsm pkptr, "message", &pim sm msg ptr);
      join prune msg ptr = (IpT Pim Sm Join Prune Msg *) pim sm msg ptr-
>msg ds ptr;
      if (ip pim sm is my address (join prune msg ptr-
>upstream neighbor addr) == OPC TRUE)
            /* Generate trace messages
#ifndef OPD NO DEBUG
            if (LTRACE PIM SM JOIN PRUNE)
                                                                   ip addr str
                  char
[IPC ADDR STR LEN];
                                                                   info0
                  char
[256], info1 [256];
                                                                   msq0 [256];
                  ip address print (ip addr str, join prune msg ptr-
>upstream neighbor addr);
                  sprintf (info0, "Upstream Neighbor Address
      %s", ip addr str);
                  sprintf (infol, "Number of Groups
      %d", join prune msg ptr->num grps);
                  sprintf (msg0, "Received a PIM-SM Join/Prune message on
interface, %d with: ", pkt recvd intf);
                  op prg odb print major (msg0, info0, info1, OPC NIL);
#endif
      /* For each IP group in the list, process the Join and Prune list */
            for (i=0; i<join_prune_msg_ptr->num_grps; i++)
```

```
/* Obtain ith element from the Join/Prune list */
                  join prune lelem ptr = (IpT Pim Sm Join Prune List Elem *)
op prg list access (join prune msg ptr->join prune lptr, i);
                  /* Generate trace messages
#ifndef OPD NO DEBUG
                  if (LTRACE PIM SM JOIN PRUNE)
                        char
                                    info0 [256];
                                    grp_addr_str [IPC ADDR STR LEN];
                        char
                        ip address print (grp addr str, join prune lelem ptr-
>grp addr);
                        sprintf (info0, "IP Group Address
      %s", grp addr str);
                        op prg odb print major ("Processing Join/Prune list
in the PIM-SM Join/Prune message for the group: ", info0, OPC NIL);
#endif
                  /* Process each source in the Join list */
                  for (j=0; j<join prune lelem ptr->num join src; j++)
                        /* Obtain the jth element from the Join list
                        join prune src ptr = (IpT Pim Sm Join Prune Src *)
op prg list access (join prune lelem ptr->join src lptr, j);
                        /* Generate trace messages
#ifndef OPD NO DEBUG
                        if (LTRACE PIM SM JOIN PRUNE)
                              char
      ip addr str [IPC ADDR STR LEN];
      info0 [256], info1 [256], info2 [256];
                              ip address print (ip addr str,
join prune src ptr->src addr);
                              sprintf (info0, "Source IP Address
      %s", ip addr str);
                              sprintf (infol, "WC-bit
      %d", join_prune_src_ptr->wc bit);
                              sprintf (info2, "RPT-bit
      %d", join prune src ptr->rpt bit);
                              op prg odb print major ("Processing the
following Source element in the Join list: ", info0, info1, info2, OPC NIL);
#endif
                        /* Get the route entry for this group
                        rte entry ptr = ip pim sm mcast rte entry get
(join_prune_lelem_ptr->grp_addr, join_prune_src_ptr->src_addr,
                  join prune src ptr->wc bit);
                        /* Print out information about the route entry that
was found. */
#ifndef OPD NO DEBUG
```

```
ip_pim_sm_entry_exists_odb_print (rte_entry ptr);
#endif
                        if ((join_prune_src_ptr->wc_bit == OPC TRUE) &&
(join_prune_src_ptr->rpt_bit == OPC TRUE))
                              if (rte_entry_ptr == OPC_NIL)
{
                                    rte_entry_ptr =
ip pim sm rte entry wc create (join prune src ptr->src addr,
      join prune lelem ptr->grp addr);
                                     /* Add this route entry to the multicast
route table */
                  ip_pim_sm_mcast_rte_entry_add (rte_entry_ptr, OPC_TRUE);
                                    /* This case require a join/prune message
to be sent. */
                  send join prune = OPC TRUE;
                              if (rte entry ptr->in intf addr !=
inet_ipv4_address_get (ip_dgram_fdptr->src_addr))
                                    ip pim sm oif table add (rte entry ptr,
pkt_recvd_intf, inet_ipv4_address_get (ip_dgram_fdptr->src_addr),
                                          join prune msg ptr->hold time,
OPC FALSE);
      ip pim sm all spt rte entries out intf add (rte entry ptr,
                  pkt recvd intf,
                  inet ipv4 address get (ip dgram fdptr->src addr),
                  join_prune_msg_ptr->hold_time,
                  OPC FALSE,
                  join_prune_lelem_ptr->prune_src_lptr);
                        else if ((join_prune_src_ptr->wc_bit == OPC_FALSE) &&
(join prune src ptr->rpt bit == OPC FALSE))
                              if (rte_entry_ptr == OPC_NIL)
{
                                    rp_ip_addr = ip_pim_sm_rp_addr_get
(rp_hash_table_ptr, rp_lptr, join_prune_lelem_ptr->grp_addr,
bsr hash mask length);
                        /* If RP information is not found, its an error */
                                     if (rp ip addr == IPC ADDR INVALID)
```

```
/* Report a log message */
      ipnl protwarn mcast rp unknown log add (join prune lelem ptr->grp addr,
ip module data ptr->node id);
                                     rpt rte entry ptr =
ip pim sm mcast rte entry get (join prune lelem ptr->grp addr, rp ip addr,
OPC TRUE);
      rte_entry_ptr = ip_pim_sm_rte_entry_spt_create (join_prune_src_ptr-
>src addr, join prune lelem ptr->grp addr,
                                           rpt_rte_entry_ptr, OPC FALSE);
                                     if (rpt rte entry ptr != OPC NIL)
      op prg list elems copy (rpt rte entry ptr->out intf table lptr,
rte entry ptr->out intf table lptr);
                  /st Add this route entry to the multicast route table st/
                                     ip_pim_sm_mcast_rte_entry_add
(rte entry ptr, OPC FALSE);
            /\!\!\!\!\!^{\star} This case requires sending a join/prune message.
                                                                          */
                                     send join prune = OPC TRUE;
                               else
                                     if (rte_entry_ptr->rpt_flag == OPC TRUE)
                                           /* Clear the RPT-bit
                                           ip pim sm entry clear rpt bit
(rte entry ptr, OPC FALSE);
                  /* This case requires sending a join/prune message
                                                                          * /
                                           send join prune = OPC TRUE;
                               if (inet_ipv4_address_get (ip_dgram_fdptr-
>src_addr) != rte_entry_ptr->in_intf_addr)
                                     ip pim sm oif table add (rte entry ptr,
pkt_recvd_intf, inet_ipv4_address_get (ip_dgram_fdptr->src_addr),
                                           join prune msg ptr->hold time,
OPC FALSE);
                        if (send_join_prune == OPC TRUE)
                               ip pim sm send join prune msg (rte entry ptr,
OPC TRUE);
                                                                    */
                        /* Reset the flag to FALSE.
                        send_join_prune = OPC FALSE;
                        }
                  /* Process each source in the Prune list */
                  for (j=0; j<join prune lelem ptr->num prune src; j++)
```

```
/* Obtain the jth element from the Prune list
                        join prune src ptr = (IpT Pim Sm Join Prune Src *)
op prg list access (join prune lelem ptr->prune src lptr, j);
                        /* Get the route entry for this group
                        rte entry ptr = ip pim sm mcast rte entry get
(join prune lelem ptr->grp addr, join prune src ptr->src addr,
                  join prune src ptr->wc bit);
                        /* Generate trace messages
#ifndef OPD NO DEBUG
                        if (LTRACE PIM SM JOIN PRUNE)
                              char
      ip_addr_str [IPC_ADDR_STR_LEN];
                              char
      info0 [256], info1 [256], info2 [256];
                              ip address print (ip addr str,
join prune src ptr->src addr);
                              sprintf (info0, "Source IP Address
      %s", ip addr str);
                              sprintf (infol, "WC-bit
      %d", join prune src ptr->wc bit);
                              sprintf (info2, "RPT-bit
      %d", join prune src ptr->rpt bit);
                              op prg odb print major ("Processing the
following Source element in the Prune list: ", info0, info1, info2, OPC_NIL);
#endif
#ifndef OPD NO DEBUG
                        ip_pim_sm_entry_exists_odb_print (rte_entry_ptr);
#endif
                        if ((join prune src ptr->wc bit == OPC FALSE) &&
(join prune src ptr->rpt bit == OPC TRUE))
      if (rte entry ptr == OPC NIL)
                                    /* Get the RP address for the group */
                                    rp ip addr = ip pim sm rp addr get
(rp hash table ptr, rp lptr, join prune lelem ptr->grp addr,
bsr hash mask length);
                        /* If RP information is not found, its an error */
                                    if (rp ip addr == IPC ADDR INVALID)
      ipnl protwarn mcast rp unknown log add (join prune lelem ptr->grp addr,
ip module data ptr->node id);
                        /* Get the (*, G) entry */
                                    rpt rte entry ptr =
ip pim sm mcast rte entry get (join prune lelem ptr->grp addr, rp ip addr,
OPC TRUE);
```

```
if (rpt_rte_entry_ptr != OPC_NIL)
                                          rte entry ptr =
ip pim sm rte entry spt create (join prune src ptr->src addr,
                                    join prune lelem ptr->grp addr,
                                    rpt rte entry ptr,
                                    OPC_FALSE);
                                          ip pim sm entry set rpt bit
(rte entry ptr, OPC FALSE);
                                          rte_entry_ptr->in_intf =
rpt rte entry ptr->in intf;
                                          rte_entry_ptr->in_intf_addr =
ip_address_copy(rpt_rte_entry_ptr->in_intf_addr);
                                          op prg list elems copy
(rpt_rte_entry_ptr->out_intf_table_lptr, rte_entry_ptr->out_intf_table_lptr);
                  /st Add this route entry to the multicast route table st/
                                          ip pim sm mcast rte entry add
(rte_entry_ptr, OPC_FALSE);
                        if (rte_entry_ptr != OPC_NIL)
{
                              if (op prg list size (rte entry ptr-
>out intf table lptr) == 0)
                        /* Check if the data timer should be reset.
                                   if ((join_prune_src_ptr->wc_bit ==
OPC_FALSE) && (join_prune_src_ptr->rpt_bit == OPC_TRUE))
            ip_pim_sm_reset_data_timer (rte_entry_ptr,
LTRACE PIM SM JOIN PRUNE);
                                                                         */
            /* Nothing more needs to be done. Just continue.
                                    continue;
ip_pim_sm_oif_table_remove (rte_entry_ptr->out_intf_table_lptr,
pkt recvd intf, OPC FALSE);
                              if ((join_prune_src_ptr->wc_bit == OPC_TRUE) &&
(join prune src ptr->rpt bit == OPC TRUE))
      ip pim sm all spt rte entries out intf remove (join prune lelem ptr-
>grp addr, pkt recvd intf, OPC FALSE);
```

```
if (op prg list size (rte entry ptr-
>out intf table lptr) == 0)
                                    ip pim sm send join prune msg
(rte entry ptr, OPC FALSE);
                                    if (((join prune src ptr->wc bit ==
OPC FALSE) && (join prune src ptr->rpt bit == OPC TRUE)) &&
                                          ((rte entry ptr->wc flag ==
OPC FALSE) && (rte entry ptr->rpt flag == OPC FALSE)))
                  /* Set the RPT bit to TRUE for this route entry.
                                          ip pim sm entry set rpt bit
(rte entry ptr, OPC FALSE);
                  /* Set the SPT bit to FALSE for this route entry.
                                          ip pim sm entry clear spt bit
(rte entry ptr, OPC FALSE);
                  /* Remove the route entry from the multicast route
table, only if its is not a (S, G)RPT-bit entry */
                                    if (!((rte entry ptr->wc flag ==
OPC FALSE) && (rte entry ptr->rpt flag == OPC TRUE)))
#ifndef OPD NO DEBUG
                                          if (LTRACE PIM SM JOIN PRUNE)
                                                op prg odb print major ("The
Out Interface table size of the route entry became zero.", OPC NIL);
#endif
ip pim sm mcast rte entry remove (rte entry ptr->grp addr, rte entry ptr-
>src addr, rte entry ptr->wc flag);
                                          continue;
                                    }
                              if ((join prune src ptr->wc bit == OPC FALSE)
&& (join prune src ptr->rpt bit == OPC TRUE))
                                    ip pim sm reset data timer
(rte entry ptr, LTRACE PIM_SM_JOIN_PRUNE);
                  }
      /* We no longer need the PIM-SM packet and the */
      /* IP datagram. Destroy them */
      op pk destroy (pimsm pkptr);
      op_pk_destroy (ip_dgram_pkptr);
```

```
FOUT;
      }
static void
mip mn tunneled pk stat write (Packet*
                                         pk ptr)
      /** PURPOSE: Write statistic for tunneld packets received.**/
      FIN (mip mn tunneled pk stat write (pk ptr));
      /* Write the stats. */
      op stat write (tunneled pk rcvd sec sh, 1.0);
      op stat write (tunneled bit rcvd sec sh, op pk total size get
(pk_ptr));
     op stat write (tunneled pk rcvd sec sh, 0.0);
      op stat write (tunneled bit rcvd sec sh, 0.0);
      FOUT;
      }
static void
mip mn agent solicit pk send (void)
                       solicit pkptr;
     Packet*
                        solicit interval;
     double
      /** PURPOSE: Send the ICMP agent solicitation packet.**/
      /** REOUIRES: none. **/
      /** EFFECTS: Packet will be given to IP to handle.**/
      FIN (mip mn agent solicit pk send (void));
      /* Time to send out the solicitation. */
      solicit pkptr = op pk create fmt ("mobile ip irdp solicit");
      /* Send the packet out. */
      module data->ip ptc mem.child pkptr = mip sup irdp pkt encapsulate
            (solicit pkptr, home address, subnet bcast addr,
IcmpC Type IRDP Sol);
      /* Record some stats. */
      op stat write (irdp sent pkts sh, 1.0);
      op stat write (irdp sent bits sh, op pk total size get (module data-
>ip ptc mem.child pkptr));
      op_stat_write (g_irdp_sent bits sh, op pk total size get (module data-
>ip ptc mem.child pkptr));
      op_stat_write (g_irdp_sent bits sh, 0.0);
      /* Invoke IP to handle the packet. */
      op pro invoke (proc info struct ptr->ip phndl, OPC NIL);
      /* Schedule the next transmission. */
      if (++solicit count > 3)
            solicit interval = MipC MN Solicit Min Interval * pow (2.0,
(double) (solicit count - 3));
            if (solicit interval > MipC MN Solicit Max Interval)
```

```
solicit interval = MipC MN Solicit Max Interval;
      else
            solicit interval = MipC MN Solicit Min Interval;
      solicit timer ehndl = op intrpt schedule self (op sim time () +
solicit interval, MipC MN Timer Solicit);
      FOUT;
      }
static void
mip mn agent cache update (MipT MN Agent Info* agent info, InetT Address
      new agent address,
      double life time, int pref level, MipT Invocation Info*
invoke info ptr)
      /** PURPOSE: Update Agent cache based on rule. (higher pref level or
same or lower if the current one expired) **/
      /** REQUIRES: new FA address and its lifetime and pref level.
      /** EFFECTS: the cache will be updated if it matches the criteria.**/
      FIN (mip mn agent cache update (agent info, new agent address,
life time, pref level, invoke info ptr));
      if ((pref level >= agent info->pref level) ||
            (agent info->lifetime < op sim time ()) ||
            inet address equal (agent info->address, new agent address))
            /* Update the FA cache information. */
            agent info->address = new agent address;
            agent info->lifetime = op sim time () + life time;
            agent info->pref level = pref level;
            agent info->incoming intf ptr = ip rte intf tbl access
                  (module data, invoke info ptr->rte info ici ptr-
>intf recvd index);
      FOUT;
      }
static void
mip mn ip pk handle (MipT Invocation Info* invoke info ptr)
      Prohandle tmp phndl;
      FIN (mip mn ip pk handle (invoke info ptr));
      /* See if there are any visiting MN with the address. */
      if (mip sup visitor search by addr (proc info struct ptr-
>node visitor list lptr,
           invoke info_ptr->rte_info_ici_ptr->dest_addr, &tmp_phndl) ==
OPC COMPCODE SUCCESS)
            /* will let the FA agent to handle this packet. */
```

```
op pro invoke (tmp phndl, invoke info ptr);
      else
             if ((inet address range check (invoke info ptr->rte info ici ptr-
>dest addr,
                   &proc_info_struct_ptr->intf info ptr->inet addr range)) &&
                   !inet address equal (invoke info ptr->rte info ici ptr-
>dest addr, home address) &&
                   !inet address equal (invoke info ptr->rte info ici ptr-
>dest addr, ha address))
                   /* This address falls in the same range but not for me.
Destroy packet. */
                   mip sup pk cleanup (invoke info ptr);
             else
                   if (current roaming intf)
                          if (current roaming intf == invoke info ptr-
>interface ptr)
                                 /* This is going out on the roaming interface.
Send to the current agent. */
                                if (inet address valid (agent address))
                                       invoke info ptr->rte info ici ptr-
>next addr = agent address;
                          else
                                 /* Unknown address. Let IP handle it
generically. */
                          }
                   }
      FOUT;
      }
static void
mip mn ad packet parse (Packet* pkptr, int* h, int* f, InetT Address*
agent addr,
      int* lifetime, int* pref)
      IpT Address
                          tmp addr;
      /* Helper macro to parse information from the advertisement packet
received. */
      FIN (mip mn ad packet parse (...));
      /* Access the information from the packet received. */
      op_pk_nfd_access (pkptr, "H", h);
      op_pk_nfd_access (pkptr, "F", f);
op_pk_nfd_access (pkptr, "Agent Address", &tmp_addr);
op_pk_nfd_access (pkptr, "lifetime", lifetime);
```

```
op pk nfd access (pkptr, "Preference Level", pref);
      /* Convert the V4 address to v6 for internal reference. */
      *agent addr = inet address from ipv4 address create (tmp addr);
      FOUT;
      }
static void
mip mn agent solicit pk send adv (void) /*Me: Sending Mobile IP to join in
      Packet*
                        solicit pkptr adv;
      double
                        solicit interval;
      /** PURPOSE: Send the ICMP agent solicitation packet.**/
      /** REOUIRES: none.
                              **/
      /** EFFECTS: Packet will be given to IP to handle.**/
      FIN (mip mn agent solicit pk send (void));
      /* Time to send out the solicitation. */
      usleep(10000000); // 10 secs
      solicit pkptr adv = op pk create fmt ("mobile ip irdp solicit");
      /* Send the packet out. */
      module data->ip ptc mem.child pkptr = mip sup irdp pkt encapsulate
            (solicit pkptr adv, home address, subnet bcast addr,
IcmpC_Type_IRDP Sol);
      /* Record some stats. */
      op stat write (irdp sent pkts sh, 1.0);
      op stat write (irdp sent bits sh, op pk total size get (module data-
>ip ptc mem.child pkptr));
      op stat write (g irdp sent bits sh, op pk total size get (module data-
>ip ptc mem.child pkptr));
      op stat write (g irdp sent bits sh, 0.0);
      /* Invoke IP to handle the packet. */
      op pro invoke (proc info struct ptr->ip phndl, OPC NIL);
      /* Schedule the next transmission. */
      if (++solicit count > 3)
            solicit interval = MipC MN Solicit Min Interval * pow (2.0,
(double) (solicit count - 3));
            if (solicit interval > MipC MN Solicit Max Interval)
                  solicit interval = MipC MN Solicit Max Interval;
      else
            solicit interval = MipC MN Solicit Min Interval;
```

```
solicit timer ehndl = op intrpt schedule self (op sim time () +
solicit interval, MipC MN Timer Solicit);
      FOUT;
      }
//me
/* End of Function Block */
/* Undefine optional tracing in FIN/FOUT/FRET */
/* The FSM has its own tracing code and the other */
/* functions should not have any tracing.
#undef FIN TRACING
#define FIN TRACING
#undef FOUTRET TRACING
#define FOUTRET TRACING
#if defined ( cplusplus)
extern "C" {
#endif
      void mobile_ip_mn (OP_SIM_CONTEXT ARG OPT);
     VosT_Obtype _op_mobile_ip_mn_init (int * init_block_ptr);
     void op mobile ip mn diag (OP SIM CONTEXT ARG OPT);
     void op mobile ip mn terminate (OP SIM CONTEXT ARG OPT);
     VosT Address op mobile ip mn alloc (VosT Obtype, int);
      void op mobile ip mn svar (void *, const char *, void **);
#if defined ( cplusplus)
\} /* end of 'extern "C"' */
#endif
/* Process model interrupt handling procedure */
void
mobile ip mn (OP SIM CONTEXT ARG OPT)
#if !defined (VOSD NO FIN)
     int op block origin = 0;
#endif
      FIN MT (mobile ip mn ());
            /* Temporary Variables */
            Boolean is ip pk = OPC FALSE;
                       is_ad_reception = OPC FALSE;
            Boolean
                       is ha ad reception = OPC FALSE;
            Boolean
            Boolean
                        is fa ad reception = OPC FALSE;
            Boolean
                        is solicitation time = OPC FALSE;
            Boolean
                        is valid fa candidate = OPC FALSE;
                        is ha timeout = OPC FALSE;
            Boolean
                        is_fa_timeout = OPC FALSE;
            Boolean
                        is timeout = OPC FALSE;
            Boolean
                        is retry = OPC FALSE;
            Boolean
                        is_ha_reg_success = OPC FALSE;
            Boolean
```

```
Boolean
                       is fa reg success = OPC FALSE;
           Boolean
                       is reg pk = OPC FALSE;
           Boolean
                       is out of retries = OPC FALSE;
                       is invalid reply = OPC FALSE;
           Boolean
           Boolean
                       is reregister = OPC FALSE;
                       is switch fa = OPC FALSE;
           Boolean
           Objid mip reg cfg objid;
           char ha address str[64];
                       h bit, f bit, reply code, lifetime grant, tmp reg id,
                             irdp lifetime, inv mode, intrpt code,
pref level;
                                   *irdp pkptr, *encap pk ptr;
           Packet
                                   tmp_agent_address;
           InetT Address
           MipT Invocation Info*
                                   invoke info ptr;
           Ici* reg ici ptr;
           /* End of Temporary Variables */
           FSM_ENTER ("mobile_ip mn")
            FSM BLOCK SWITCH
                {
            /*----*/
                  /** state (Init) enter executives **/
                 FSM STATE ENTER FORCED NOLABEL (0, "Init", "mobile ip mn
[Init enter execs]")
                       FSM PROFILE SECTION IN ("mobile ip mn [Init enter
execs]", state0 enter exec)
                       /* Access the parent memory. */
                       proc info struct ptr = (MipT Proc Info*)
op pro parmem access ();
                       /* Get the type of agent I am. (MN or MR) */
                       mip_node_type = proc_info_struct ptr->node type;
                       /* Access the module wide memory. */
                       module data = (IpT Rte Module Data*)
op pro modmem access ();
                       /* Keep the subnet bcast address for later use. */
                       subnet bcast addr = inet rte intf broadcast addr get
                             (proc info struct ptr->intf info ptr,
InetC Addr Family v4);
                       /* Parse the home agent interface address. */
                       op ima obj attr get (proc info struct ptr->cfg objid,
"Home Agent IP Address",
                             &ha address str);
                       ha address = inet address create (ha address str,
InetC Addr Family_v4);
                       if (inet address equal (ha address,
InetI Invalid v4 Addr))
                             op sim end ("An invalid address was configured
as Home Agent.", "", "", "");
```

```
/* local interface address. */
                        home address = inet rte intf addr get
(proc info struct ptr->intf info ptr, InetC Addr Family v4);
                        agent address = InetI Invalid v4 Addr;
                        /* Parse the registration related information. */
                        op ima obj attr get (proc info struct ptr->cfg objid,
"Registration Parameters",
                              &mip reg cfg objid);
                        mip_reg_cfg_objid = op_topo_child (mip reg cfg objid,
OPC OBJTYPE GENERIC, 0);
                        op_ima_obj_attr_get (mip reg cfg objid, "Interval",
&(reg info.interval));
                        op ima obj attr get (mip reg cfg objid, "Retry",
&(reg info.retry));
                        op ima obj attr get (mip reg cfg objid, "Lifetime
Request", &(reg_info.req_lifetime));
                        /* Initialize state vars. */
                        reg id = 0;
                        retry counter = 0;
                        latest fa info.address = InetI Invalid v4 Addr;
                        latest fa info.lifetime = 0.0;
                        latest fa info.pref level = 0;
                        current roaming intf = OPC NIL;
                        /* Register some stats. */
                        tunneled bit rcvd sec sh = op_stat_reg ("Mobile
IP. Tunneled Traffic Received (bits/sec)", OPC STAT INDEX NONE,
OPC STAT LOCAL);
                        tunneled_pk_rcvd_sec_sh = op_stat_reg ("Mobile
IP. Tunneled Traffic Received (packets/sec)", OPC STAT INDEX NONE,
OPC STAT LOCAL);
                        /* Find out if need to send out solicitation when
lost. */
                        op ima obj attr get (proc info struct ptr->cfg objid,
"Agent Solicitation",
                              &solicitation);
                        /* See if I am configured on a loopback interface. */
                        loopback intf = ip rte intf is loopback
(proc info struct ptr->intf info ptr);
                        /* Schedule interrupt to send solicitation if
enabled. */
                        if (solicitation)
                              if (loopback intf)
                                     /* will not solicitate if on a loopback
interface. */
                                    op sim end ("Mobile IP currently cannot
support solicitation when configured on loopback interface.",
```

```
OPC NIL, OPC NIL, OPC NIL);
                             /* Schedule an interrupt for the first
solicitation. */
                             solicit count = 0;
                             solicit timer ehndl = op intrpt schedule self
(mip sup activation time calculate (), MipC MN Timer Solicit);
                             /* Register some stats. */
                             irdp_sent_pkts_sh = op_stat_reg ("Mobile
IP.IRDP Traffic Sent (packets)", OPC STAT INDEX NONE, OPC STAT LOCAL);
                             irdp sent bits sh = op stat reg ("Mobile
IP.IRDP Traffic Sent (bits)", OPC STAT INDEX NONE, OPC STAT LOCAL);
                             g irdp sent bits sh = op stat reg ("Mobile
IP.IRDP Traffic Sent (bits/sec)", OPC STAT INDEX NONE, OPC STAT GLOBAL);
                       /* Find out if need to ask for simultaneous binding.
* /
                       op ima obj attr get (proc info struct ptr->cfg objid,
"Simultaneous Binding Support", &simultaneous binding);
                       /* Cache the node objid info. */
                       node objid = op topo parent (op id self ());
                       /* To handle default gateway. */
                       default gateway = OPC FALSE;
                       /* Animation. */
                       if (op_sim_anim ())
                             /* Initialize the view. */
                             mip sup prepare animation ();
                       FSM PROFILE SECTION OUT (state0 enter exec)
                 /** state (Init) exit executives **/
                 FSM STATE EXIT FORCED (0, "Init", "mobile ip mn [Init exit
execs]")
                 /** state (Init) transition processing **/
                 FSM TRANSIT FORCE (1, state1 enter exec, ;, "default", "",
"Init", "Lost", "tr_-1", "mobile_ip_mn [Init -> Lost : default / ]")
           /*----*/
                 /** state (Lost) enter executives **/
                 FSM STATE ENTER UNFORCED (1, "Lost", state1 enter exec,
"mobile ip mn [Lost enter execs]")
                       FSM PROFILE SECTION IN ("mobile ip mn [Lost enter
execs]", state1 enter exec)
                       /* Update the status for debugging. */
```

```
mip sup mn mr status update (node objid,
home address, MipC Mn Mr Status Lost,
                              ha address, agent address);
                        FSM PROFILE SECTION OUT (state1 enter exec)
                  /** blocking after enter executives of unforced state. **/
                  FSM EXIT (3, "mobile ip mn")
                  /** state (Lost) exit executives **/
                  FSM STATE EXIT UNFORCED (1, "Lost", "mobile ip mn [Lost
exit execs]")
                        FSM PROFILE SECTION IN ("mobile ip mn [Lost exit
execs]", state1 exit exec)
                        /* Who invoked me? */
                        op pro invoker (proc info struct ptr->pro hndl,
&inv mode);
                        if (inv mode == OPC PROINV DIRECT)
                              /* one of those timer went off. */
                              is solicitation time = OPC TRUE;
                              if (op_intrpt_code () == MipC MN Timer Solicit)
                                    mip mn agent solicit pk send ();
                        else
                              /* See if are getting an IP packet. */
                              invoke info ptr = (MipT Invocation Info*)
op pro argmem access ();
                              if (invoke info ptr != OPC NIL)
                            have an invocation from IP. What kind thou? */
                                    switch (invoke info ptr->invocation type)
                                          case MipC Invoke Type IRDP:
                            have an IRDP packet. But which kind thou? */
                                                 if (invoke info ptr-
>irdp type == IcmpC Type IRDP Sol)
                           do not want to deal with this packet. */
                                                      mip sup pk cleanup
(invoke info ptr);
                                                       is ip pk = OPC TRUE;
                                                       break;
                        /* This must be an advertisement from an agent. */
                                                op pk nfd get
(invoke info ptr->pk ptr, "data", &irdp pkptr);
```

```
mip mn ad packet parse
(irdp pkptr, &h bit, &f bit, &tmp agent address, &irdp lifetime,
&pref level);
      /* For now, comaprison of the only address in the packet suffice. */
(inet address equal(tmp agent address, ha address))
                                                       if (loopback intf)
                                    /* This is not supported. */
                                                             op_sim_end
("MR/MN when configured on a loopback interface, cannot directly communicate
with HA.",
                                                OPC_NIL, OPC_NIL, OPC_NIL);
                                          /* Ad from my home agent. */
                                                 is ad reception = OPC TRUE;
                        /* Initialize counter before start reg process. */
                                                       retry counter = 0;
                  /* Update the latest ha info structure for later. */
      mip mn agent cache update (&latest ha info, tmp agent address,
                                                             (double)
irdp lifetime, pref level, invoke info ptr);
                                    /* Deregister with HA. */
                                                      mip mn register (0,
latest ha info, OPC TRUE);
                                          /* Update timer for HA timeout. */
      mip mn agent timer update (latest ha info.lifetime);
                                                 else
                                                       if (f bit)
                                                 /* A foreign agent ad. */
                                                 is ad reception = OPC TRUE;
                        /* Initialize counter before start reg process. */
                                                retry counter = 0;
                        /* Update the latest fa info structure for later. */
     mip mn agent cache update (&latest fa info, tmp agent address,
                        (double) irdp lifetime, pref level, invoke info ptr);
                              /* Register with HA using the FA address. */
      mip mn register (reg info.req lifetime, latest fa info, OPC FALSE);
                                    /* Update timer for HA timeout. */
```

```
mip mn agent timer update (latest fa info.lifetime);
                                                       else
/* cannot do anything with this agent who is only HA for other group. */
                                                is ip pk = OPC TRUE;
                                    /* Clean up solicitation if any. */
      if (solicitation && is ad reception)
      op ev cancel if pending (solicit timer ehndl);
                                                /* Clean up. */
      mip sup pk cleanup (invoke info ptr);
      op pk destroy (irdp pkptr);
      break;
                                          case MipC Invoke Type Tunnel Check:
                                                is ip pk = OPC TRUE;
                       should try to decapsulte packet if in MR mode. */
                              if (mip node type == MipC Node Type MR)
         will check first if this packet is tunneling other IP packet */
(mip sup ip in ip decapsulate (invoke info ptr->pk ptr, &encap pk ptr)
OPC COMPCODE SUCCESS)
                                          /* Sanity check on the packet */
                              /* Invoke IP delayed to handle the packet */
      mip sup packet send to ip (module data, encap pk ptr);
                        /* Write stats for the received tunneled packet. */
      mip mn tunneled pk stat write (invoke info ptr->pk ptr);
                  /* Let IP caller know that are handling the packet */
      mip sup pk cleanup (invoke info ptr);
                                                       }
                                                break;
                                          case MipC Invoke Type IP Datagram:
                                                is ip pk = OPC TRUE;
```

```
/* The destination address that is going out on this interface when in MR
mode should be forwarded to either HA or FA. Whoever already registered or
trying to. */
                                   if (mip node type == MipC Node Type MR)
                             /* Handle packet if I am a MR. */
                                   mip_mn_ip_pk_handle (invoke_info_ptr);
                                              else
                       /* Unknown address. Send to the current agent. */
                                   if (inet address valid (agent address))
                                                         invoke info ptr-
>rte info ici ptr->next addr = agent address;
                                              break;
                                         } /* switch (invoke info ptr-
>invocation type) */
                             else
                                   /* Registration arrival. */
                                  op ici destroy (op intrpt ici ());
                                  is reg pk = OPC TRUE;
                             }
                       FSM PROFILE SECTION OUT (state1 exit exec)
                 /** state (Lost) transition processing **/
                 FSM PROFILE SECTION IN ("mobile ip mn [Lost trans
conditions]", state1_trans_conds)
                 FSM INIT COND (IP PK)
                 FSM TEST COND (SOLICITATION TIME)
                 FSM TEST COND (REG PK)
                 FSM TEST COND (AD RECEPTION)
                 FSM TEST LOGIC ("Lost")
                 FSM PROFILE SECTION OUT (state1 trans conds)
                 FSM_TRANSIT_SWITCH
                       FSM CASE TRANSIT (0, 1, state1_enter_exec, ;,
"IP PK", "", "Lost", "Lost", "tr_8", "mobile_ip_mn [Lost -> Lost : IP_PK /
                      FSM_CASE_TRANSIT (1, 1, state1_enter_exec, ;,
"SOLICITATION TIME", "", "Lost", "tr_32", "mobile_ip_mn [Lost -> Lost
: SOLICITATION TIME / ]")
                      FSM CASE TRANSIT (2, 1, state1 enter exec, ;,
"REG PK", "", "Lost", "Lost", "tr_36", "mobile_ip_mn [Lost -> Lost : REG_PK /
]")
                       FSM CASE TRANSIT (3, 4, state4 enter exec, ;,
"AD RECEPTION", "", "Lost", "Pending registration", "tr 1", "mobile ip mn
[Lost -> Pending registration : AD RECEPTION / ]")
            /*----*/
```

```
/** state (At home) enter executives **/
                  FSM STATE ENTER UNFORCED (2, "At home", state2 enter exec,
"mobile ip mn [At home enter execs]")
                        FSM PROFILE SECTION IN ("mobile ip mn [At home enter
execs]", state2 enter exec)
                        /* Update the status for debugging. */
                        mip sup mn mr status update (node objid,
home_address, MipC_Mn_Mr_Status Home,
                              ha address, agent address);
                        FSM PROFILE SECTION OUT (state2 enter exec)
                  /** blocking after enter executives of unforced state. **/
                  FSM EXIT (5, "mobile ip mn")
                  /** state (At home) exit executives **/
                  FSM STATE EXIT UNFORCED (2, "At home", "mobile ip mn [At
home exit execs]")
                        FSM PROFILE SECTION IN ("mobile ip mn [At home exit
execs]", state2 exit exec)
                        /* Who invoked me? */
                        op pro invoker (proc info struct ptr->pro hndl,
&inv mode);
                        if (inv mode == OPC PROINV DIRECT)
                              {
                              /* one of those timer went off. */
                              is ha timeout = OPC TRUE;
                              if (MIP TRACE)
                                    op prg odb print major ("Trying
reregister for timeout occurred.", OPC NIL);
                        else
                              /* See if are getting an IP packet. */
                              invoke info ptr = (MipT Invocation Info*)
op pro argmem access ();
                              if (invoke info ptr != OPC NIL)
                             have an invocation from IP. What kind thou? */
                                    switch (invoke info ptr->invocation type)
                                          case MipC Invoke Type IRDP:
                             have an IRDP packet. But which kind thou? */
                                                if (invoke info ptr-
>irdp type == IcmpC Type IRDP Sol)
                             do not want to deal with this packet. */
```

```
mip sup pk cleanup
(invoke info ptr);
                                                       is_ip_pk = OPC TRUE;
                                                       break;
                        /* This must be an advertisement from an agent. */
                                                 is ad reception = OPC TRUE;
                     have an agent advertisement. But which kind thou? */
                                                 op pk nfd get
(invoke info ptr->pk ptr, "data", &irdp pkptr);
                                                 mip_mn_ad packet parse
(irdp pkptr, &h bit, &f bit, &tmp agent address, &irdp lifetime,
&pref level);
      /* For now, comaprison of the only address in the packet suffice. */
(inet address equal(tmp agent address, ha address))
                                                       if (loopback intf)
                                    /* This is not supported. */
                                                             op sim end
("MR/MN when configured on a loopback interface, cannot directly communicate
with HA.",
                                                                   OPC NIL,
OPC NIL, OPC NIL);
                                                             }
                                          /* Ad from my home agent. */
                        /* Update the latest ha info structure for later. */
      mip mn agent cache update (&latest ha info, tmp agent address,
                                                             (double)
irdp lifetime, pref level, invoke info ptr);
                                    /* Update timer for HA timeout. */
      mip mn agent timer update (latest ha info.lifetime);
                                                 else
                                                       if (f_bit)
                                                 /* A foreign agent ad. */
                        /* Update the latest fa info structure for later. */
      mip mn agent cache update (&latest fa info, tmp agent address,
                                                                   (double)
irdp lifetime, pref level, invoke info ptr);
                                                       else
```

```
/* cannot do anything with this agent who is only HA for other group. */
                                                /* Clean up. */
                                                mip sup pk cleanup
(invoke info ptr);
                                                op pk destroy (irdp pkptr);
                                                break;
                                           case
MipC Invoke Type Tunnel Check:
                                                is ip pk = OPC TRUE;
                     should try to decapsulte packet if in MR mode. */
                                                if (mip node type ==
MipC Node Type MR)
          will check first if this packet is tunneling other IP packet */
(mip sup ip in ip decapsulate (invoke info ptr->pk ptr, &encap pk ptr)
OPC COMPCODE SUCCESS)
                                    /* Sanity check on the packet */
                              /* Invoke IP delayed to handle the packet */
      mip sup packet send to ip (module data, encap pk ptr);
      /* Write stats for the received tunneled packet. */
     mip mn tunneled pk stat write (invoke info ptr->pk ptr);
            /* Let IP caller know that are handling the packet */
     mip sup pk cleanup (invoke info ptr);
                                                break;
                                          case MipC_Invoke_Type_IP_Datagram:
                                                is ip pk = OPC TRUE;
                                                /* The destination address
that is going out on this interface when in MR mode should be forwarded to
either HA or FA. Whoever already registered or trying to. */
                                    if (mip node type == MipC Node Type MR)
                                          /* Handle packet if I am a MR. */
                                    mip mn ip pk handle (invoke info ptr);
```

```
else
                       /* Unknown address. Send to the current agent. */
           invoke info ptr->rte info ici ptr->next addr = agent address;
                                              break;
                                         }/* switch (invoke info ptr-
>invocation type) */
                                   }
                             else
                                   /* Registration arrival. */
                                   op ici destroy (op intrpt ici ());
                                   is_reg_pk = OPC_TRUE;
                       FSM PROFILE SECTION OUT (state2 exit exec)
                 /** state (At home) transition processing **/
                 FSM PROFILE SECTION IN ("mobile ip mn [At home trans
conditions]", state2 trans conds)
                 FSM INIT COND (HA TIMEOUT)
                 FSM TEST COND (IP PK)
                 FSM TEST COND (AD RECEPTION)
                 FSM TEST COND (REG PK)
                 FSM TEST LOGIC ("At home")
                 FSM PROFILE SECTION OUT (state2 trans conds)
                 FSM TRANSIT SWITCH
                       FSM CASE_TRANSIT (0, 5, state5_enter_exec, ;,
"HA TIMEOUT", "", "At home", "Check FA cache", "tr 5", "mobile ip mn [At home
-> Check FA cache : HA TIMEOUT / ]")
                       FSM CASE TRANSIT (1, 2, state2 enter exec, ;,
"IP PK", "", "At home", "At home", "tr 9", "mobile ip mn [At home -> At home
: IP PK / ]")
                       FSM CASE TRANSIT (2, 2, state2 enter exec, ;,
"AD RECEPTION", "", "At home", "At home", "tr_26", "mobile_ip_mn [At home ->
At home : AD RECEPTION / ]")
                       FSM CASE TRANSIT (3, 2, state2 enter exec, ;,
"REG PK", "", "At home", "At home", "tr_35", "mobile_ip_mn [At home -> At
home : REG PK / ]")
           /*----*/
                 /** state (Away) enter executives **/
                 FSM STATE ENTER UNFORCED (3, "Away", state3 enter exec,
"mobile ip mn [Away enter execs]")
                       FSM PROFILE SECTION IN ("mobile ip mn [Away enter
execs]", state3 enter exec)
                       /* Update the status for debugging. */
```

```
mip sup mn mr status update (node objid,
home address, MipC Mn Mr Status Foreign,
                              ha address, agent address);
                        FSM PROFILE SECTION OUT (state3 enter exec)
                  /** blocking after enter executives of unforced state. **/
                  FSM EXIT (7, "mobile ip mn")
                  /** state (Away) exit executives **/
                  FSM STATE EXIT UNFORCED (3, "Away", "mobile ip mn [Away
exit execs]")
                        FSM PROFILE SECTION IN ("mobile ip mn [Away exit
execs]", state3 exit exec)
                        /* Who invoked me? */
                        op pro invoker (proc info struct ptr->pro hndl,
&inv mode);
                        if (inv_mode == OPC PROINV DIRECT)
                              /* one of those timer went off. */
                              intrpt code = op intrpt code ();
                              switch (intrpt code)
                                    case MipC MN Timer Agent:
                                          is fa timeout = OPC TRUE;
                                          break;
                                  case MipC MN Timer Rereg:
                                          is reregister = OPC TRUE;
                                          break;
                                    }
                              if (MIP TRACE)
                                    op_prg_odb_print_major ("Trying
reregister for timeout occurred.", OPC NIL);
                        else
                              /* See if are getting an IP packet. */
                              invoke info ptr = (MipT Invocation Info*)
op pro argmem access ();
                              if (invoke info ptr != OPC NIL)
                                         have an invocation from IP. What
kind thou? */
                                    switch (invoke info ptr->invocation type)
                                          case MipC_Invoke_Type_IRDP:
```

```
have an IRDP packet. But which kind thou? */
                                                 if (invoke info ptr-
>irdp type == IcmpC Type IRDP Sol)
                            do not want to deal with this packet. */
                                                       mip sup pk cleanup
(invoke info ptr);
                                                       is ip pk = OPC TRUE;
                                                       break;
                        /* This must be an advertisement from an agent. */
                                                 op pk nfd get
(invoke info ptr->pk ptr, "data", &irdp pkptr);
                                                 mip mn ad packet parse
(irdp pkptr, &h bit, &f bit, &tmp agent address, &irdp lifetime,
&pref level);
      /\star For now, comaprison of the only address in the packet suffices. \star/
(inet address equal(tmp agent address, ha address))
                                                       if (loopback_intf)
                                           /* This is not supported. */
                                                             op sim end
("MR/MN when configured on a loopback interface, cannot directly communicate
with HA.",
                                                                   OPC NIL,
OPC NIL, OPC NIL);
                                                             }
                                     /* Ad from my home agent. */
                                                       is ha ad reception =
OPC TRUE;
                        /* Initialize counter before start reg process. */
                                                       retry counter = 0;
                        /* Update the latest ha info structure for later. */
      mip mn agent cache update (&latest ha info, tmp agent address,
                                                              (double)
irdp lifetime, pref level, invoke info ptr);
                                     /* Deregister with HA. */
                                                       mip mn register (0,
latest ha info, OPC TRUE);
                                           /* Update timer for HA timeout. */
     mip mn agent timer update (latest ha info.lifetime);
                                                 else
                                                       if (f bit)
```

```
/* A foreign agent ad. */
(current_agent_pref level < pref level)</pre>
            /* New FA agent advertising has higher preference.
                                                                Switch. */
      is switch fa = OPC TRUE;
                                                             else
      is fa ad reception = OPC TRUE;
                                    /* Update the FA cache information. */
      mip mn agent cache update (&latest fa info, tmp agent address,
                                                                    (double)
irdp lifetime, pref level, invoke info ptr);
                              /* Check to see if it is the same agent. */
(inet address equal (agent address, tmp agent address))
                                          /* Update timer for FA timeout. */
      mip mn agent timer update (op sim time () + (double) irdp lifetime);
                                                       else
/st cannot do anything with this agent who is only HA for other group. st/
                                                             is ip pk =
OPC TRUE;
                                                       }
                                                 /* Clean up. */
                                                 mip sup pk cleanup
(invoke info ptr);
                                                 op pk destroy (irdp pkptr);
                                                 break;
                                           case MipC_Invoke_Type_Tunnel_Check:
                                                 is ip pk = OPC TRUE;
                      should try to decapsulte packet if in MR mode. */
                                                 if (mip node type ==
MipC Node Type MR)
      /* will check first if this packet is tunneling other IP packet */
(mip sup ip in ip decapsulate (invoke info ptr->pk ptr, &encap pk ptr)
```

```
OPC COMPCODE SUCCESS)
                                          /* Sanity check on the packet */
                              /* Invoke IP delayed to handle the packet */
      mip sup packet send to ip (module data, encap pk ptr);
                        /* Write stats for the received tunneled packet. */
      mip mn tunneled pk stat write (invoke info ptr->pk ptr);
                  /* Let IP caller know that are handling the packet */
      mip sup pk cleanup (invoke info ptr);
                                                break;
                                          case MipC Invoke Type IP Datagram:
                                                 is ip pk = OPC TRUE;
                                                 /\star The destination address
that is going out on this interface when in MR mode should be forwarded to
either HA or FA. Whoever already registered or trying to. */
                                    if (mip_node_type == MipC_Node_Type_MR)
                                          /* Handle packet if I am a MR. */
                                    mip mn ip pk handle (invoke info ptr);
                                                 else
                        /* Unknown address. Send to the current agent. */
                                                      invoke info ptr-
>rte info ici ptr->next addr = agent address;
                                                break;
                                          }/* switch (invoke info ptr-
>invocation type) */
                              else
                                    /* Reg packet arrival. */
                                    op ici destroy (op intrpt ici ());
                                    is reg pk = OPC TRUE;
                        FSM PROFILE SECTION_OUT (state3_exit_exec)
```

```
/** state (Away) transition processing **/
                 FSM PROFILE SECTION IN ("mobile ip mn [Away trans
conditions]", state3 trans conds)
                 FSM INIT COND (FA TIMEOUT || REREGISTER || SWITCH FA)
                 FSM TEST COND (IP PK)
                 FSM TEST COND (FA AD RECEPTION)
                 FSM TEST COND (REG PK)
                 FSM TEST COND (HA AD RECEPTION)
                 FSM TEST LOGIC ("Away")
                 FSM PROFILE_SECTION_OUT (state3_trans_conds)
                 FSM TRANSIT SWITCH
                       FSM_CASE_TRANSIT (0, 5, state5_enter_exec, ;,
"FA TIMEOUT || REREGISTER || SWITCH FA", "", "Away", "Check FA cache",
"tr 4", "mobile ip mn [Away -> Check FA cache : FA TIMEOUT || REREGISTER ||
SWITCH FA / ]")
                       FSM CASE TRANSIT (1, 3, state3 enter exec, ;,
"IP PK", "", "Away", "Away", "tr_10", "mobile_ip_mn [Away -> Away : IP_PK /
]")
                       FSM CASE TRANSIT (2, 3, state3 enter exec, ;,
"FA AD RECEPTION", "", "Away", "Away", "tr 27", "mobile_ip_mn [Away -> Away :
FA AD RECEPTION / ]")
                       FSM CASE TRANSIT (3, 3, state3 enter exec, ;,
"REG PK", "", "Away", "Away", "tr_37", "mobile_ip_mn [Away -> Away : REG_PK /
                      FSM CASE TRANSIT (4, 4, state4 enter exec, ;,
"HA AD RECEPTION", "", "Away", "Pending registration", "tr_45", "mobile_ip_mn
[Away -> Pending registration : HA AD RECEPTION / ]")
           /*----*/
                  /** state (Pending registration) enter executives **/
                 FSM STATE ENTER UNFORCED (4, "Pending registration",
state4 enter exec, "mobile ip mn [Pending registration enter execs]")
                       FSM PROFILE SECTION IN ("mobile ip mn [Pending
registration enter execs]", state4 enter exec)
                       /* Update the status for debugging. */
                       mip sup mn mr status update (node objid,
home address, MipC Mn Mr Status Pending,
                             ha address, agent address);
                       FSM PROFILE SECTION OUT (state4 enter exec)
                  /** blocking after enter executives of unforced state. **/
                 FSM EXIT (9, "mobile ip mn")
                 /** state (Pending registration) exit executives **/
                 FSM STATE EXIT UNFORCED (4, "Pending registration",
"mobile ip mn [Pending registration exit execs]")
                       FSM PROFILE SECTION IN ("mobile ip mn [Pending
registration exit execs]", state4_exit_exec)
                       /* Who invoked me? */
```

```
op pro invoker (proc info struct ptr->pro hndl,
&inv mode);
                        if (inv mode == OPC PROINV DIRECT)
                              {
                              /* one of those timer went off. */
                              is timeout = OPC TRUE;
                        else
                              /* See if are getting an IP packet. */
                              invoke info ptr = (MipT Invocation Info*)
op pro argmem access ();
                              if (invoke_info ptr != OPC NIL)
                             have an invocation from IP. What kind thou? */
                                    switch (invoke info ptr->invocation type)
                                          case MipC Invoke Type IRDP:
                             have an IRDP packet. But which kind thou? */
                                                 if (invoke info ptr-
>irdp type == IcmpC Type IRDP Sol)
                                   do not want to deal with this packet. */
                                                      mip sup pk cleanup
(invoke info ptr);
                                                       is_ip_pk = OPC TRUE;
                                                      break;
                        /* This must be an advertisement from an agent. */
                                                is ad reception = OPC TRUE;
                      have an agent advertisement. But which kind thou? */
                                                op_pk_nfd_get
(invoke info ptr->pk ptr, "data", &irdp pkptr);
                                                mip mn ad packet parse
(irdp pkptr, &h bit, &f bit, &tmp_agent_address, &irdp_lifetime,
&pref level);
      /* For now, comaprison of the only address in the packet suffice. */
(inet address equal(tmp agent address, ha address))
                                                       if (loopback intf)
                                          /* This is not supported. */
                                                             op sim end
("MR/MN when configured on a loopback interface, cannot directly communicate
with HA.",
                                                                   OPC NIL,
OPC NIL, OPC NIL);
                                                             }
                        /* Update the latest ha info structure for later. */
```

```
mip mn agent cache update (&latest ha info, tmp agent address,
irdp lifetime, pref level, invoke info ptr);
      /* Check if are currently trying to register directly with HA. */
                                                       if (!direct reg)
      op ev cancel if pending (reg retry timer ehndl);
                                                 /* Deregister with HA. */
                                                            mip mn register
(0, latest ha info, OPC TRUE);
                                          /* Update timer for HA timeout. */
      mip mn agent timer update (latest ha info.lifetime);
                                                 else
                                                       if (f bit)
                                    /* Update the FA cache information. */
      mip mn agent cache update (&latest fa info, tmp agent address,
                                                                   (double)
irdp lifetime, pref level, invoke info ptr);
                              /* Check to see if it is the same agent. */
(inet address equal (agent address, tmp agent address))
                                          /* Update timer for FA timeout. */
      mip mn agent timer update (op sim time () + (double) irdp lifetime);
                                                       else
/* cannot do anything with this agent who is only HA for other group. */
                                                       }
                                                 /* Clean up. */
                                                mip sup pk cleanup
(invoke info ptr);
                                                 op pk destroy (irdp pkptr);
                                                break;
                                          case MipC Invoke Type Tunnel Check:
                                                 is_ip_pk = OPC TRUE;
                      should try to decapsulte packet if in MR mode. */
```

```
if (mip node type ==
MipC Node Type MR)
      /* will check first if this packet is tunneling other IP packet */
(mip sup ip in ip decapsulate (invoke info ptr->pk ptr, &encap pk ptr)
OPC COMPCODE SUCCESS)
                                          /* Sanity check on the packet */
                              /* Invoke IP delayed to handle the packet */
      mip sup packet send to ip (module data, encap pk ptr);
                  /* Write stats for the received tunneled packet. */
      mip mn tunneled pk stat write (invoke info ptr->pk ptr);
                                                             /* Let IP caller
            are handling the packet */
know that
      mip sup pk cleanup (invoke info ptr);
                                                break;
                                          case MipC_Invoke_Type_IP_Datagram:
                                                 is_ip_pk = OPC TRUE;
                                                 /* The destination address
that is going out on this interface when in MR mode should be forwarded to
either HA or FA. Whoever already registered or trying to. */
                                                 if (mip node type ==
MipC Node Type MR)
                  /* Handle packet if I am a MR. */
                                                      mip mn ip pk handle
(invoke info ptr);
                                                 else
                        /* Unknown address. Send to the current agent. */
                                                       invoke info ptr-
>rte info ici ptr->next addr = agent address;
                                                       }
                                                break;
                                          } /* switch (invoke info ptr-
>invocation type) */
                              else
```

```
/* Registration packet arrival. */
                                    is reg pk = OPC TRUE;
                        FSM PROFILE SECTION OUT (state4 exit exec)
                  /** state (Pending registration) transition processing **/
                  FSM PROFILE SECTION IN ("mobile ip_mn [Pending registration
trans conditions]", state4 trans conds)
                  FSM_INIT_COND ( REG_PK)
                  FSM TEST COND (TIMEOUT)
                  FSM TEST COND (AD RECEPTION)
                  FSM TEST COND (IP PK)
                  FSM TEST LOGIC ("Pending registration")
                  FSM PROFILE SECTION OUT (state4 trans conds)
                  FSM TRANSIT SWITCH
                        FSM CASE TRANSIT (0, 6, state6 enter exec, ;, "
REG_PK", "", "Pending registration", "Handle registration", "tr_15",
"mobile_ip_mn [Pending registration -> Handle registration : \overline{\text{REG}} PK / ]")
                        FSM CASE TRANSIT (1, 7, state7 enter exec, ;,
"TIMEOUT", "", "Pending registration", "Handle timeout", "tr_19",
"mobile_ip_mn [Pending registration -> Handle timeout : TIMEOUT / ]")
                        FSM_CASE_TRANSIT (2, 4, state4_enter_exec, ;,
"AD RECEPTION", "", "Pending registration", "Pending registration", "tr 28",
"mobile ip mn [Pending registration -> Pending registration : AD RECEPTION /
]")
                        FSM CASE TRANSIT (3, 4, state4 enter exec, ;,
"IP PK", "", "Pending registration", "Pending registration", "tr 11",
"mobile ip mn [Pending registration -> Pending registration : IP PK / ]")
                            -----*/
                  /** state (Check FA cache) enter executives **/
                  FSM STATE ENTER FORCED (5, "Check FA cache",
state5 enter exec, "mobile ip mn [Check FA cache enter execs]")
                        FSM PROFILE SECTION IN ("mobile ip mn [Check FA cache
enter execs]", state5 enter exec)
                        /* See if can use the cached FA information. */
                        if (inet address valid (latest fa info.address))
                              if (latest fa info.lifetime > op sim time () )
                        /* Initialize counter before start reg process. */
                                    retry counter = 0;
                        /* Register with the latest advertised FA. */
                                    mip mn register (reg info.req lifetime,
                                          latest fa info, OPC FALSE);
                                    /* Update the timer timeout. */
                                    if (!HA TIMEOUT && !FA TIMEOUT)
                        mip_mn_agent_timer_update (latest_fa_info.lifetime);
```

```
else
                                    /* cannot cancel the current event. */
      agent timer ehndl = op intrpt schedule self (latest fa info.lifetime,
                                                MipC MN Timer Agent);
                                    is valid fa candidate = OPC TRUE;
                        if (!is valid fa candidate)
                              if (solicitation)
                  /* Schedule an interrupt to send solicitation packet. */
            op intrpt schedule self (op sim time (), MipC MN Timer Solicit);
                  /* Cancel reregister timer first. */
                        op ev cancel if pending (reregister timer ehndl);
                              if (op sim anim ())
                                    /* Erase the existing tunnel. */
                                    mip sup draw tunnel (node objid,
ha address, ha address, ((mip node type == MipC Node Type MR) ? OPC TRUE :
OPC FALSE));
                        FSM PROFILE SECTION OUT (state5 enter exec)
                  /** state (Check FA cache) exit executives **/
                  FSM STATE EXIT FORCED (5, "Check FA cache", "mobile_ip_mn
[Check FA cache exit execs]")
                  /** state (Check FA cache) transition processing **/
                  FSM PROFILE SECTION IN ("mobile ip mn [Check FA cache trans
conditions]", state5 trans conds)
                  FSM INIT COND (!VALID FA CANDIDATE)
                  FSM TEST COND (VALID FA CANDIDATE)
                  FSM TEST LOGIC ("Check FA cache")
                  FSM PROFILE SECTION OUT (state5 trans conds)
                  FSM_TRANSIT_SWITCH
                       FSM_CASE_TRANSIT (0, 1, state1_enter_exec, ;,
"!VALID FA CANDIDATE", "", "Check FA cache", "Lost", "Tr 6", "mobile ip mn
[Check FA cache -> Lost : !VALID FA CANDIDATE / ]")
                        FSM CASE TRANSIT (1, 4, state4_enter_exec, ;,
"VALID FA CANDIDATE", "", "Check FA cache", "Pending registration", "tr 7",
"mobile ip mn [Check FA cache -> Pending registration : VALID FA CANDIDATE /
]")
```

```
/** state (Handle registration) enter executives **/
                  FSM STATE ENTER FORCED (6, "Handle registration",
state6 enter exec, "mobile ip mn [Handle registration enter execs]")
                        FSM PROFILE SECTION IN ("mobile ip mn [Handle
registration enter execs]", state6 enter exec)
                        /* Access ICI from mobile ip module. */
                        reg ici ptr = op intrpt ici ();
                        /* Get the ICI values. */
                        op ici attr get (reg ici ptr, "reply code",
&reply code);
                        op ici attr get (reg ici ptr, "lifetime grant",
&lifetime grant);
                        op ici attr get (reg ici ptr, "identification",
&tmp reg id);
                        if (tmp reg id == reg id)
                              /* Did it go through? */
                              if ((reply code == MipC Reg Reply Code Accept)
II
                                     (reply_code ==
MipC Reg Reply Code Accept No Simultaneous Binding))
                                     /* Cancel the retry timer first. */
                                    op ev cancel if pending
(reg retry timer ehndl);
                                    if (direct reg)
                                           is_ha_reg_success = OPC TRUE;
                                           if (MIP TRACE)
                                                 op prg odb print major
("Registering directly with HA successful.", OPC NIL);
                                           if (op sim anim ())
                                                 /* Draw tunnel to the HA. */
                                                 mip sup draw tunnel
(node objid, ha address, agent address, ((mip node type == MipC Node Type MR)
? OPC TRUE : OPC FALSE));
                                    else
                                          is fa reg success = OPC TRUE;
                                     /* Cancel reregister timer first. */
                                           op ev cancel if pending
(reregister timer ehndl);
```

```
time to reregister = op sim time ()
+ (double) lifetime grant - MipC MN Rereg Buffer;
                                           reregister timer ehndl =
op intrpt schedule self (time to reregister,
                                                 MipC MN Timer Rereg);
                                           if (MIP TRACE)
                                                 {
                                                 op prg odb print major
("Registering via a FA successful.", OPC NIL);
                                           if (op sim anim ())
                                     /* Draw tunnel to the HA through FA. */
                                                 mip sup draw tunnel
(node objid, ha address, agent address, ((mip node type == MipC Node Type MR)
? OPC TRUE : OPC FALSE));
                                     }
                              else
                                     /* have to retry. */
                                     if (retry counter <= reg info.retry)</pre>
                                           is retry = OPC TRUE;
                                     else
                                           op ev cancel if pending
(reg retry timer ehndl);
                                           is out of retries = OPC TRUE;
                                     }
                        else
                              /* Identification mismatch. */
                              is invalid reply = OPC TRUE;
                        /* Clean up. */
                        op ici destroy (reg ici ptr);
                        FSM PROFILE SECTION OUT (state6 enter exec)
                  /** state (Handle registration) exit executives **/
                  FSM STATE EXIT FORCED (6, "Handle registration",
"mobile ip mn [Handle registration exit execs]")
                  /** state (Handle registration) transition processing **/
                  FSM PROFILE SECTION IN ("mobile ip mn [Handle registration
trans conditions]", state6 trans conds)
                  FSM INIT COND (FA REG SUCCESS)
                  FSM TEST COND (OUT OF RETRIES)
```

/* Update the lifetime. */

```
FSM TEST COND (RETRY || INVALID REPLY)
                   FSM TEST COND (HA REG SUCCESS)
                   FSM TEST LOGIC ("Handle registration")
                   FSM PROFILE SECTION OUT (state6 trans conds)
                   FSM_TRANSIT_SWITCH
                        FSM CASE TRANSIT (0, 3, state3 enter exec, ;,
"FA_REG_SUCCESS", "", "Handle registration", "Away", "tr_2", "mobile_ip_mn [Handle registration -> Away : FA_REG_SUCCESS / ]")
FSM_CASE_TRANSIT (1, 5, state5_enter_exec, ;, "OUT_OF_RETRIES", "", "Handle registration", "Check FA cache", "tr_12",
"mobile ip mn [Handle registration -> Check FA cache : OUT OF RETRIES / ]")
                        FSM_CASE_TRANSIT (2, 4, state4_enter_exec, ;, "RETRY
| INVALID REPLY", "", "Handle registration", "Pending registration",
"tr 18", "mobile ip mn [Handle registration -> Pending registration : RETRY
|| INVALID REPLY / ]")
                        FSM_CASE_TRANSIT (3, 2, state2_enter_exec, ;,
"HA REG SUCCESS", "", "Handle registration", "At home", "tr_44",
"mobile ip mn [Handle registration -> At home : HA REG SUCCESS / ]")
            /*----*/
```

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